Incremental Evaluation of Schema-directed XML Publishing

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Abstract

When large XML documents published from a database are maintained externally, it is inefficient to repeatedly recompute them when the database is updated. Vastly preferable is incremental update, as common for views stored in a data warehouse. However, to support schema-directed publishing, there may be no simple query that defines the mapping from the database to the external document. To meet the need for efficient incremental update, this paper studies two approaches for incremental evaluation of ATGs [4], a formalism for schema-directed XML publishing. The reduction approach seeks to push as much work as possible to the underlying DBMS. It is based on a relational encoding of XML trees and a nontrivial translation of ATGs to SQL 99 queries with recursion. However, a weakness of this approach is that it relies on high-end DBMS features rather than the lowest common denominator. In contrast, the build-up approach pushes only simple queries to the DBMS and performs the bulk of the work in middleware. It capitalizes on the tree-structure of XML views to minimize unnecessary recomputations and leverages optimization techniques developed for XML publishing. While implementation of the reduction approach is not yet in the reach of commercialresizing, we have implemented the build-up approach and experimentally evaluated its performance compared to recomputation.

1. Introduction

XML publishing by middleware [11, 8, 16] or with direct DBMS support [7] has been well studied, and techniques from this work are rapidly being introduced into commercial products [25, 28]. In some applications, small portions of a database are extracted into "digestible" XML documents, for example the messages needed to execute or respond to requests using a web-services protocol. However, in many applications including mediation, archiving and web site management, large XML documents may need to be exported. In this case, the cached XML document can obviously be seen as a view of the database instance. For all the reasons that efficient incremental maintenance of views in the database (see, e.g., [14]) is important, it may also make sense to incrementally update published XML documents, even when they are externally cached by a middleware system. However, to

\[ I \vdash T \]

\[ \sigma \]

\[ \Delta T \]

\[ I \land \Delta T \]

\[ T \land \Delta T : D \]

Figure 1: Incremental evaluation

our knowledge, no previous work has considered incremental update of XML documents published from relational data.

In this paper, we consider the particular case of incremental update of XML documents created by schema-directed XML publishing middleware. The idea of schema-directed publishing is to extract data from a relational database and construct an XML view that conforms to a predefined schema, in this case a DTD. The need for this is evident in practice: enterprises typically agree on a common schema for data exchange; thus, a specification for publishing views that ensures schema conformance is an obvious benefit for developers. In response to this need, the ATG (Attribute Translation Grammar) formalism has been developed [4]. An ATG is a mapping \( \sigma : R \rightarrow D \) associated with a relational schema \( R \) and a predefined (possibly recursive) DTD \( D \). Given an instance \( I \) of \( R \), \( \sigma \) computes an XML view \( T = \sigma (I) \) such that \( T \) conforms to \( D \). To accomplish this transformation, \( D \) is treated as a set of productions. Each production is annotated by \( \sigma \) with a set of semantic rules, and a single semantic attribute is defined for each element type. These rules govern the production of child elements for each element based on the data present in the attribute.

To discuss incremental update, we assume that a set of changes to an XML tree \( T \) can be encapsulated as \( \Delta T \), and that an operator \( \oplus \) represents the application of these updates. Given these assumptions, the incremental evaluation problem for ATGS can be stated as follows: given an ATG \( \sigma : R \rightarrow D \), a relational instance \( I \) of \( R \), the XML view \( T = \sigma (I) \), and changes \( \Delta T \) to \( I \), compute XML changes \( \Delta T \) to \( T \) such that \( T \oplus \Delta T = \sigma (I \land \Delta T) \). These relationships are illustrated in Fig. 1. Note further that, if \( I \) is correct, \( T \oplus \Delta I \) is trivially guaranteed to conform to the predefined DTD \( D \). In contrast to recomputing the new view from scratch, incremental evaluation of ATGS can, in principle, improve performance substantially by applying only the changes \( \Delta T \) to the old view \( T \). To realize the improvement requires a) an efficient algorithm that computes the XML changes \( \Delta T \) in response to relational changes \( \Delta I \), and b) an efficient implementation of the tree-update operator \( \oplus \).

Our first contribution is a reduction framework for incremental maintenance of XML views defined by ATGS. For a given ATG definition \( \sigma \), the framework consists of a) a set of virtual relations associated with the semantic attributes of \( \sigma \), b) queries defining the recursive relationship of these virtual relations to encode ATG-produced XML documents c) a mapping based on these recursive queries from ATGS to SQL 99 queries with linear recursion [24], and d) mechanisms for computing XML changes \( \Delta T \) from relational changes \( \Delta I \)
and for updating the external view $T$ with $\Delta T$.

Previous work [11, 8] for XML publishing has developed mappings from XML view queries to SQL queries that compute "root-to-leaf paths in a relational encoding of the XML tree. These techniques, however, cannot be directly applied to generating or updating AVG-generated documents due to a lack of support for recursion. In the context of XML sharding, the authors of [17] showed that linear recursion of SQL 99 is sufficient to support XPath queries over shredded XML data, even when the shredding schema is recursive, and suggested that this approach could be applied for publishing queries. The reduction framework extends the prior work by showing the connection between the recursive XML views needed for schema-directed publishing and SQL 99 views.

Much as [11, 8] seek to push XML publishing work to the DBMS, a primary goal of the reduction framework is to push incremental work to the DBMSs, thereby taking advantage of sophisticated capabilities for query optimization, execution, and incremental view update. However, three practical issues complicate the use of DBMS resources in support of incremental work. The first issue involves the DBMS features required for the reduction approach. Middleware that seeks to work with a wide variety of products supporting the ODBC interface must take a "lowest common denominator" approach to the functionality required from the DBMSs.

But SQL 99 recursion and incremental update of views are separate, advanced features of only the most sophisticated commercial products; furthermore the reduction approach also depends on incremental update of materialized views defined using SQL 99 recursion; in other words, a combination of both features. Second, to effectively push down the work required to incrementally update an external view, one must have access not only to a materialized view, but more importantly, to a stream of updates to that view. One way to obtain this functionality would be to define triggers on the materialized view, but this is disallowed by at least one commercial DBMS with materialized view support, and needs not be supported in general for the same function well.

Finally, if the publishing queries are even mildly complex, the combined recursive queries may become extremely complex. As a result, they may not be effectively optimized by all platforms supporting with... recursive for the same reasons that not all DBMS platforms can effectively optimize complex non-recursive publishing queries [11].

In response to this, our second contribution is to propose an alternative approach, referred to as the bud-cut approach to incremental AVG evaluation, that requires less sophistication from the DBMS. Further, we develop certain optimization techniques which capitalize on the tree-structure of XML views. The bud-cut mechanism propagates relational changes to XML in three phases: generation, completion and garbage collection. The bud-cut generation phase determines the impact of $\Delta I$ on existing parent-child relations in the XML view, i.e., insertions (buds) and deletions (cuds), by evaluating a fixed number of incrementalized SQL queries. Following this, the bud completion phase iteratively computes newly inserted subtrees top-down by pushing SQL queries to the relational engine. Finally, deleted subtrees are removed by a garbage collection process.

