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Lossless Selection Views under Conditional Domain Constraints

Ingo Feinerer, Enrico Franconi, and Paolo Guagliardo

Abstract—A set of views defined by selection queries splits a database relation into sub-relations, each containing a subset of the original rows. This decomposition into horizontal fragments is lossless when the initial relation can be reconstructed from the fragments by union. In this paper, we consider horizontal decomposition in a setting where some of the attributes in the database schema are interpreted over a specific domain, on which a set of special predicates and functions is defined. We study losslessness in the presence of integrity constraints on the database schema. We consider the class of conditional domain constraints (CDCs), which restrict the values that the interpreted attributes may take whenever a certain condition holds on the non-interpreted ones, and investigate lossless horizontal decomposition under CDCs in isolation, as well as in combination with functional and unary inclusion dependencies.

Index Terms—Selection, views, losslessness, constraints, CDC, consistency, separability

1 INTRODUCTION

The problem of updating a database through a set of views consists in propagating updates issued on the views to the underlying base relations over which the view relations are defined, so that the changes to the database reflect exactly those to the views. This is a classical problem in database research, known as the view update problem ([1], [2], [3]), which in recent years has received renewed and increasing attention ([4], [5], [6], [7], [8], [9]).

View updates can be consistently propagated in an unambiguous way under the condition that the mapping between database and view relations is lossless, which means that not only do the view relations depend on the database relations, but also the converse is true. However, just knowing that such an “inverse” dependency exists is not yet sufficient to effectively propagate the changes from the views to the database. What is essential to know is how, in some constructive way, the database relations depend on the view relations. This amounts to being able to define each database relation in terms of the views by means of a query, in much the same way the latter are defined from the former [10]. In such a context, database decompositions [11] play an important role, because their losslessness is associated with the existence of an explicit reconstruction operator that, as the name suggests, prescribes how a database relation can be rebuilt from the pieces, called fragments, into which it has been decomposed.

Horizontal decomposition is the process of splitting a given relation into sub-relations on the same attributes and of the same arity, each containing a subset of the rows of the original relation. For example, consider the relation R shown in Fig. 1, recording data about the employees of a company: their name (EMP), the department (DEP) and the position (POS) in which they are employed, and their income (e.g., euros per month) consisting of a fixed salary (SAL) plus a variable bonus (BON). In Fig. 1, the relation R is decomposed into three fragments: V1 selects the rows of R with employees working as managers in departments other than ICT, V2 selects the rows of R with employees who get strictly less than 4,000 as bonus, and V3 selects the rows of R with employees who do not work as managers. This kind of decomposition is lossless when the original relation can be reconstructed from the fragments by union; in other words, the reconstruction operator for horizontal decomposition is the union. In the example of Fig. 1, the set of views {V1, V2, V3} constitutes a lossless horizontal decomposition of R, as the union of V1, V2 and V3 contains all (and only) the rows of R. Each proper subset of {V1, V2, V3} constitutes a lossy decomposition of R, because each view selects at least one row that is not selected by any of the others; e.g., the union of V1 and V2 does not contain the third row of R, (Linda, Finance, Consultant, 5,000, 4,000), which is selected only by V3.

Observe that the horizontal decomposition specified by the definitions of views V1, V2 and V3 in Fig. 1 is lossless for the given relation R, but this is not the case for every relation (over the same attributes). For instance, the tuple (Sam, ICT, Manager, 6,000, 5,000) is not selected by any of these views; indeed, every relation containing a row for an employee who works as a manager in the ICT department and receives a bonus greater than 4,000 would not be losslessly decomposed by V1, V2 and V3. In the presence of integrity constraints, however, things may be different, because some tuples, such as the one above, might not be allowed in the input relation.

The study of horizontal decomposition ([12], [13], [14], [15]) has mostly focused on settings where data values can only be compared for equality. However, most real-world
applications make use of data values coming from domains with a richer structure (e.g., ordering) on which a variety of other restrictions besides equality can be expressed (e.g., that of being within a range or above a threshold). Examples are the attributes \( \text{SAL} \) and \( \text{BON} \) of the relation in Fig. 1, the dimensions, weights and prices in the database of a shipping company, or the various amounts (credits, debits, interest and exchange rates, etc.) recorded in a banking application. It is therefore of practical interest to consider a scenario where some of the attributes in the database schema are interpreted over a specific domain, such as the reals or the integers, on which a set of predicates (e.g., smaller/greater than) and functions (e.g., addition and subtraction) are defined, according to a first-order language \( \mathcal{C} \).

In the present paper, we consider horizontal decomposition in a setting with interpreted attributes, in which fragments are defined by selection queries consisting of a condition on the non-interpreted attributes, expressed by a Boolean combination of equalities, and a condition on the interpreted attributes, expressed by a formula in \( \mathcal{C} \). In particular, we study the losslessness (w.r.t. every input relation) of horizontal decompositions specified in this way, in the presence of integrity constraints on the database schema. We work under the pure universal relation assumption (URA) [11], that is, we restrict ourselves to a database schema consisting of only one relation symbol, as customary in the study of database decomposition.

### 1.1 Contribution and Outline

In Section 2, we introduce a class of integrity constraints called conditional domain constraints (CDCs). By means of a formula in \( \mathcal{C} \), a CDC restricts the values that the interpreted attributes can take whenever a certain condition is satisfied by the non-interpreted ones. Depending on the expressive power of \( \mathcal{C} \), CDCs can capture constraints that naturally arise in practice; for example, in the scenario of Fig. 1, it may be required that employees in the ICT department have a total income (i.e., salary plus bonus) of at most 5,000, that employees working as managers get a bonus of at least 2,000, and that employees never receive a bonus greater than their salary. These constraints can be expressed as:

\[
\text{DEP} = \text{"ICT"} \Rightarrow \text{SAL} + \text{BON} \leq 5,000 \quad (1a)
\]

- \( \text{POS} = \text{"Manager"} \Rightarrow \text{BON} \geq 2,000 \quad (1b) \)
- \( \text{SAL} - \text{BON} \geq 0. \quad (1c) \)

As we shall see, the views of Fig. 1 losslessly decompose every relation satisfying the above CDCs.

In our investigation, we do not commit to any specific language \( \mathcal{C} \) and we simply assume that \( \mathcal{C} \) is closed under negation.

In Section 3, we characterise consistent sets of CDCs in terms of satisfiability in \( \mathcal{C} \). Whenever the satisfiability of sets of formulae in \( \mathcal{C} \) is decidable, our characterisation directly gives a decision procedure for checking whether a set of CDCs is consistent. This is the case, e.g., for the so-called Unit Two Variable Per Inequality fragment of linear arithmetic over the integers, whose formulae (referred to as UTVPIs) consist of at most two variables and variables have unit coefficients, as well as for Boolean combinations of such formulae. We prove that deciding consistency is \( \mathsf{NP} \)-complete for both of these languages.

In Section 4, we characterise lossless horizontal decomposition under CDCs in terms of unsatisfiability in \( \mathcal{C} \). Whenever the satisfiability of sets of formulae in \( \mathcal{C} \) is decidable, this characterisation gives a decision procedure for checking whether a horizontal decomposition is lossless under CDCs. We show that this problem is \( \mathsf{co-NP} \)-complete when \( \mathcal{C} \) is the language of either UTVPIs or Boolean combinations of UTVPIs.

In Section 5, we study lossless horizontal decomposition under CDCs in combination with traditional integrity constraints. We show that functional dependencies (FDs) do not interact with CDCs and can thus be allowed without any restriction, whereas this is not the case for unary inclusion dependencies (UINDs). We provide a domain propagation rule to derive a set of CDCs that fully captures the interaction between a given set of UINDs and opportunely restricted CDCs w.r.t. lossless horizontal decomposition, which makes possible to employ the general technique for deciding losslessness also in the presence of UINDs. In addition, we consider restricted combinations of CDCs with both FDs and UINDs.

We conclude in Section 6 with a discussion of the results, relevant related work and future research directions.

## 2 Preliminaries

We start by introducing the necessary notation and notions that will be used throughout the paper. We assume some familiarity with formal logic and its application to database theory.

**Basics.** An \( n \)-tuple is an ordered list of \( n \) elements, where \( n \) is a positive integer. We denote tuples by overlined lowercase letters (e.g., \( \overline{t} \)) and we write them as comma-separated sequences enclosed in parentheses; the \( k \)th element of a tuple \( \overline{t} \) is denoted by \( t[k] \).

A schema is a finite set \( S \) of relation symbols, also called a relational signature. Each relation symbol \( S \) has a positive arity \( |S| \) indicating the total number of positions in \( S \), which are partitioned into interpreted and non-interpreted ones.
Relation symbols of arity $n$ are called $n$-ary; we indicate that $|S| = n$ by writing $S/n$.

Let $\text{dom}$ be a possibly infinite set of arbitrary values, and let $\text{idom}$ be a set of values from a specific domain (e.g., the integers $\mathbb{Z}$) on which a set of predicates (e.g., $\leq$) and functions (e.g., $\oplus$) are defined, according to a first-order language $\mathcal{L}$ closed under negation. An instance over a schema $S$ associates each $S \in S$ with a relation $S^I$ of appropriate arity on $\text{dom} \cup \text{idom}$, called the extension of $S$ under $I$, such that the values for the interpreted and non-interpreted positions of $S$ are taken from $\text{idom}$ and $\text{dom}$, respectively. The set of elements of $\text{dom} \cup \text{idom}$ occurring in an instance $I$ is called the active domain of $I$, denoted by $\text{adom}(I)$. An instance is finite if its active domain is, and all instances in this paper are assumed to be finite. A fact is given by the association, denoted by $R(I)$, between a relation symbol $R$ and a tuple $\overline{t}$ of values of appropriate arity; an instance can be represented as a set of facts.

Constraints. A language over a relational signature $S$ is a set of first-order logic (FOL) formulae over $S$ with constants $\text{dom} \cup \text{idom}$ under the standard name assumption (i.e., the interpretation of each constant is the constant’s name itself). A formula in some language $\mathcal{L}$ is called an $\mathcal{L}$-formula. The sets of constants and relation symbols that occur in a formula $\varphi$ are denoted by $\text{const}(\varphi)$ and $\text{sig}(\varphi)$, respectively; we extend $\text{const}(\cdot)$ and $\text{sig}(\cdot)$ to sets of formulae in the natural way.

A constraint is a closed formula (that is, without free variables) in some language. For a set $\Gamma$ of constraints, we say that an instance $I$ over $\text{sig}(\Gamma)$ is a model of (or satisfies) $\Gamma$, and write $I \models \Gamma$, to indicate that the relational structure $\langle \text{dom}(I) \cup \text{const}(\Gamma), I \rangle$ makes every formula in $\Gamma$ true under the standard FOL semantics. We write $I \models \varphi$ as short for $I \models \{ \varphi \}$, and say that $I$ satisfies $\varphi$. A set of constraints $\Gamma$ entails (or logically implies) a constraint $\varphi$, written $\Gamma \models \varphi$, if every finite model of $\Gamma$ also satisfies $\varphi$. All sets of constraints in this paper are finite.

Propositional theories. A propositional variable is a variable whose value can be either $\top$ (true) or $\bot$ (false). A propositional formula is a Boolean combination of propositional variables, including the two special propositional variables $\top$ and $\bot$, whose values are always $\top$ and $\bot$, respectively. A propositional theory is a set of propositional formulae. We denote the set of propositional variables occurring in a propositional formula $P$ by $\text{var}(P)$ and we extend $\text{var}(\cdot)$ to propositional theories in the natural way. A valuation of a set of propositional variables (also called a truth-value assignment) assigns a truth-value (i.e., either $\top$ or $\bot$) to each propositional variable in the set. The truth-value $\alpha(P)$ of a propositional formula $P$ under a valuation $\alpha$ of its propositional variables is determined by the standard semantics of the Boolean connectives. We say that $\alpha$ satisfies (or makes true) $P$, and write $\alpha \models P$, if $\alpha(P) = \top$. Given a propositional theory $\Pi$, a valuation of $\text{var}(\Pi)$ satisfies $\Pi$, written $\alpha \models \Pi$, if $\alpha$ satisfies every propositional formula in $\Pi$.

2.1 Horizontal Decomposition

We consider a source schema $R$, consisting of a single relation symbol $R$, and a decomposed schema $V$, disjoint with $R$, of view symbols with the same arity as $R$. We formally define horizontal decomposition as follows.

Definition 1. Let $R = \{ R \}$ and $V = \{ V_1, \ldots, V_n \}$. Let $\Delta$ be a set of constraints over $R$ and let $\Sigma$ be a set of exact view definitions, one for each $V_i \in V$, of the form $\forall x. V_i(x) \leftrightarrow \varphi(x)$, where $\varphi$ is a safe-range formula over $R$. Then, $\Sigma$ is a horizontal decomposition of $R$ into $V$ under $\Delta$ if $\Delta \cup \Sigma \models \forall x. V_i(x) \rightarrow \varphi(x)$ for every $V_i \in V$. We say that $\Sigma$ is lossless if $\Delta \cup \Sigma \models \forall x. R(x) \leftrightarrow V_1(x) \land \cdots \land V_n(x)$.

For the sake of simplicity, w.l.o.g. we assume that the first $|R|$ positions of $R$ and of every $V \in V$ are non-interpreted, while the remaining $|R| - |V|$ positions are interpreted. Under this assumption, instances over $R \cup V$ associate each relation symbol with a subset of the Cartesian product $\text{dom}^n \times \text{idom}^{n-k}$, where $n = |R|$ and $k = |\Sigma|$. Unless otherwise specified, when we speak of a tuple $\overline{t}$ we implicitly assume that $\overline{t}$ is of arity $n$ and that the first $k$ values of $\overline{t}$ are from $\text{dom}$ while the rest are from $\text{idom}$. W.l.o.g. we also assume that a variable associated with the $i$th position of $R$ is named $x_i$, if $i \leq k$, and $y_{i-k}$ otherwise. By default, $\top$ and $\bot$ denote the tuples $(x_1, \ldots, x_k)$ and $(y_1, \ldots, y_{n-k})$, respectively.

Since every $\mathcal{L}$-formula is over variables associated with interpreted positions, we write $\phi[\overline{t}]$ to indicate that $\phi$ is a $\mathcal{L}$-formula whose free variables are among the variables in $\overline{t}$. For a tuple $\overline{t}_y$ of $n - k$ values from $\text{idom}$, we denote by $\phi[\overline{t}_y]$ the result of replacing every occurrence of $\varphi$ in the free variable $y_i$ with the value $\overline{t}_y[i]$. We say that $\overline{t}_y$ is a solution to $\phi$ if $\phi[\overline{t}_y]$ is true under the semantics of $\mathcal{L}$. In such a case, we also say that the assignment $\beta$ associating each $y_i$ with $\overline{t}_y[i]$ satisfies $\phi$, and we write $\beta \models \phi$.