The bud-cut approach has several properties. a) $T$ can be updated in parallel with ongoing computation of $\Delta T$ during bud completion. b) It minimizes unnecessary recomputation via a caching strategy not considered in prior work for maintaining recursive views such that each new subtree in the XML view is computed at most once no matter how many times it occurs in the XML view, and furthermore, the computation maximally reuses subtrees of the old XML view. c) It incorporates optimization techniques that have proved effective in XML publishing but are not supported by DSSs, e.g., query merging [11]. d) Since the tree is computed level-by-level in this phase, it is possible to return partial results to a user navigating the tree while computation is ongoing, and such computation can even be deferred according to a lazy evaluation strategy [7]. e) It does not require materialization of the view in the DBMS. Finally, of course the bud-cut approach does not require the DBMS to support either SQL 99 or incremental view updates.

We have implemented the bud-cut approach. We use the implementation to investigate the impact of $\Delta I$ and $\Delta T$ on the performance of incremental update, as well as the improvement obtained over full recomputation for small updates. Further, we investigate the impact on performance of the subtree reuse optimization mentioned above, and find that its impact is greatest when only a portion of $I$ is published in $T$, and when there is a moderate degree of locality in the updates appearing in $\Delta I$. Unfortunately but not surprisingly, the reduction approach is unrealistic since current commercial relational systems do not implement incremental update of recursive queries.

The algorithms and the bud-cut mechanism can be extended to accommodate multiple data sources, i.e., for XML integration studied by [3], which is a generalization of AVG's by supporting multi-source SQL queries and XML constraints. They can also be used for incremental maintenance of XML views generated by other systems, such as [11, 8].

Organization. Section 2 reviews AVGs. Section 3 describes data structures for external XML views. Section 4 provides the reduction approach, followed by the bud-cut approach in Sect. 5. Section 6 presents experimental results. Section 7 addresses related work and Sect. 8 concludes the paper.

2. Background

In this section we first introduce a running XML publishing example used in the rest of the paper. We then review DTDs and present a refinement of AVGs as defined by [4].

Example 2.1: Consider a registrar database specified by the relational schema $R_0$ below (with keys underlined):

course(cno, title, dept), project(cno, title, dept), student(sno, name), enroll(sno, cno), prereq(cno, cno2).

The database maintains student data, enrollment records, course data classified into regular courses and projects, and a relation prereq which gives the prerequisite hierarchy of courses where a tuple $(c1, c2)$ in prereq indicates that $c2$ is a prerequisite of $c1$.

The office of registrar maintains an XML view for the CS department, which contains data of CS courses registration, extracted from the registrar database. The view is required to conform to the DTD $D_0$ below (the definition of elements whose type is PCDATA is omitted):

```
<ELEMENT db (course*)>
<ELEMENT course (cno, title, type, prereq, takenby*)>
<ELEMENT type (regular | project)>
<ELEMENT prereq (course*)>
<ELEMENT takenby (student*)>
```
The idea of attribute translation grammars (ATGs) is to treat the DTD as a grammar and recursively fire productions from the grammar to create an XML document. We now briefly review the syntax and semantics of ATGs (see [4]).

An ATG \( \sigma : R \rightarrow D \) specifies a mapping from instances of the source relational schema \( R \) to documents of the target DTD as follows: a) For each element type \( A \) of \( D \), \( \sigma \) defines a semantic attribute \( \$A \) whose value is a single relational tuple of a fixed arity and type; intuitively, \( \$A \) controls the generation of \( A \) elements in the XML view, and is used to pass data downward as the document is produced. b) For each production \( p = A \rightarrow \alpha \) in \( D \), \( \sigma \) specifies a set of semantic rules, \( \text{rule}(p) \). These rules specify the computation of the \( B \) children of an \( A \) element for each type \( B \) in \( \alpha \).

Given a database \( I \) of \( R \), the ATG \( \sigma \) is evaluated top-down starting at the root \( r \) of \( D \). A partial tree \( T \) is initialized with a single node of type \( r \), and this node is marked unexpanded; we refer to unexpanded nodes as buds. The tree \( T \) is then grown by repeatedly selecting a bud \( b \) (of some element type \( A \)), evaluating the semantic rules associated with \( A \) and marking \( b \) expanded. Specifically, we find the production \( p = A \rightarrow \alpha \) in \( D \), and the generate the children of \( b \) by evaluating \( \text{rule}(p) \) and using the value of the attribute \( \$A \) of \( b \). The rules \( \text{rule}(p) \) are defined and evaluated based on the form of \( \alpha \) as follows:

1. If \( \alpha = B_1 \ldots B_n \) then a node tagged \( B_i \) is created for each \( i \in [1,n] \) as a child of \( b \). The tuple value of \( B_i \) associated with the new \( B_i \) child is determined by projection from \( \$A \). That is, \( B_i = (\$A.a_i^1, \ldots, \$A.a_i^n) \) in \( \text{rule}(p) \) for \( i \in [1,n] \), where \( a_i^j \) is a field of the tuple \( \$A \).
2. If \( \alpha = B_1 + \ldots + B_n \) then \( \text{rule}(p) \) is defined by \( (B_1, B_2, \ldots, B_n, v) \) and \( \text{case } f(\$A) \) of

\[
\begin{array}{l}
1: (\$A, \text{null}, \ldots, \text{null}) \\
\cdots \\
n: (\text{null}, \ldots, \text{null}, \$A)
\end{array}
\]

where \( f \) is a function that maps \$A to natural numbers in \([1,n]\). That is, based on the conditional test, a node is created for exactly one child \( B_i \). The value of the parent attribute \$A is passed down to that child. No \( B_i \) child is created if \( i \neq j \) and \( B_j \) (the special value \text{null} is ignored. We assume that the function \( f \) is simple enough to determine whether it is in the range \([1,n]\).

3. If \( \alpha = B^* \), then \( \text{rule}(p) \) is defined by \( B \rightarrow Q(\$A) \), where \( Q \) is an \text{sql} query over a database of \( R \), and it treats \$A as a constant parameter. For each distinct tuple \( t \) returned by \( Q(\$A) \), a \( B \) child is generated, carrying \( t \) as the value of its \$B attribute. To help ensure that only finite documents are created, only references to attributes and constants, but not expressions are allowed in the select-list of \( Q \).

4. If \( \alpha = \text{str} \) then the rule specifies formatting of the values of \$B for presentation (string/PCDATA). Such rules are not shown or discussed further.

5. If \( \alpha = e \), then \( \text{no } r(p) \) is defined and no action is taken.

The element children of the node \( b \) become new buds and are also processed. The process proceeds until the partial tree cannot be further expanded, i.e., it has no unexpanded node. The fully expanded XML tree do not include attributes \$A, which are only used to control the tree generation.

Example 2.2: The ATG \( \sigma_0 \) given in Fig. 3 defines the XML view described in Example 2.1. Here \( \text{rule(course)} \).
$\text{db} \rightarrow \text{course}$
$\text{Course} \leftarrow \text{Q}_1$
$\text{Q}_1$: select distinct c.cno, c.title, 1 as tag from course c where c.dept = '__CS'
union
select distinct p.cno, p.title, 2 as tag from project p where p.dept = '__CS'

$\text{course} \rightarrow \text{cno, title, type, prereq, takenBy}$
$\text{Course} \rightarrow \text{Course, cno}$
$\text{Student} \rightarrow \text{Student, cno}$
$\text{StakenBy} \leftarrow \text{Course, cno}$
$\text{type} \rightarrow \text{regular + project}$
$\text{Regular, Project} = \text{case type of}$
$1: (\text{type}, \text{null})$
$2: (\text{null}, \text{type})$

$\text{prereq} \rightarrow \text{course}$
$\text{Course} \leftarrow \text{Q}_2[\text{prereq}]$
$\text{Q}_2[c]:$ select distinct c.cno, c.title, 1 as tag from prereq p, course c
where p.cno = c and p.cno = c.cno

$\text{takenBy} \rightarrow \text{student}$
$\text{Student} \leftarrow \text{Q}_3[\text{StakenBy}]$
$\text{Q}_3[c]:$ select distinct s.ssn, s.name from enroll e, student s
where e.cno = c and e.ssn = s.ssn

![Figure 3: Example ATG σ](image)

Theorem 2.1: Let $\sigma : R \rightarrow D$ be an arbitrary ATG with exception handling. Then over all instances $I$ of $R$, $\sigma$ terminates and $\sigma(I)$ conforms to the DTD $D$. □

Theorem 2.1 shows that an ATG $\sigma : R \rightarrow D$ is actually a total function: given a database $I$ of $R$, $\sigma(I)$ computes an XML documents of $D$ that is unique up to reordering of $B$-elements produced by productions of the form $A \rightarrow B^*$. In the sequel we consider only ATGs with exception handling and refer to them simply as ATGs. Note that our incremental techniques also work on ATGs without exception handling.

3. External Trees

Incremental update makes sense only in the context of an externally maintained tree available to the client. A variety of potential implementations for the external tree exist, from a native storage system like Berkeley XML DB to an in-memory implementation of DOM. While in the current work we have used our own implementation of an in-memory tree in C++, we expect the data structures and algorithms we propose to apply to other implementations.

Node Identity. We assume that we can associate a compact, unique value with each tuple value taken on by a semantic attribute in $\sigma(I)$. We abstract away the implementation of this identity value by assuming, without loss of generality, the existence of a Skolem function $\text{genId}$(see e.g., [9]) that, given the tuple value of a semantic attribute $A$, computes $\text{id}_A$ that is unique among all identities associated with all semantic attributes (for example, it might encode the type and a unique value within that type).

Tree vs. Graph Representations. An important property of an ATG $\sigma : R \rightarrow D$ is that, for any database $I$ of $R$ and type $A$ of $D$, an $A$-element $T_A$ in the XML view $\sigma(I)$ is uniquely determined by the value of the semantic attribute $A$ at the root of $T_A$. Thus the ATG defines a function $\text{str}$ such that, given an element type $A$ and a value $t$ of $\text{str}(A, t)$ returns a subtree rooted at a node tagged $A$ and carrying $t$ as its attribute.
Since a subtree $st(A, S, A)$ may appear at different places in the XML view $\sigma(I)$, if the middleware system is managing the external view, it may be more efficient to represent $\sigma(I)$ as a graph such that a single copy of $st(A, S, A)$ is stored and shared by its multiple occurrences. Indeed, if $\sigma(I)$ is stored as a tree, for example by using an available implementation of DOM [2], the tree may be exponentially larger than the graph representation. In light of this the current implementation of ATGS adopts the graph representation, but supports client navigation on the graph as if it were a tree. The trade-off of the space efficiency is that the support of user navigation is complicated, as the path from a particular node to the root of the tree would be dependent on the route navigated; further, a mechanism must be provided to check for duplicates along this path to provide the exception handling semantics described above.

**Data Structures.** The external data structures used to represent the XML document are depicted in Fig. 4. The tree $T$ is stored in a hash index $H$ and a subtree pool $S$. Each entry of $H$ is of the form $(A, id_A, ptr)$, where $A$ is an element type, $id_A$ is the unique id of a value of $S$, and $ptr$ is a pointer to the root node of the subtree $st(A, S, A)$ in $S$. The subtree pool $S$ consists of entries $(A, id_A, L)$, where $(A, id_A)$ represents a node $v$ of $T$, and $L$ is a list $[(B_1, id_{B_1}), \ldots, (B_n, id_{B_n})]$ representing the children of $v$ such that each $(B_i, id_{B_i})$ is an entry in $H$. Observe that in the graph representation, there is a one-to-one mapping from $H$ entries to the nodes in $S$.

**Handling Updates.** In this paper, we allow $I$ to be any group of updates to the underlying DBMS that preserves the consistency of the database (integrity constraints). XML updates $\Delta I$ generated from $I$ by one of the techniques described in the next two sections are represented as $(E^+, E^-)$, where $E^+$ is a set of edges to be inserted into the tree $T$, and $E^-$ is a set of edges to be deleted from $T$. The edges are represented as $(id_A, id_B)$, where $id_A$ and $id_B$ are the ids of the parent $A$-element and the child $B$-element respectively.

The system processes insertions as follows for each tuple $(id_A, id_B)$ in $E^+$, 1) find the $H$ entry $(A, id_A)$ and the pointer to the node $(A, id_A, L)$ in the subtree pool $S$; 2) insert $(B, id_B)$ into $L$ if it is not already in $L$, and 3) if $H$ does not have an entry $(B, id_B)$ in $H$, create an entry in $H$ and a node $(B, id_B, [])$ in $S$ with an empty child-list $[]$. Note that if there is already an entry $(B, id_B)$ in $H$, the new edge is actually a cross edge to the existing $B$ node. A subtle issue concerns the order of the children list $L$ when a $B$ node is inserted into $L$. Note that $L$ is constrained by the production $A \rightarrow \alpha$: if $\alpha$ is $B_1, \ldots, B_n$, the elements of $L$ are in the same order as their element types in the production; and if $\alpha$ is $B^+$, $L$ elements are ordered by a default order on the tuple values of the $S$ attributes of these elements.