Source constraints. The class of integrity constraints we consider on the source schema $R$ is that of conditional domain constraints, which restrict the admissible values at interpreted positions by means of formulae in $\mathcal{L}$, when a certain condition holds on the non-interpreted ones. Formally, a CDC is a formula of the form

$$\forall x, \overline{y}. (R(x, \overline{y}) \land \lambda(x)) \rightarrow \delta(\overline{y}),$$

where $\lambda(x)$ is a Boolean combination of equalities $x = a$, with $a$ from $\mathcal{T}$ and $a$ from $\text{idom}$, and $\delta(\overline{y}) \in \mathcal{C}$. We use $x \neq a$ as short for $\neg(x = a)$ and, for ease of notation, we write (2) simply as $\lambda(x) \rightarrow \delta(\overline{y})$. Here, we make use of a more general variant of the CDCs introduced in [16], where the condition $\lambda(x)$ was limited to a conjunction of possibly negated equalities. In general, (2) is more expressive than a CDC of the form used in [16], as in the latter negation is allowed only atomically, that is, in front of equalities, and so disjunction cannot be expressed. However, there is no difference in expressivity between the two variants when considering sets of CDCs, because the antecedent of (2) can always be rewritten in disjunctive normal form (DNF) and the CDC split into a set of CDCs having the same consequent and each disjunct as antecedent.

View definitions. The view symbols in $V$ are defined by selection queries with conditions on both interpreted and non-interpreted positions. Formally, each $V \in V$ is defined by a formula of the form

1. For details on the syntactic notion of range restriction, corresponding to the semantic notion of domain independence, refer to [11].
2. Sometimes, by abuse of terminology, we say that an assignment $\beta$ is a solution to a $\mathcal{L}$-formula $\phi$, with the obvious meaning.
\[ \forall \bar{x}, \bar{y}. \; V(\bar{x}, \bar{y}) \leftrightarrow (R(\bar{x}, \bar{y}) \land \lambda(\bar{x}) \land \sigma(\bar{y})), \]  
\( (3) \)

where \( \lambda(\bar{x}) \) is as in (2) and \( \sigma(\bar{y}) \in \mathcal{C} \). In the following, we write (3) simply as \( V: \lambda(\bar{x}) \land \sigma(\bar{y}) \). View definitions of this form clearly generalise those in [16], where \( \lambda(\bar{x}) \) is limited to a conjunction of possibly negated equalities and disjunction cannot be expressed. While this has an impact on end-users, who can define more expressive views, there is no difference between the two formalisms w.r.t. losslessness, in that any view symbol \( V \) defined by (3) can be split into a set of views, defined by formulæ of the form used in [16], that together select exactly the same tuples as \( V \); given \( \lambda(\bar{x}) \) in DNF, each of these view definitions has the same selection condition as \( V \) on the interpreted attributes and a disjunct of \( \lambda(\bar{x}) \) as selection condition on the non-interpreted ones.

The technique we will present in Section 4 for checking whether a set of selections of the form (3) is lossless can also be applied when (some of) the selections have the form \( V: \lambda(\bar{x}) \lor \sigma(\bar{y}) \) by considering, in place of each such selection, the two selections \( V': \lambda(\bar{x}) \) and \( V'': \sigma(\bar{y}) \).

Running example. To clarify the notation and illustrate the concepts introduced so far, we now give an example that will be used also in the rest of the paper. It is based on the source schema of Fig. 1, the CDCs (1a)-(1c) informally described in Section 1 and the views previously specified in Fig. 1 by means of SQL statements.

Example 1. Let \( R \) be a relation symbol of arity 5, whose positions are associated with attributes EMP, DEP, POS, SAL, BON, in this order, with the last two interpreted over the integers. Differently from the example of Fig. 1, for simplicity we assume that salaries and bonuses are given in thousands of euros/month. Let \( a = "ICT" \) and \( b = "Manager" \), and consider the following set \( D \) of CDCs: \( \{x_2 = a \rightarrow y_1 + y_2 \leq 5; x_3 = b \rightarrow y_2 \geq 2; T \rightarrow y_1 = y_2 \geq 0\} \). Let \( V = \{V_1, V_2, V_3\} \), and let \( \Sigma \) be the horizontal decomposition given by \( V_1: x_2 \neq a \land x_3 = b, V_2: y_2 < 4, \) and \( V_3: x_3 \neq b \).

Specific languages. The techniques we will present for deciding whether a set of CDCs is consistent (Section 3) and whether a horizontal decomposition is lossless under CDCs (Section 4) give actual algorithms when satisfiability in \( \mathcal{C} \) is decidable; in the case of losslessness, \( \mathcal{C} \) is additionally required to be closed under negation. Thus, even though our investigation is in general independent of the choice of \( \mathcal{C} \), from a practical point of view it makes sense to consider concrete languages that enjoy both of the above properties. Two prominent such languages are Unit Two-Variable Per Inequality formulæ (UTVPIs) and Boolean combinations thereof. UTVPIs, a.k.a. Generalised 2 SAT (G2SAT) formulæ [17], are a fragment of linear arithmetic over the integers. Formally, a UTVI formula has the form \( ax + by \leq d \), where \( a \) and \( y \) are integer variables, \( a, b \in \{ -1, 0, 1 \} \) and \( d \in \mathbb{Z} \).

UTVPIs can express comparisons between two variables and between a variable and an integer, as well as compare the sum or difference of two variables with an integer. As integers allow to represent also real numbers with fixed precision, UTVPIs may be sufficient for most applications. The CDCs and the view definitions of Example 1 can be expressed when \( \mathcal{C} \) is the language of UTVPIs.

Whether a set of UTVPIs is satisfiable can be checked in polynomial time ([18], [19], [20]). We refer to a Boolean combination of UTVPIs as BUTVP; deciding the satisfiability of a set of BUTVPs is NP-complete [21].

3 Consistent Sets of CDCs

Before turning our attention to horizontal decomposition, we first deal with the relevant problem of determining whether a set of CDCs is consistent, that is, whether it has a non-empty model. It is important to make sure that the integrity constraints over the source schema are consistent, as every horizontal decomposition is meaninglessly lossless when there are in fact no legal relations to decompose.

In this section, we will characterise the consistency of a set of CDCs in terms of satisfiability in \( \mathcal{C} \), where \( \mathcal{C} \) is not required to be closed under negation. The consistency problem for CDCs is the decision problem that takes as input a set \( \Delta \) of CDCs and answers the question: "Is \( \Delta \) consistent?" We will show that when \( \mathcal{C} \) is the language of either UTVPIs or BUTVPs this problem is NP-complete. The technique employed here provides the basis for the approach we follow in Section 4 in the study of lossless horizontal decomposition.

Observe that, given their form, CDCs affect only one tuple at a time, and so whether an instance satisfies a set of CDCs depends on each tuple of the instance in isolation from the others. Indeed, a set of CDCs is consistent precisely if it is satisfiable on an instance consisting of only one tuple, therefore we can restrict our attention to single tuples. Moreover, we are not really interested in the actual values of a tuple at non-interpreted positions; what we need to know is simply whether such values satisfy the conditions in the antecedent of each CDC or not. To this end, with each equality between a variable \( x_i \) and a constant \( a \) we associate a propositional variable \( p_{ij}^x \), whose truth-value indicates whether the value in the \( i \)-th position is \( a \). To each valuation of such propositional variables corresponds the (possibly infinite) set of tuples satisfying the equalities associated with the names of the propositional variables. For example, a valuation assigning true to \( p_1^x \) and false to \( p_2^x \) identifies all the tuples in which the value of the first element is \( a \) and the value of the second is different from \( b \). A bit more care is needed with valuations of propositional variables that refer to the same position (i.e., have the same subscript) but to different constants (i.e., have different superscripts). For example, \( p_1^x \) and \( p_2^x \) (with \( a \neq b \)) should never be both evaluated to true.

As we shall see, checking whether a set \( \Delta \) of CDCs is consistent amounts to first building a propositional theory by replacing the equalities with the corresponding propositional variables, and then looking for a valuation \( \alpha \) such that:

- any two propositional variables referring to the same position but to different constants are not evaluated both to true; and
- the set of \( \mathcal{C} \)-formulæ that "apply" under \( \alpha \) is satisfiable.

Definition 2. Let \( \Delta = \{ \phi_1, \ldots, \phi_r \} \) be a set of CDCs over \( R \). For each \( \phi_i \in \Delta \), recalling it has the form (2), we construct

3. Since CDCs are universally-quantified closed implicational formulæ, any set thereof is always trivially satisfied by the empty instance.
prop(ϕ) = P → v_i, \quad (4)

where P is a propositional formula (possibly ⊤) obtained from the condition λ(π) in the antecedent of ϕ by replacing each equality x_i = a between a variable x_i and a ∈ idom with the propositional variable p^i_a, and v_i is a fresh propositional variable associated with the ℂ-formula δ(π), denoted by idf(v_i), in the consequent of ϕ. We denote \{prop(ϕ) | ϕ ∈ ∆\} by Π_∆ and we call it the propositional theory associated with ∆.

We consider the set var(Π_∆) of propositional variables occurring in Π_∆ partitioned into var_r(Π_∆) = {var(P) | (P → v_i) ∈ Π_∆} and var_a(Π_∆) = var(Π_∆) \ var_r(Π_∆).

For a pair of distinct propositional variables p^i_a and p^j_b associated with the same position i but distinct dom constants a and b, we consider the propositional formula p^i_a ∧ p^j_b → ⊥, called the axiom of unique value for p^i_a and p^j_b, intuitively stating that two distinct constants are not allowed in the same position. The axioms of unique value for a set of propositional variables consist of the axiom of unique value for each pair of distinct propositional variables p^i_a and p^j_b in the set. A tuple T is consistent with a valuation α if, for every propositional variable p^i_a, it holds that T[i] = a precisely if α(p^i_a) = T. In general, given a valuation α of a set of propositional variables, by construction there exists a tuple consistent with α if and only if α satisfies the corresponding axioms of unique value for that set.

Definition 3. Let ∆ be a set of CDCs over R. The auxiliary theory Π_aux for Π_∆ consists of the axioms of unique value for var_r(Π_∆).

Example 2. The propositional theory associated with Π_∆ is Π_∆ = \{ p^2_a → v_1, p^3_a → v_2, T → v_3 \} for which var_r(Π_∆) = \{ p^2_a, p^3_a \} and var_a(Π_∆) = \{ v_1, v_2, v_3 \}. The auxiliary theory for Π_∆ is Π_aux = ∅. The association between the propositional variables in var_r(Π_∆) and the set of UTVPIs from the CDCs in Π_∆ is idf = \{ v_1 → y_1 + y_2 ≤ 5, v_2 → y_2 ≥ 2, v_3 → y_1 − y_2 ≥ 0 \}.

Given a set ∆ of CDCs and a valuation α of var_r(Π_∆), we say that a CDC ϕ ∈ ∆ is applicable under α if α makes the l.h.s. of prop(ϕ) true. We can use α to “filter” Π_∆ and construct a set consisting of the consequent of each CDC in Π_∆ that is applicable under α. This set contains the ℂ-formulae that must be necessarily satisfied by the values at interpreted positions of every tuple consistent with α, that is, whose values at non-interpreted positions satisfy the antecedents of the CDCs applicable under α.

Definition 4. Let ∆ consist of CDCs over R, and let α be a valuation of var_r(Π_∆). The α-filtering of Π_∆ is the set

Π^u_∆ = \{ idf(v) | (P → v) ∈ Π_∆, α(P) = T \}, \quad (5)

consisting of ℂ-formulae associated with propositional variables that occur in some propositional formula of Π_∆ whose l.h.s. holds true under α.

The main result of this section characterises the consistency of a set of CDCs in terms of satisfiability in ℂ. We remark again that the result holds in general for any language ℂ, not necessarily closed under negation. This requirement will become essential only in the upcoming Sections 4 and 5.

Theorem 1. Let ∆ be a set of CDCs over R, and let Π_aux be the auxiliary theory for Π_∆. Then, ∆ is consistent if and only if there exists a valuation α of var_r(Π_∆) satisfying Π_aux and such that Π^u_∆ is satisfiable.

Whenever the satisfiability of sets of ℂ-formulae is decidable, Theorem 1 gives an algorithm to check whether a set of CDCs is consistent, as we illustrate below in our running example, where ℂ is the language of UTVPIs.

Example 3. With respect to Π_∆ of Example 2, consider the valuation α = \{ p^2_a → T, p^3_b → F \}, for which we have Π^u_∆ = \{ y_1 + y_2 ≤ 5, y_1 − y_2 ≥ 0 \}. Obviously, α satisfies the (empty) auxiliary theory Π_aux for Π_∆. In addition, Π^u_∆ is satisfiable as, e.g., \{ y_1 → 3, y_2 → 2 \} is a solution to every UTVPI in it.

We will now give the proof of Theorem 1, for which we first need to prove a technical lemma. Let n = |R| and k = ||R|| with each tuple T is associated the assignment β: \{ y_1, \ldots, y_{n−k} \} → idom, which we refer to as the assignment induced by the interpreted positions of T, such that β(y_{i−k}) = T[i] for every i ∈ {k + 1, \ldots, n}. Intuitively, the following lemma shows that any tuple that is consistent with a valuation α satisfies a set of CDCs precisely if the assignment induced by its interpreted positions satisfies the α-filtering.

Lemma 1. Let ∆ be a set of CDCs over R, and let α be a valuation of var_r(Π_∆). Let T be consistent with α, and let β be the assignment induced by the interpreted positions of T. Then, \{ R(T) \} ∈ ∆ if and only if β satisfies Π^u_∆.

Proof. Let n = |R| and k = ||R||.

Claim 1. Let ϕ ∈ ∆ and prop(ϕ) = P → v. Then, α(P) = T iff λ(π) is true under \{x_1 → T[1], \ldots, x_k → T[k]\}.

Proof. Since T is consistent with α, for i ∈ \{1, \ldots, k\} we have that T[i] = a if and only if α(p^i_a) = T. Claim 1 follows.

Claim 2. For each prop(ϕ) = P → v with ϕ ∈ Π_∆, it is the case that I ⊧ α if and only if α(P) = T and β(ϕ) = idf(v).

Proof. As ϕ is a CDC, I ⊧ ϕ if and only if the antecedent λ(π) of ϕ holds true under \{x_1 → T[1], \ldots, x_k → T[k]\} and the consequent δ(π) of ϕ is not true under \{ y_1 → T[k+1], \ldots, y_{n−k} → T[n]\}. In turn, this is the case if and only if both α(P) = T (by Claim 1) and β does not satisfy idf(v) = δ(π) (by construction).

We prove Lemma 1 by showing that I ⊧ ∆ if and only if β does not satisfy Π^u_∆.

“if” Assume β ⊧ Π^u_∆, that is, there is some ℂ-formula ψ ∈ Π^u_∆ not satisfied by β. By construction of Π^u_∆, ψ is the consequent of a CDC ϕ ∈ ∆ such that prop(ϕ) = P → v, with ψ = idf(v) and α(P) = T. Thus, as β does not satisfy ψ, by Claim 2 I ⊧ φ, and therefore I ⊧ ∆.