The system carries out deletions similarly: for each tuple $(id_A, id_B)$ in $E^-$, 1) find the node $(A, id_A, L)$ in the subtree pool via the $H$-index entry $H(A, id_A)$, and 2) cuts the edge $(id_A, id_B)$ by removing $id_B$ from the children list $L$. It should be noted that during the processing of relational updates $\Delta I$, no entries are physically deleted from $H$ because a) an entry may be shared by multiple subtrees of $T$, and b) as will be seen shortly, our incremental system minimizes unnecessary recomputations by reusing subtrees that have been disconnected from the tree by cuts.

**Garbage Collection.** Upon the completion of the computation of the new XML view $\sigma(I \oplus \Delta I)$, a garbage collection process runs in the background to remove $H$-entries and nodes in the subtree pool that are not linked to $\sigma(I \oplus \Delta I)$. Conceptually this can be done via a top-down traversal of the new XML tree, removing nodes that are no longer reachable from the root. In our implementation we associate with each $H$-entry a count keeping track of the number of nodes linked to it, and maintain the counters when processing insertions and deletions; the garbage collection process removes unused nodes, i.e., those $H$-entries with count $= 0$ and their corresponding nodes in the subtree pool.

### 4. Pushing Incremental Work to the DBMS

In this section we present our reduction mechanism for incremental evaluation of ATGS. The approach is based on 1) a relational encoding of XML trees via a set of (interrelated) virtual attribute and edge relations and 2) a translation of an ATGS $\sigma : R \rightarrow D$ to a set $M_{Q_e}$ of SQL 99 queries utilizing the with...recursive construct, which computes the attribute and edge relations of the XML views defined by $\sigma$.

Given this the reduction approach to incrementally maintaining an XML tree $T$ computed by ATGS $\sigma(I)$ works as follows: 1) encode $T$ with the attribute and edge relations, 2) map $\sigma$ to SQL 99 queries $M_{Q_e}$, 3) define an incrementally updated materialized view in the source DBMS for each query of $M_{Q_e}$, 4) in response to relational updates $\Delta I$, utilize the usual functionality to capture the incremental changes made to each of the materialized views which can be directly transformed into the XML changes $\Delta T$ (i.e., $E^+$ and $E^-$), and 5) propagate the changes $\Delta T$ to the external tree as described in the previous section. This leads to a convenient approach to incremental evaluation of ATGS by pushing as much work as possible to the underlying DBMS.

We next focus on the relational encoding of XML trees and translation of ATGS to SQL 99 queries.

#### 4.1 A Relational Encoding of XML Trees

An XML view $\sigma(I)$ is encoded via node (semantic attribute) and edge relations.
Figure 6: Edge relation generating queries.

Attribute Relations. The attribute relations are to capture the values taken on by the semantic attributes defined in \( \sigma(I) \). To avoid confusion, "attributes" of relational tables will be uniformly referred to as "columns".

Recall that \( \sigma \) associates with each element type \( T \) a tuple-like variable \( B \). For each such variable \( B \), let \( \text{gen}_B \) be a relation with columns matching the arity and type of \( B \), along with a column for \( \text{id}_B \) if an existing group of attributes does not serve this role. Further, in the context of a database instance \( I \), assume that \( \text{gen}_B \) is populated with all the (non-null) values taken on by \( B \) during an evaluation of \( \sigma \) on \( I \).

We define each attribute relation, \( \text{gen}_B \), in terms of a query \( Q_{\text{gen}_B} \) involving the other attribute relations and the relations of \( I \). To define \( Q_{\text{gen}_B} \), we first rewrite sqI queries embedded in \( \sigma \) to queries that take \( \text{gen}_A \) instead of a single tuple \( A \), as a parameter. Specifically, consider productions \( A \rightarrow a \) in which \( B \) appears on its right-hand side (rhs).

(1) For productions of the form \( A \rightarrow B’ \) with associated semantic rule \( S \rightarrow Q(A) \), \( B’ \in \text{gen}_B \) is the union, over all values of \( A \) in \( \text{gen}_A \) of \( Q(A) \). In a manner similar to the level-at-a-time processing of \([4]\), this query can be rewritten as \( Q_i \), which takes no parameters but additionally accesses \( \text{gen}_A \). This is accomplished, roughly by a) adding \( \text{gen}_A \) to the from list of \( Q \), and b) replacing references to \( A \) in \( Q \) with the corresponding references to \( \text{gen}_A \).

(2) If \( B \) appears as some \( B_i \) when \( a \) is \( B_1, \ldots, B_n \), then \( Q’ \) is a simple selection query from \( \text{gen}_A \) that projects fields according to \( f_i(A) \) (see Section 2).

(3) If \( B \) appears as some \( B_i \) in \( a \) is \( B_1 + \cdots + B_n \) then \( Q’ \) can be written as \( \text{select } * \text{ from } \text{gen}_A \text{ where } f_i(a) = i \).

Note that we assume further that \( f(a) \) is computable in the dialect of sqI used to express \( Q \).

It is now simple to generalize this construction to handle the case where \( B \) appears on the rhs of multiple rules in \( \sigma \). Suppose that it appears on the rhs of rules \( p_1, p_2, \ldots, p_m \), and that \( Q_i \) is defined per the discussion above. Then \( \text{gen}_A \) is formed by taking the (distinct) union of all \( Q_i \) queries.

For example, Fig. 5 shows the definition of two attribute-relation generating queries when the resulting construction is applied to the rhs of Fig. 3.

We denote by \( \text{AR}(\sigma) \) the set of all \( Q_{\text{gen}_A} \) queries for \( \sigma \).

Maintaining Edges. We next describe how to capture the edge relations of the xml view. Let \( \text{edge}_AB \) be a relation with two columns, \( \text{id}_A \) and \( \text{id}_B \). We create such a relation if \( B \) appears on the rhs of some production for \( A \). We overload the Skolem function \( \text{gen}_id \) described in Sect. 3 to compute the unique id \( \text{id} \) from a relational tuple \( a \).

We now discuss how to derive queries to define these edge relations in terms of attribute relations and base relations. As before, consider first productions of the form \( A \rightarrow B’ \), where \( B’ \leftarrow Q(A) \) is the associated semantic rule. In this case, \( \text{edge}_AB \) is the set of pairs \((a, b)\) such that \( a \in \text{gen}_A \) with \( Q1(\text{cmo.title,...}) \) as (Q1 from Figure 3) and \( \text{Q2(idc.idp, pq.idp, pq.idc, srel}) \)

with \( \text{Q1(\text{cmo.title,...}) as (Q1 from Figure 3)} \) and \( \text{Q2(idc.idp, pq.idp, pq.idc, srel}) \)

Figure 7: with ... recursive for recursive schema

and \( b \in Q(a) \), where \( ia = \text{gen}_id(a) \) and \( ib = \text{gen}_id(b) \). To derive a query \( Q_{\text{edge}_AB} \) for an edge relation \( \text{edge}_AB \), we can employ the same rewriting as for the attribute relations, with the following change: the select list of the attribute relation is replaced with \((ia, ib) \).