“Only if” Assume I ⊧ ∆. Then, there exists some φ ∈ ∆ which is not satisfied by I. Since prop(ϕ) = P → v, Claim by 2 α(P) = T and β ⊧ idf(v). Hence, idf(v) ∈ Π^u_∆. Therefore, β ⊧ Π^u_∆.
Proof of Theorem 1. Let \( n = |R| \) and \( k = ||R||. \)

“If” Let \( \alpha \) and \( \beta \) be such that \( \alpha \models \Pi_{\text{aux}} \) and \( \beta \models \Pi_{\text{aux}}^\Delta. \) Then, as \( \alpha \models \Pi_{\text{aux}} \), it is never the case that two distinct propositional variables in \( \forall_{\text{prop}}(\Pi_{\Delta}) \) associated with the same position are both true under \( \alpha. \) Thus, there exists a tuple \( \tau \) consistent with \( \alpha \) and such that \( \beta \) is the assignment induced by its interpreted positions. Therefore, as \( \beta \models \Pi_{\text{aux}}^\Delta, \) the instance \( \{R(\tau)\} \) is a model of \( \Delta \) by Lemma 1.

“Only if” Assume that \( \Delta \) is consistent, that is, it has a non-empty model. In particular, as every formula in \( \Delta \) is in one tuple, there is a tuple \( \tau \) such that the instance \( I = \{R(\tau)\} \) is a model of \( \Delta. \) Take \( \alpha \) as follows: for every propositional variable \( p \in \forall_{\text{prop}}(\Pi_{\Delta}) \), \( \alpha(p) = \top \) if \( p = p_i^\Delta \) and \( I[i] = a_i \), otherwise \( \alpha(p) = F. \) By construction, \( \alpha \models \Pi_{\text{aux}} \) and \( \tau \) is consistent with \( \alpha \). Therefore, as \( I \models \Delta, \) the assignment \( \beta \) induced by the interpreted positions of \( \tau \) satisfies \( \Pi_{\text{aux}}^\Delta \) by Lemma 1.

The satisfiability problem for \( \mathcal{C} \) takes as input a set \( \Gamma \) of \( \mathcal{C} \)-formulae and answers the question: “Is \( \Gamma \) satisfiable?”

Lemma 2. The satisfiability problem for \( \mathcal{C} \) linearly reduces to the consistency problem for CDCs.

Proof. Let \( \Gamma = \{\phi_1, \ldots, \phi_n\} \) be a set of \( \mathcal{C} \)-formulae. Then, take \( \Delta = \{\top \rightarrow \phi_i \mid \phi_i \in \Gamma\}, \) let \( \Pi_{\Delta} = \{\top \rightarrow v_i \mid i = 1, \ldots, n\} \) and \( \Pi_{\text{aux}} = \{v_i \rightarrow \phi_i \mid i = 1, \ldots, n\}. \) The auxiliary theory for \( \Pi_{\Delta} \) is \( \Pi_{\text{aux}} = \emptyset. \) As \( \forall_{\text{prop}}(\Pi_{\Delta}) = \emptyset, \) the only valuation of \( \forall_{\text{prop}}(\Pi_{\Delta}) \) is \( \alpha = \emptyset, \) which satisfies \( \Pi_{\text{aux}} \) and for which \( \Pi_{\text{aux}}^\Delta = \{id(\{i\}): v_i \in \forall_{\text{prop}}(\Pi_{\Delta})\} = \Gamma. \) Thus, by Theorem 1, the set \( \Delta \) of CDCs is consistent iff \( \Gamma \) is satisfiable. The reduction is linear in the size of \( \Gamma. \)

With regard to the consistency problem for CDCs whose consequents are either UTVPIs or BUTVPIs, we have the following complexity results.

Theorem 2. When \( \mathcal{C} \) is the language of either BUTVPIs or UTVPIs, the consistency problem for CDCs is NP-complete.

Proof sketch. Constructing the propositional theories \( \Pi_{\Delta} \) and \( \Pi_{\text{aux}} \) requires linear time, checking that a valuation \( \alpha \) of \( \forall_{\text{prop}}(\Pi_{\Delta}) \) satisfies \( \Pi_{\text{aux}} \) takes polynomial time, and checking that an assignment from the variables in \( \forall_{\text{prop}}(\Pi_{\Delta}) \) to integers satisfies \( \Pi_{\text{aux}}^\Delta \) (whose construction takes linear time) can be done in polynomial time, whether \( \Pi_{\text{aux}}^\Delta \) consists of either UTVPIs or BUTVPIs. Hence, in light of Theorem 1, we can verify a given solution to the consistency problem, when \( \mathcal{C} \) is the language of either UTVPIs or BUTVPIs, in polynomial time.

The NP-hardness of the consistency problem when \( \mathcal{C} \) is the language of BUTVPIs follows by Lemma 2 from the fact that the satisfiability problem for BUTVPIs is NP-hard.

The NP-hardness of the consistency problem when \( \mathcal{C} \) is the language of UTVPIs can be shown by a reduction from SAT.

\[ \text{Proof of Theorem 1.} \]

\[ \text{Proof of Theorem 2.} \]

\section{Lossless Selections under CDCs}

The technique described in the previous section can be opportunistically extended and applied for checking whether a set of selection views of the form (3) is lossless under CDCs, that is, whether every source relation satisfying the given CDCs can be reconstructed by union from the fragments into which it is decomposed by the given view definitions.

In this section, we will characterise lossless horizontal decomposition in terms of unsatisfiability in \( \mathcal{C} \), where \( \mathcal{C} \) is closed under negation. The losslessness problem in \( \mathcal{C} \) is the decision problem that takes as input a horizontal decomposition \( \Sigma \) specified by selections of the form (3) and a set \( \Delta \) of CDCs and answers the question: “Is \( \Sigma \) lossless under \( \Delta? \)” We will show that this problem is co-NP-complete when \( \mathcal{C} \) is the language of either UTVPIs or BUTVPIs. For these languages, our characterisation provides an exponential-time algorithm for deciding the losslessness of \( \Sigma \) under \( \Delta \), by means of a number of unsatisfiability checks in \( \mathcal{C} \) which is exponentially bounded by the size of \( \Delta. \)

By definition, a horizontal decomposition \( \Sigma \) of \( R \) into \( V_1, \ldots, V_n \) is lossless under a set \( \Delta \) of CDCs over \( R \) if \( R^\Delta = V_1^\Delta \cup \cdots \cup V_n^\Delta \) for every model \( I \) of \( \Delta \cup \Sigma. \) As the extension of each view symbol is always included in the extension of \( R, \) the problem is equivalent to checking that there is no model \( I \) of \( \Delta \cup \Sigma \) where a tuple \( \tau \in R^\Delta \) does not belong to any \( V_i^\Delta. \) In turn, this means that for each definition in \( \Sigma, \) which has the form (3), the values in \( \tau \) at non-interpreted positions do not satisfy \( \lambda, \) or the values in \( \tau \) at interpreted positions do not satisfy the \( \mathcal{C} \)-formula \( \sigma. \)

The formulae in \( \Sigma \) apply to one tuple at a time and, as already observed in Section 3, so do CDCs; therefore we can again focus on single tuples. With each equality we associate, as before, a propositional variable whose truth-value determines whether the equality is satisfied. Given a valuation \( \alpha \), we consider the set consisting of \( \mathcal{C} \)-formulae in the r.h.s. of all the CDCs that are applicable under \( \alpha \) and the negation of the selection condition \( \delta(\overline{y}) \) of each view definition in \( \Sigma \) whose selection condition \( \lambda(\overline{x}) \) is satisfied by \( \alpha. \) Then, checking losslessness is equivalent to checking that there exists no valuation \( \alpha \) for which the above set of \( \mathcal{C} \)-formulae is satisfiable. Indeed, from such a valuation and the corresponding assignment of values from \( \text{idom} \) satisfying the relevant \( \mathcal{C} \)-formulae, we can obtain a tuple that provides a counterexample to losslessness.