The definition of \( Q_{\text{edge}_AB} \) is similar for productions of other forms.

The set of all \( Q_{\text{edge}_A} \) queries for \( \sigma \) is denoted by \( ER(\sigma) \).

As an example, Fig. 6 shows two edge-relation generating queries derived from the on A of Fig. 3.

4.2 Computing Attribute and Edge Relations with sqI 99 Linear Recursion

As the running example illustrates, the attribute relations are potentially mutually recursive, and the edge relations depend on the attribute relations. Furthermore, each attribute relation can be related to itself and other attribute relations through a variety of paths, raising the possibility that concluding evaluation of updates to the attribute relations will be excessively complex. Fortunately, this is not the case.

While the attribute relations are recursive, there are substantial constraints on this recursion. Consider the references in the query \( Q_{\text{gen}_A} \) to other attribute relations \( B, C, \ldots \). First, observe that such a relationship cannot be made via negation (in the Datalog sense), that is, the reference to \( B \) cannot appear in a not exists clause in \( Q_{\text{gen}_A} \)

Second, note that references to two different attributes relations for example \( B \) and \( C \) must be made in two different subqueries of \( Q_{\text{gen}_A} \), such that these subqueries are joined only by the top-level union operator introduced by the construction. Finally since expressions are not allowed in the select clauses of view-definition queries, aggregate expressions which depend on attribute relations will not appear.

Given these restrictions, the computation of the attribute and/or edge relations can be accomplished with the sqI 99 with...recursive construct. Let \( Q_0 \) be a set of edge and attribute relation-generating queries from \( \sigma \) which either includes all of the attribute relations or all of the edge relations. Let \( G_0 \) be a graph in which there is a node for each query in \( Q_0 \), and an edge from \( Q_0 \) to \( Q_1 \) if \( Q_0 \) refers to the virtual relation \( \text{gen}_A \) in the from clause. Let \( c_1, \ldots, c_n \) be the connected components of this graph. We now merge components and generate possibly recursive queries along the same lines as \([17]\). First, singleton components with only a single incoming edge is merged with the source component. Let the result of this process be \( c_1, \ldots, c_n \). For each such component \( c_i \), if it is acyclic, then a set of merge-queries...
Figure 8: Middleware Architecture

\[ M_{Q_c}, \] for the component can be defined. The number of queries produced will depend on the choices made for query merging [4, 11], and DOM structures can be handled with an embedded with clause as suggested by [17].

If the component is cyclic, a single recursive query \( M_{Q_c} \) is defined to compute the outer union of all the virtual relations in \( c \) with the following steps:

1. The schema of \( M_{Q_c} \) is the union of all columns appearing in any query of \( c \) plus an additional \( \text{rel} \) attribute intended to encode which virtual relation a given tuple in \( M_{Q_c} \) represents.
2. The initial condition for \( M_{Q_c} \) is the union of the queries for nodes with edges incident on \( c \) in \( G_Q \).
3. The query \( M_{Q_c} \) is the union of all the queries of \( c \). In each subquery, if a mention to a virtual attribute relation \( gen_A \) appears it is replaced with a reference to \( M_{Q_c} \), and a conjunctive condition is added to the where clause to ensure that \( \text{rel} = \{ \text{"I"} \} \); similarly for edge relations.

Applying this algorithm to the ATG of Fig. 3, we get two components after merging. One component consists of \( Q_{edge\_predecessor} \) and \( Q_{old\_predecessor} \), and is shown as the linear-recursive QC2 in Fig. 7. All the remaining queries are in the other, which is not recursive.

We refer to the set of queries \( M_{Q_c} \) generated for an ATG \( \sigma \) as \( M_{Q_c}^{\sigma} \). It is easy to verify the correctness of the mapping: for any ATG \( \sigma : R \rightarrow D \) and any relational instance \( I \) of \( R \), \( M_{Q_c}(I) \) computes \( \text{gen}_A \) and \( \text{edge}_A'B \) for all attributes \$A and parent-child edges in the XML view of \( I \).

In response to relational updates \( \Delta I \), each edge relation \( \text{edge}_A'B \) defined by squ. 99 queries \( M_{Q_c} \) is updated via insertions and deletions. Insertions and deletions over all edge relations are collected into two sets \( E^+ \) and \( E^- \), respectively. The two sets \( E^+ \) and \( E^- \) are sent to the middleware maintaining the external tree, as described in the last section. As mentioned in the introduction, this approach assumes that the underlying DBMS incrementally computes edge changes \( E^+ \) and \( E^- \) and further that these incremental changes can be captured, for example with triggers.

5. Bud-Cut Incremental Evaluation

As discussed in the introduction, the reduction approach is not practical for middleware-based XML publishing since it depends on a combination of features not yet found even in the most advanced commercial DBMSs. While middleware should depend on only the most common features. This observation motivates us to propose the bud-cut approach to incrementally evaluating ATG \( \sigma : R \rightarrow D \), which does not require the underlying DBMS to support with \( ...\) recursive.

A middleware system based on the bud-cut mechanism is depicted in Fig. 8. The system interacts with an underlying DBMS and maintains a hash index \( H \) and a subtree pool for the external xml view \( T \) (\( \sigma(I) \)) as described in Sect. 3. It responds to a relational update \( \Delta I \) in three phases. The first phase bud-cut generation identifies the portions of the existing tree that will be affected by the updates \( \Delta I \) and propagates \( \Delta I \) to XML changes to the existing nodes in \( T \). Nodes created in this phase are \textit{buds}, \textit{i.e.}, they are marked unexpanded; and nodes to be removed are \textit{cuts}. The second phase, bud completion, constructs subtrees under buds (and generates new buds), taking advantage of ATG properties to avoid recomputation and reusing existing subtrees when possible. The third phase, \textit{garbage collection}, runs after bud completion is finished and removes unreachable subtrees.

While requiring several round-trips between the middleware and the DBMS, this approach is able to exploit other optimizations by taking advantage of the specifics of ATG semantics and XML views. In particular, as observed in Sect. 3, since the value of a subtree \( \text{sv}(A, t) \), is determined by the tuple value \( t \) of the semantic attribute of the subtree’s root node, changes to the children of the existing nodes in the tree can be computed by a fixed set of non-recursive queries. Furthermore, any new subtree \( \text{sv}(A, t) \) can be reused and thus needs to be computed at most once. In addition, the bud-cut approach allows the update process of the external tree \( T \) to run in parallel with the computation of XML changes \( \Delta T \), thus improving the response time. This last point leads to a variety of options for \textit{lazy evaluation} of incremental updates. That is, the middleware can optionally defer complete processing of updates until the subtrees affected by those updates are accessed. When the external view is materialized as an in-memory tree, it is possible to support tree navigation concurrent with processing [7, 23].

5.1 Implementation of Delta SQL Queries

Assuming the existence of virtual attribute relations, the first step is to derive two sets of incremental (cos)recursive squ. queries from edge-generating queries \( E\!(\sigma) \). These two forms of incremental queries will be used in the bud-cut generation phase and the bud-completion phase, respectively.

Incremental queries of the first form are derived as follows. For each query \( Q_{edge\_A'B} \) in \( E\!(\sigma) \), we define a \textit{bud generating} query formed by adding all the columns of the symmetric \( \#B \) to the select list of \( Q_{edge\_A'B} \). An incremental form of this bud-generating query, \( \Delta Q_{edge\_A'B} \), is then created by using a counting method like [15]. In a nutshell, [15] associates a \textit{count} with each tuple in a view to keep track of the number of alternative derivations of the tuple via the view, and computes \textit{updates} to the view, namely, insertions and deletions of tuples, by incrementing or decrementing the counts of its tuples in response to changes to base relations.

Since the method of [15] is assumed to execute in the DBMS, it assumes access to both \( I \) and \( I \oplus \Delta I \). However, our middleware system is separate from the DBMS and can only access \( I \oplus \Delta I \). To find \( I \) for a given relation, \( R_i \), we assume...
Figure 10: Incremental query for bud-completion.

The existence of a change table $\Delta I$, that holds $\Delta I$ restricted to $R_i$ as well as a count cnt with the value either "+1" or "-1" to indicate an insert or delete. Further, we assume that $\Delta I$ has already been applied to $I_i$, the instance of $R_i$ in $I$.

In order to simulate the pre-image of $I_i$, we define a view $\text{old}_R$: for each relation $R_i \in R$ which simply selects tuples of $I_i$ that do not appear in $\Delta I$. It can be implemented efficiently if $\Delta I$ is small and is indexed on the key of $R_i$.

For example, Fig. 9 gives the incremental bud-generating query for $edge_{prereq,course}$, which is derived from the query $Q_{edge_{prereq,course}}$ by using the counting method of [15]. It computes updates to $edge_{prereq,course}$ in response to changes $A_{course}$ and $prereq$. Observe that in addition to edges of $edge_{A,B}$, the query also returns the corresponding $\#B$ values. A query $Q_{edge_{A,B}}$ is typically a union of several queries each incorporating the effects of changes in one of the base relations referenced by $Q_{edge_{A,B}}$.

Incremental queries of the second form are a mild extension of $ER(r)$. For each $Q_{edge_{A,B}}$ in $ER(r)$, we define an incremental bud-completing query $Q_{bud_{edge_{A,B}}}(Bud_A)$ to compute changes to $edge_{A,B}$ in response to a set $Bud_A$ of insertions to the attribute relation $gen_A$. The query is derived from $Q_{edge_{A,B}}$ by substituting $Bud_A$ for $gen_A$ and by adding all the columns of the attribute $\$B$ to the select list of $Q_{edge_{A,B}}$. It computes the edges of the newly inserted $A$-nodes in $Bud_A$. In contrast to the incremental bud-generating query $Q_{edge_{A,B}}$, the incremental query $Q_{bud_{edge_{A,B}}}(Bud_A)$ takes $Bud_A$ as a parameter and does not access $gen_A$; furthermore, it uses the new relation $I \otimes \Delta I$ without requiring the old $I$. For example, Fig. 10 shows the incremental bud-completing query of $Q_{bud_{edge_{A,B}}}(Bud_A)$, which assumes that course has been updated by $A_{course}$.

5.2 Bud-Cut Generation

Having derived the incrementalized SQL queries, this phase proceeds as follows:

1. The incremental bud-generating queries $Q_{edge_{A,B}}$ are executed in the DBMS, yielding changes to the edge relations a set of insertions $E^+$ and deletions $E^-$. The sets ($E^+, E^-$) are transmitted to the middleware and posted against the tree $T$, as described in Sect. 3. Moreover, the $\$A$-value of each new $A$-node added to $T$ is also added to $Bud_A$, a set of buds of type $A$. A $A$-node is marked unexpanded if its $\$A$ value is in the $Bud_A$ set.

2. The bud-cut approach does not require the attribute relations to be materialized in the DBMS since they can be computed from the hash table and the subtree pool in the middleware. To speed up the response time with a tradeoff of space, we treat these relations as views maintained by the DBMS. In the latter case, the changes computed in step (1) above also update these views; furthermore, in step (2) tuples in $Bud_A$ are sent to the DBMS and inserted into $gen_{A,A}$ for all element types $A$ in the ATG.

It should be remarked that the bud-cut-generation phase updates the parent-child edges of all the existing nodes in the XML view $T$ in response to relational changes $\Delta I$. Thus all the edge deletions are handled in this phase, by evaluating the incrementalized SQL (bud-generating) queries of $ER(r)$ once. Note further that nodes disconnected by the edge cuts are not removed from $H$-index or the tree but instead remain in the subtree pool pending reuse during bud generation.

Garbage collection of the nodes that have been disconnected from the tree will be handled by a background process after the bud-completion phase, as described in Sect. 3.

A complicated ATG may lead to a large number of incremental bud-generating queries. For an given small relational change $\Delta I$, however, one only needs to evaluate those that refer to a relation affected by $\Delta I$. Techniques [21, 6] for identifying queries irrelevant to $\Delta I$ can further reduce the number of bud-generating queries that need to be evaluated.

Example 5.1: Recall the ATG $\sigma_0$ defined in Fig. 3. Assume that an XML tree $\sigma_0(I)$ of a relational database $I$ is maintained by the middleware. Consider relational changes $\Delta I$: deletion of a CS course from the course relation, along with deletions from $prereq$ and $enroll$ accordingly. Given this the bud-cut-generation phase computes only $E^+$, which consists of deletions of edges (db.course), (course.prereq), etc. These changes are made to the XML tree $\sigma_0(I)$ in this phase.

Note that no buds are generated, i.e., $Bud_A$ is empty for $A$ in the ATG. In other words, although the relational changes have impact to the XML tree at an arbitrary level, they are captured in the bud-cut-generation phase by evaluating a fixed number of incrementalized SQL queries.

On the other hand, if $\Delta I$ is to update the name fields of some $student$ tuples, then it involves both deletions and insertions, i.e., $Bud_{student}$ is nonempty and consists of $student$ buds. The bud-cut-generation phase handles deletions and the next phase, the bud-completion phase, proceeds to generate subtrees of new $student$ buds.

5.3 Bud Completion

The previous bud-cut-generation phase creates sets $Bud_A$ consisting of the $\$A$-attribute values of new $A$-nodes for each element type $A$ in the ATG, and the bud-completion phase is to produce the subtrees of these buds. The process may further generate new buds, but does not incur deletions.