Similarly to what we did in Section 3 for sets of CDCs, we build a propositional theory associated with a given horizontal decomposition.

Definition 5. Let \( \Sigma = \{\phi_1, \ldots, \phi_n\} \) be a horizontal decomposition. For each \( \phi_i \in \Sigma, \) which has the form (3), we build

\[ \text{prop}(\phi_i) = P \rightarrow v_i^\prime, \]

in which \( v_i^\prime \) is either a fresh propositional variable associated (by means of idf) with the \( \mathcal{C} \)-formula \( \delta(\overline{y}) \), if any, occurring in \( \phi_i \), or \( \bot \) otherwise.\(^5\) We denote \{prop(\phi) \mid \phi \in \Sigma \} \) by \( \Pi_\Sigma \) and we call it the propositional theory associated with \( \Sigma. \)

We consider the set \( \forall_{\text{prop}}(\Pi_\Sigma) \) of propositional variables occurring in \( \Pi_\Sigma \) partitioned into \( \forall_{\text{prop}}(\Pi_\Sigma) = \{\forall_{\text{prop}}(P) \mid (P \rightarrow v_i) \in \Pi_\Sigma \text{ and } \forall_{\text{prop}}(\Pi_\Sigma) = \forall_{\text{prop}}(\Pi_\Sigma) \setminus \forall_{\text{prop}}(\Pi_\Sigma). \)

\(^5\) This is because the constraints in \( \Sigma \) may not specify a \( \mathcal{C} \)-formula.
Given a set $\Delta$ of CDCs over $R$ and a horizontal decomposition $\Sigma$ of $R$, the propositional theory associated with $\Delta \cup \Sigma$ is $\Pi = \Pi_\Delta \cup \Pi_\Sigma$, where $\Pi_\Delta$ and $\Pi_\Sigma$ are the propositional theories of Definitions 2 and 5 associated with $\Delta$ and $\Sigma$, respectively. The set $\text{var}(\Pi) = \text{var}(\Pi_\Delta) \cup \text{var}(\Pi_\Sigma)$ of propositional variables occurring in $\Pi$ is partitioned into $\text{var}_\tau(\Pi) = \text{var}_\tau(\Pi_\Delta) \cup \text{var}_\tau(\Pi_\Sigma)$ and $\text{var}_v(\Pi) = \text{var}_v(\Pi_\Delta) \cup \text{var}_v(\Pi_\Sigma)$.

**Definition 6.** Let $\Delta$ be a set of CDCs over $R$ and let $\Sigma$ be a horizontal decomposition of $R$. The auxiliary theory $\Pi_{aux}$ for $\Pi = \Pi_\Delta \cup \Pi_\Sigma$ consists of the propositional formulæ in $\Pi_\Sigma$ whose r.h.s. is $\bot$ and the axioms of unique value for $\text{var}_\tau(\Pi)$.

Observe that the above is a proper extension of Definition 3: whenever $\Sigma$ is empty, the auxiliary theory for $\Pi$ coincides with the auxiliary theory for $\Pi_\Delta$.

**Example 4.** The propositional theory associated with $\Sigma$ of Example 1 is $\Pi_\Sigma = \{ \neg p_1^2 \land p_3^2 \rightarrow \bot, \tau \rightarrow v_2', \neg p_3^2 \rightarrow \bot \}$. Let $\Pi = \Pi_\Delta \cup \Pi_\Sigma$, where $\Pi_\Delta$ is the propositional theory already given in Example 2. The association between the propositional variables in $\text{var}_\tau(\Pi)$ and UTVPi is idf of Example 2 extended with $v_1' \rightarrow y_1 < 4$, and the auxiliary theory for $\Pi$ is $\Pi_{aux} = \{ \neg p_2^2 \land p_3^2 \rightarrow \bot, \neg p_3^2 \rightarrow \bot \}$.

**Definition 7.** Let $\Sigma$ be a horizontal decomposition, and let $\alpha$ be a valuation of $\text{var}_v(\Pi_\Delta)$. The $\alpha$-filtering of $\Pi_\Delta$ is the set

$$\Pi_{\alpha}^\Sigma = \{ \neg \text{idf}(v') \mid (P \rightarrow v') \in \Pi_\Delta, \alpha(P) = \top, v' \neq \bot \}$$

consisting of the negation of $\mathcal{C}$-formulæ associated with propositional variables that occur in some propositional formulæ of $\Pi_\Delta$ whose l.h.s. holds true under $\alpha$.

Observe that in (7), differently from (5), $\mathcal{C}$-formulæ are negated. This is because a counter-instance $I$ to losslessness is such that $V_i \times \cdots \times V_n \# \emptyset$ and $R^I$ has only one tuple; therefore, whenever the formula $\lambda(I)$ in the selection that defines a view symbol is satisfied by $I$, the $\mathcal{C}$-formula $\delta(I)$, if any, is not. On the other hand, the $\mathcal{C}$-formula in the consequent of a CDC must hold whenever the condition in the antecedent is satisfied.

For a valuation $\alpha$ of $\text{var}_v(\Pi_\Delta)$, the $\alpha$-filtering of $\Pi_\Delta$ is the set $\Pi_{\alpha}^\Sigma = \Pi_{\alpha}^\Delta \cup \Pi_\Sigma^\alpha$, which, as $\mathcal{C}$ is closed under negation, consists of $\mathcal{C}$-formulæ.

The main result of this section is the following characterization of lossless horizontal decomposition in terms of unsatisfiability in $\mathcal{C}$.

**Theorem 3.** Let $\Sigma$ be a horizontal decomposition of $R$, let $\Delta$ be a set of CDCs over $R$, and let $\Pi_{aux}$ be the auxiliary theory for $\Pi = \Pi_\Delta \cup \Pi_\Sigma$. Then, $\Sigma$ is lossless under $\Delta$ if and only if the $\alpha$-filtering $\Pi_{\alpha}^\Sigma = \Pi_{\alpha}^\Delta \cup \Pi_\Sigma^\alpha$ of $\Pi$ is unsatisfiable for every valuation $\alpha$ of $\text{var}_v(\Pi_\Delta)$ satisfying $\Pi_{aux}$.

Whenever the satisfiability of $\mathcal{C}$-formulæ is decidable, Theorem 3 provides an algorithm for deciding whether a given horizontal decomposition is lossless. We illustrate this in our running example with UTVPi.

**Example 5.** Consider $\Pi$ and $\Pi_{aux}$ from Example 4. The only valuation of $\text{var}_v(\Pi_\Delta)$ satisfying $\Pi_{aux}$ is $\alpha = \{ p_2^2 \rightarrow \top, p_3^2 \rightarrow \bot \}$, for which the $\alpha$-filtering of $\Pi$ is

$$\Pi_{\alpha}^\Sigma = \{ y_1 + y_2 \leq 5, y_2 \geq 2, y_1 - y_2 \geq 0 \} \cup \{ y_2 \geq 4 \}$$

Note that $y_2 \geq 4$ in $\Pi_{\alpha}^\Sigma$ is $\neg \text{idf}(v_2')$, that is, the negation of $y_2 < 4$. The set $\Pi_{\alpha}^\Sigma = \Pi_{\alpha}^\Delta \cup \Pi_\Sigma^\alpha$ is unsatisfiable because from $y_1 + y_2 \leq 5$ and $y_2 \geq 4$ we get $y_1 \leq 1$, which together with $y_1 - y_2 \geq 0$ yields $y_2 \leq 1$, in conflict with $y_2 \geq 2$. So, the horizontal decomposition $\Sigma$ is lossless under $\Delta$.