The processing of bud completion is conducted by Algorithm eval, given in Fig. 11. The algorithm takes $Bud_A$'s from the generation phase as input, and processes each non-empty $Bud_A$ based on the production $p = A \rightarrow a$ and its associated rule $p$ (case 1-4), generating children of these $A$-elements. For example, in case 3, it processes a given set $Bud_A$ as follows. First it finds the edges of the $A$-nodes in $Bud_A$ that are to be inserted into $edge_{A,B}$ by evaluating the incremental bud-completing query $Q_{bud_{edge_{A,B}}}(Bud_A)$ in the DBMS. Then it updates the children lists of these $A$-nodes in the subtree pool $T$, taking advantage of the $H$-index and $id_{A}$. Then, for each $B$ child $id_{B}$ of such an $A$-bud $id_{A}$ it invokes the procedure process($B$, $id_{B}$) to inspect whether there already exists an entry for the node $(B, id_{B})$ in the hash index $H$. If so, it simply adds a cross edge $(id_{A}, id_{B})$ instead of recomputing the subtree of $(B, id_{B})$; otherwise, it creates a new $H$-entry for $(B, id_{B})$ and adds the attribute value $\$B$ of $id_{B}$ to $Bud_B$. This yields $Bud_B$, the set of new $B$-buds that will be further expanded at the next level of the tree. Note that to populate $Bud_B$ we need the attribute values $\$B$ of these newly inserted $B$ nodes.

Observe the following properties of Algorithm eval.
Figure 11: Algorithm eval

(1) At each iteration (level), any paths in the XML view that begin at the root and do not encounter a bud are guaranteed to be correct, and thus partial results of the new XML view $\sigma(I + \Delta T)$ can be exposed to the users in parallel with the computation of $\Delta T$.

(2) It minimizes unnecessary computations via the procedure process: it reuses subtrees that has been computed either by Algorithm eval or earlier for the old view $\sigma(I)$. Thus each bud is computed at most once. The reuse of the previous computations is possible since no nodes are removed from the $H$-index or subtree pool at this stage.

(3) The bud-completion phase does not need materialization of the attribute relations, since all the $\delta A$-values that $\Delta Q_{Edge,A,B}$, $\Delta Q_{Edg,Edge}$ needs in $Bud_A$, not in gen_A. Note that only case (3) in Algorithm eval needs to access the DBMS.

(4) The procedure process also ensures that Algorithm eval always terminates, no matter whether the input ATG has exception handling or not. Specifically, if a descendant bud of an $A$-node $v$ with attribute value $\delta A$ is also an $A$ node with the same $\delta A$ value, the node $v$ is not created since there is already an entry for (A, $\delta A$) in the $H$-index; instead, any edge to $v$ is treated as a cross edge and linked to $v$. Thus no infinite computation is incurred.

Example 5.2: Consider again the ATG of relational changes $\Delta I$ that update the name fields of some student tuples. The bud-cut-generation phase generates a set $Bud_{new}$ consisting of these updated student tuples, generates new student buds for these tuples, and redirects affected edges to these new nodes. The bud-completion phase completes the subtree of these new students by creating new name fields and relinking the same fields for the students. This is done in one iteration of the outer loop of Algorithm eval, although the students may appear at arbitrary levels of the XML view.

Now consider changes $\Delta I$ consisting of insertion of a new CS course to the course relation, along with insertions of its prerequisites to prereq and insertions to enroll for the students who have enrolled in the course. The bud-cut-generation phase generates a set $Bud_{new}$ consisting of the newly inserted course tuple, creates a new node $c$ representing the course, and adds an edge from the root to $c$. The bud-completion phase takes $Bud_{new}$ as input and constructs the subtree of $c$. Since all the prerequisites of $c$ are already in the XML view, the subtree of $c$ can be completed in two iterations of the outer loop of Algorithm eval. In the first iteration only $Bud_{new}$ is nonempty, and the algorithm creates new $\text{new, title, prereq, type}$ and $\text{takeby}$ children of $c$. The second iteration completes their subtrees by reusing the existing course and student nodes. Since no new nodes are generated in the second step, the iteration terminates, although the subtree under $c$ may have a depth greater than two due to its prerequisites hierarchy.

5.4 Implementation Issues

Overlapping Phases. While the first and second phase have been presented as completely separate, in practice it may be advantageous to "pipeline" them. In particular, consider the case where $\Delta I$ affects base tables for both $\Delta Q_{Edge,A,B}$ and $\Delta Q_{Edge,Edge}$. It makes sense to compute $\Delta Q_{Edge,A,B}$ first, allowing the $A$ nodes in the tree to be expanded and then $\delta B$ to be incremented before $\Delta Q_{Edge,Edge}$ is executed. This follows the observation in [15] that incremental queries should be evaluated according to their "stratum number", which reflects dependency relationship.

Query Merging. Query merging has been proved useful for systems to publish relational data in XML [11, 4]. The idea is to reduce the number of queries issued to the DBMS by merging multiple queries into a single, larger query via outer-join/outer-union. Query merging can help decrease the communication costs between the middleware and DBMS, while also potentially diminishing query processing time and execution overheads.

The generation of subtrees under buds is similar to the original process of ATG generation [4], and query merging can be used in the bud-completion phase to merge multiple edge-generating queries of the form $\Delta Q_{Edge,A,B}(Bud_A)$ and $\Delta Q_{Edge,Edge}(Bud_B)$ into one. However, the reuse of existing subtrees is an important difference, and direct application of the merging techniques for ATGs [4] may lead to unne-
6. Experimental Results

Our experimental evaluation focuses on the effect of database and update size on the performance of our approach and on the effectiveness of the proposed subtree caching strategy.

We build the source database based on the schema of Example 2.2, but with an additional `taughtby` table giving instances of professors teaching courses. The database size is given in terms of the number of courses, with scaling factors for other tables given in Fig. 12. A smaller pool of professors is used to initially build the tree. In effect, f controls the probability a new prerequisite will exist in the tree already. For the recursive part of the schema, between one and three random prerequisites with lower cons from the pool are generated for each course. The maximum depth of this recursive part is limited to 8 levels of XML nodes.

An update consists of inserting or deleting a course and its associated prerequisites. Fixed-size batches of meaningful updates, \( \Delta I \), are generated by ensuring that inserts are not duplicates and that deleted courses are present. Since deletes in our system generally execute much faster than inserts, \( \Delta I \) is constrained to have 50% inserts unless stated otherwise. The locality of updates is controlled by selecting the updates in \( \Delta I \) from a random pre-generated "universe," \( S \), of courses. When \( S \) is smaller, the updates in \( \Delta I \) are more related. The parameters and their default values are summarized in Fig. 13. Experiments are performed by running sequences of such batches for two workloads \( W_1 \) and \( W_2 \), and using \( \text{bud-cut} \) to propagate the changes to the external tree. The first workload, \( W_1 \), consists of executing such batched updates against our example ATG, \( \sigma_0 \). A second workload \( W_2 \) modifies \( \sigma_0 \) to restrict prerequisites to courses taught by instructors from the computer science department. Thus, \( W_2 \) introduces a join to `taughtby` in the recursive part of the view.