We will now give the proof of Theorem 3, for which we first need to prove two additional lemmas. In the following, and in the rest of the paper, let $\varphi$ denote the formulæ $\forall x, y. R(x, y) \iff \lor v \in V(x, y)$, and recall that a horizontal decomposition $\Sigma$ is lossless under $\Delta$ if and only if $\Delta \cup \Sigma \models \varphi$.

**Lemma 3.** Let $\Sigma$ be a horizontal decomposition of $R$ and let $\Delta$ be a set of CDCs over $R$. Then, $\Delta \cup \Sigma \models \varphi$ if and only if there exists a tuple $T$ such that the instance $I = \{ R(T) \}$ is a model of $\Delta \cup \Sigma$.

**Proof.** The “if” is trivial. For the “only if”, assume that $\Delta \cup \Sigma \models \varphi$. Then, there exists a model $J$ of $\Delta \cup \Sigma$ such that $J \models \varphi$, that is, $R^J \neq V_1 \times \cdots \times V_n$. The extension of each $V_i$ always contains a subset of the tuples in the extension of $R$ under every instance, hence there must be $\bar{T} \in R^J$ such that $\bar{T} \notin V_i$ for every $i \in \{ 1, \ldots, n \}$. Let $I = \{ R(\bar{T}) \}$; as every constraint in $\Delta \cup \Sigma$ is in one tuple and $J \models \Delta \cup \Sigma$, we have that $J \models \Delta \cup \Sigma$. □

The next lemma is more technical: intuitively, it shows that any tuple that is consistent with a valuation $\alpha$ satisfying the auxiliary theory provides a counterexample to losslessness if and only if the assignment induced by its interpreted positions satisfies the $\alpha$-filtering.

**Lemma 4.** Let $\Sigma$ be a horizontal decomposition of $R$, let $\Delta$ be a set of CDCs over $R$, and let $\Pi_{aux}$ be the auxiliary theory for $\Pi = \Pi_\Delta \cup \Pi_\Sigma$. Let $\alpha$ be a valuation of $\text{var}_v(\Pi_\Delta)$, let $\bar{T}$ be a tuple consistent with $\alpha$, and let $\beta$ be the assignment induced by the interpreted positions of $\bar{T}$. Whenever $\alpha \models \Pi_{aux}$, we have that $\{ R(\bar{T}) \} \models \Delta \cup \Sigma$ if and only if $\beta$ satisfies $\Pi_{\alpha}^\Sigma$.

**Proof.** Let $n = |R|$ and $k = ||R||$.

**Claim 1.** Let $\text{prop}(\phi) = P \rightarrow v'$, with $\phi \in \Sigma$. Then, $I \not\models \phi$ iff $\alpha(P) = \top$ and, whenever $v' \neq \bot$, $\beta \models \text{idf}(v')$.

**Proof.** Since $\phi \in \Sigma$ has the form (3), $I \not\models \phi$ iff $\lambda(I)$ is true under $\{ x_1^{[1]} \rightarrow \bar{T}[1], \ldots, x_k^{[k]} \rightarrow \bar{T}[k] \}$ and $\delta(I)$, if any, holds true under $\{ y_1 \rightarrow \bar{T}[k + 1], \ldots, y_{n-k} \rightarrow \bar{T}[n] \}$. As $v' \neq \bot$ iff $\phi$ contains a $\mathcal{C}$-formula $\delta(I) = \text{idf}(v')$, the claim follows by construction of $\alpha$ and $\beta$.

Assume $\alpha \models \Pi_{aux}$. We will show that $I \not\models \Delta \cup \Sigma$ if and only if $\beta$ does not satisfy $\Pi_{\alpha}^\Sigma$.

“$\Rightarrow$”. Assume $\beta \not\models \Pi_{\alpha}^\Sigma$. Then, there is a $\mathcal{C}$-formula $\psi \in \Pi_{\alpha}^\Sigma$ that is not satisfied by $\beta$. By construction of $\Pi_{\alpha}^\Sigma$, either $\psi$ or its negation appear in some $\phi \in \Delta \cup \Sigma$, depending on whether $\phi \in \Delta$ or $\phi \in \Sigma$, respectively. If
φ ∈ Δ, then prop(φ) = P → v with α(P) = T, so I ̸|= φ (by Claim 2 in the proof of Lemma 1). If φ ∈ Σ, then prop(φ) = P → v′ with v′ ̸|= v and α(P) = T, hence I ̸|= φ by Claim 1. In either case I ̸|= Δ ∪ Σ.

“only if”. Assume I ̸|= Δ ∪ Σ. Then, there is some φ ∈ Δ ∪ Σ that is not satisfied by I. If φ is in Δ, by Lemma 1 β |= Πaux, hence β |= Πaux. If φ is in Σ, prop(φ) = P → v′; as I ̸|= φ, by Claim 1 α(P) = T and v′ ̸|= v implies β |= idf(v′). Suppose v′ = L, then prop(φ) is in Πaux, and, since α |= Πaux, we obtain α(P) = F, which is a contradiction. So, v′ ̸|= v and β |= idf(v′). In turn, we have that β |= idf(v′) and ¬idf(v′). Therefore, β ̸|= Πaux.

Proof of Theorem 3. Let n = |R| and k = |R|L. We will show that Δ ∪ Σ ̸|= ϕ if and only if there exist α and β satisfying Πaux and Πaux, respectively.

“if”. Let α and β be such that α |= Πaux and β |= Πaux. Since α ̸|= Πaux, no two distinct propositional variables in varp(Π) associated with the same position are both true under α. Hence, there is a tuple t such that α(t) is a model of Δ ∪ Σ. Let β be the assignment induced by the interpreted positions of t, and let β be the valuation such that, for each p ∈ varp(Π), α(p) = T if p = p and t[i] = a, and α(p) = F otherwise. We will show that α satisfies Πaux and, in turn, β |= Πaux by Lemma 4, since I ̸|= Δ ∪ Σ. By construction, α satisfies every propositional formula in Πaux of the form p ∧ p with p, p ∈ varp(Π) and p ̸= p. All other propositional formulae in Πaux have the form prop(φ) = P → v, where φ is a constraint in Σ that does not contain a Π-formula σ(ϕ). As I ̸|= Δ ∪ Σ, the condition λ(π) in each such φ ∈ Σ is not true under some {x1 7 → t[1], . . . , xk 7 → t[k]}. Therefore, prop(φ) = P → v is true under α as α(P) = F by construction of α.

The unsatisfiability problem for C is the complement of the satisfiability problem for C.[19]

Proof. Let Σ = {ϕ1, . . . , ϕn} be a set C-formulae. We will show how to construct a horizontal decomposition that is lossless under Δ = Σ precisely if I is unsatisfiable. To this end, take Σ = {V i : 7 ϕi | ϕi ∈ G} and observe that, as C is closed under negation, 7 ϕi ∈ C. Thus, Σ consists of selections of the form (3), where σ = 7 ϕi and λ = T. Therefore, Σ is indeed a horizontal decomposition.

Let Π = Πaux = Πaux = Πaux = Πaux = Πaux = Πaux. Since varp(Π) = Σ, the only valuation of varp(Π) is α = Σ, which satisfies Πaux and for which Πaux = {¬idf(v) | v ∈ varp(Πaux)} = Π. Therefore, by Theorem 3, Σ is lossless under Δ = Σ if and only if I is unsatisfiable. The reduction is linear in the size of Σ.

With regard to the losslessness problem in the languages of UTVPI and BTVPIs, we have the following complexity results.

Theorem 4. When C is the language of either BTVPIs or UTVPIs, the losslessness problem in C is co-NP-complete.