The experiments were run on a system with a 2GHz Pentium 4 processor with 1G byte of RAM; they were conducted with a large (256MB) warm DMs buffer cache. While incremental update may reasonably avoid disk access since the relevant data was just updated, fully cached operation for document publishing is also reasonable, at least for published XML data up to a few tens of megabytes in size given modern database configurations. Similarly, the external XML tree (see Sect. 3) is fully cached in RAM on the same system. Each experiment ran five times and the average is reported here. All numbers obtained for each average are within 6% of the average value. Unless otherwise noted, we used the proposed caching strategy.

We compare the cost of incremental update against the cost of full tree recreation for a variety of database sizes. We scale the database by varying \( |I| \) from 100 to 10K, and setting \( |w| \) to 4% of \( |I| \). For this experiment, locality is not considered (\( S = I \), \( f = 100\% \)). Figure 14(a) shows the resulting time with \( |I| \) on a log-scale for workloads \( W_1 \) and \( W_2 \). Not surprisingly, incremental update far outperforms full recreation for small updates once the overall database size gets appreciably large. Note that incremental update for \( W_2 \) is faster than for \( W_1 \) since the tree is significantly smaller (and thus less update work in the tree is needed).

Fixing \( |I| = 10K \), we vary \( |w| \) and study the behavior of incremental update vs. full tree recreation. The result is presented in Fig. 14 (b). As this figure shows, for \( W_1 \) and \( W_2 \), \( \text{bud-cut} \) scales nicely with the size of the change in the source database, \( |w| \). We note that the \( \text{bud-cut} \) approach is sensitive to the selectivity of the \( \text{ATG} \) queries, and low selectivity leads to substantial activity on the hash index in the subtree pool. Consequently, the incremental update and the full tree recreation for \( W_1 \) and \( W_2 \) cross when 10% and 80% of the database is updated, respectively.

To investigate the impact of the subtree pool, we evaluate the performance impact of turning it off. We fix \( |w| = 100 \) and vary \(|S| \) from 100 to 600 (decreasing the locality of updates as \( |S| \) grows) for \( W_1 \). (In this experiment, the percentage of inserts in \( |w| \) is allowed to vary.) Finally, we consider two values of \( f \), 50% and 80%, to control the probability that a given subtree will already appear in the tree. The results are shown in Fig. 15(a), where each point represents the average of 20 such experiments, with garbage collection fired after each run. As expected, the impact of subtree caching is greatest with smaller values for \(|S| \) and \( f \). With \( f = 50\% \), the impact is substantial across the range. However, the curves for \( f = 80\% \) are rather flat, since most cached subtrees are small due to "natural caching" by the rest of the tree.

While in Fig. 14(b) we compare the size of the change in the source database to the time taken, it is also interesting to consider the size of the change in the output XML tree, and compare this to the time taken for the incremental update. An experiment to investigate this relationship is shown in Fig. 15(b) (for workload \( W_1 \)). To capture this we output the XML tree before and after the incremental update \( w \). We then use X-Diff [32] to produce an edit script converting the old tree to the new tree, and measure its size as an approximation of \( |\Delta T| \). Furthermore, \( |w| \) remains 100, \( f \) is 50% and \(|S| \) is set to 300. This experiment shows that \( \text{bud-cut} \) also behaves reasonably well with respect to output size, and further, that the subtree caching optimization (in combination with the graph-oriented storage model) tends to be more effective exactly when an incremental update to
the database causes a large change in the output tree.

7. Related Work

Incremental computation has proved useful in many areas (see [27] for a survey). In particular, incremental view maintenance algorithms have been extensively studied for relational, datalog (see [13] for a collection of readings) and object-oriented views (see, e.g., [12, 18]). On the one hand, these algorithms are not directly applicable to incremental evaluation of XML publishing. Unlike traditional database views, XML views (ATGs) are defined by associating a collection of SQL queries with a (possibly recursive) DTD. These views are stored outside of DBMS, and their incremental maintenance is to propagate relational group updates to external XML trees. On the other hand, our incremental evaluation algorithms leverages previous work for traditional databases. Our reduction approach relies on the support of incremental maintenance of SQL views by the underlying DBMS, and our top-down mechanism makes use of the counting algorithm of [15] to incrementally evaluate SQL queries.

We now draw the analogy of our top-down ATG evaluation algorithm to incremental evaluation of recursive queries. A number of algorithms have been developed for evaluating recursive datalog queries and Prolog programs, notably [15, 29]. These algorithms typically involve a first phase where conservative deletions are conducted, followed by a second phase where insertions are performed, which may restore information deleted in the first phase. In contrast, our algorithm avoids recomputation between the two phases to an extent by reusing subtrees computed earlier and by deferring removal of disconnected subtrees. For example consider an extreme case where the conservative deletion involves deletion of the root of the tree, and thus the entire tree, whereas in the second phase the tree is reassembled with somewhat differing children of the root. While this may require reconstruction of the entire tree for the previous algorithms [15], our algorithm is capable of making maximum reuse of existing subtrees without doing unnecessary recomputations. Another difference between our algorithm and previous ones is that the conservative deletion phase repeatedly executes deletion operations for an unbounded number of times [15, 29], whereas by taking advantages of the semantics of XML trees, our bud-cut-generation phase does all the deletions by evaluating a fixed number of incrementalized SQL queries.

It is worth remarking that the possibility of encoding ATGs with SQL 99 queries was first suggested by [17], but [17] did not show how the encoding should be defined. It should be mentioned that there is a big gap between XML trees computed by an ATG and their relational encoding, while the former can be exponentially large in the size of the relational database, the latter is bounded by a polynomial in the database size. Thus an SQL 99 encoding of an ATG alone does not suffice to compute XML trees defined via ATG.

There have also been incremental maintenance algorithms for semi-structured views [1, 31, 33]. These algorithms are developed for views defined with a nonrecursive query over graph structures, and cannot be applied to ATG evaluation.

Incremental update of external views has been studied for data warehousing [22, 30, 34] and web site maintenance [19, 20]. Techniques developed in this work are complementary to ours addressing e.g., compensating queries to cope with updates to decoupled sources, maintaining mediated views via constrained rules, scheduling updates for a group of views to maximize data quality, and deciding where to materialize web views with respect to access and update patterns.

There has also been a host of work on the self-maintainability of relational and datalog views (e.g., [21, 65]). We plan to study techniques for detecting updates irrelevant to XML views in order to optimize XML view maintenance.
8. Conclusion

Finally, the strategy commonly used by XML publishing middleware that pushes simple queries to the DBMS and conducts the rest of the work for composing a view at the middleware switches back to query processing for distributed relational databases [5]. Again recursive XML views introduce new challenges not encountered in the relational context.

9. References