Proof sketch. Constructing the propositional theories II and Πaux takes linear time, checking whether a valuation α of varp(Π) satisfies Πaux requires polynomial time, and checking that an assignment of integers to the variables in Π satisfies II (whose construction takes linear time) can be done in polynomial time, whether II consists of BTVPIs or UTVPIs. Hence, in light of Theorem 3, we can verify a given solution to the complement of the losslessness problem, in either language, in polynomial time. Therefore, the losslessness problem is co-NP in both cases.

The co-NP-hardness in the case of BTVPIs follows by Lemma 5 from the fact that the satisfiability problem for BTVPI-formulae is NP-hard and so its complement is, in turn, co-NP-hard. The co-NP-hardness in the case of UTVPIs can be shown by a reduction from UNSAT.

5 Adding FDs and UIDNs

So far, we have considered lossless horizontal decompositions under CDCs in isolation; in this section, we extend our study to the case in which the integrity constraints over the source schema are combinations of CDCs with more traditional database constraints. This investigation is vital to understand whether, how and to what extent the techniques we described in Section 4 can be applied to existing database schemas on which a set of integrity constraints other than CDCs is already defined.

Here, we focus on two well-known classes of integrity constraints, namely functional dependencies and unary inclusion dependencies [11]. Under certain restrictions—as we shall see—their interaction with CDCs can be fully captured, w.r.t. lossless horizontal decomposition, in terms of CDCs. It is important to remark that we consider restrictions solely on the CDCs, so that existing integrity constraints need not be modified in any way in order to allow for CDCs.

Let us recall that an instance I satisfies a UIDN R[i] ⊆ R[j] if every value in the ith column of R appears in the jth column of R'. The following example shows that, if we allow CDCs together with constraints from another class, such as UIDNs, their interaction may influence the losslessness of horizontal decomposition.

Example 6. Let R and V be relation symbols of arity 2, whose positions are interpreted over the integers. Let Σ be the horizontal decomposition defined by V : y1 > 3, and let Δ be a set of integrity constraints on R consisting of the CDC T → y2 > 3 and the UIDN R[1] ⊆ R[2]. It is easy to see that Δ entails T → y1 > 3. Therefore, Σ is lossless under Δ because V selects all of the tuples in R, which is clearly not the case without the UIDN.

We now introduce a general property, separability, that will constitute the main technical tool for the subsequent analysis of combinations of CDCs with FDs and UIDNs.
Informally, a class of constraints is separable from CDCs if, after making explicit the result of their interaction, which is captured by a suitable set of inference rules, we can disregard constraints from that class and focus solely on CDCs, as far as lossless horizontal decomposition is concerned. In what follows, for a set \( \Delta \) of constraints we denote by \( \text{cdc}(\Delta) \) the maximal subset of \( \Delta \) consisting solely of CDCs.

**Definition 8 (Separability).** Let \( C \) be a class of integrity constraints, let \( S \) be a finite set of sound inference rules\(^7\) for \( C \) extended with CDCs, and let \( \Delta \) consist of CDCs and \( C \)-constraints. We say that the \( C \)-constraints are \( S \)-separable in \( \Delta \) from the CDCs if every horizontal decomposition is lossless under \( \Delta \) exactly when it is lossless under \( \text{cdc}(\Delta') \), where \( \Delta' \) denotes the \( S \)-closure of \( \Delta \).\(^8\) We say that the \( C \)-constraints are separable if there is some \( S \) for which they are \( S \)-separable.

Thus, to check whether a horizontal decomposition \( \Sigma \) is lossless under an \( S \)-separable combination \( \Delta \) of CDCs and other constraints, one can proceed as follows:

1. compute the deductive closure \( \Delta' \) of \( \Delta \) w.r.t. \( S \), which makes explicit the interaction between CDCs and the other constraints in \( \Delta \) by adding entailed constraints;
2. by using the technique of Section 4, check whether \( \Sigma \) is lossless under \( \text{cdc}(\Delta') \), that is, the set obtained by discarding from \( \Delta' \) all of the constraints that are not CDCs.

Observe that \( S \)-separability implies \( S' \)-separability for every sound \( S' \supseteq S \).

### 5.1 Functional Dependencies

We begin our investigation of separability by showing that FDs do not interact with CDCs and so, as far as the losslessness of horizontal decomposition is concerned, they can be freely allowed in combination with them.

**Theorem 5.** Let \( \Delta \) be a set of CDCs and FDs. Then, the FDs are \( \emptyset \)-separable in \( \Delta \) from the CDCs.

**Proof.** We will prove that a horizontal decomposition is lossless under \( \Delta \) if and only if it is lossless under \( \text{cdc}(\Delta) \), “if”.

We have that \( \text{cdc}(\Delta) \subseteq \Delta \), and, in turn, \( \Delta \) entails \( \text{cdc}(\Delta) \); therefore \( \text{cdc}(\Delta) \models \bar{\varphi} \) implies \( \Delta \models \bar{\varphi} \).\(^9\)

“only if”. Whenever a horizontal decomposition is not lossless, by Lemma 3 there is a witness instance \( I \) with only one tuple. Since the violation of an FD involves at least two tuples, \( I \) satisfies all of the FDs in \( \Delta \).\(^10\) \( \square \)

### 5.2 Unary Inclusion Dependencies

Since in general it is not possible to compare values from \( \text{dom} \) with values from \( \text{idom} \), we consider only UINDs of the form \( R[i] \subseteq R[j] \) where positions \( i \) and \( j \) are either both non-interpreted or both interpreted. We refer to the UINDs in the former case as X-UINDs and in the latter as Y-UINDs.

Let \( n = |R| \) and \( k = |R| \); we write \( R[x]_i \subseteq R[x]_j \) with \( i, j \in \{1, \ldots, k\} \) to denote the X-UIND \( R[i]_i \subseteq R[j]_j \) and we write \( R[y]_i \subseteq R[y]_j \) with \( i, j \in \{1, \ldots, n - k\} \) to denote the Y-UIND \( R[i] \subseteq R[j] \).

#### 5.2.1 UINDs on Interpreted Attributes

First, we study the interaction between Y-UINDs (that is, UINDs at interpreted positions) and a restricted form of CDCs, which we shall introduce shortly. This interaction is captured by the following domain propagation rule:

\[
\mathbb{T} \rightarrow \delta(y_i) \quad R[y]_i \subseteq R[y],
\]

whose soundness is easily shown below.

**Theorem 6.** Let \( \Delta \) be a set of CDCs and UINDs. If \( \Delta \models \forall \varphi. R(\bar{x}, \bar{y}) \not\rightarrow \delta(\bar{y}) \) and \( \Delta \models R[y]_i \subseteq R[y] \), then \( \Delta \models \forall \varphi. R(\bar{x}, \bar{y}) \not\rightarrow \delta(\bar{y}) \).

**Proof.** If \( \Delta \) is inconsistent, the claim follows trivially. Thus, let \( I \) be a model of \( \Delta' \); hence \( I \) satisfies the CDC \( \forall \varphi. R(\bar{x}, \bar{y}) \not\rightarrow \delta(\bar{y}) \) and the UIND \( R[y]_i \subseteq R[y] \). If \( R' = \emptyset \), then trivially \( I \models \forall \varphi. R(\bar{x}, \bar{y}) \not\rightarrow \delta(\bar{y}) \). So, let \( R' \not= \emptyset \) and suppose \( I \not= \forall \varphi. R(\bar{x}, \bar{y}) \not\rightarrow \delta(\bar{y}) \). Then, there exists \( \bar{t} \in R' \) for which \( \delta(\bar{t}[i + k]) \) holds true, with \( k = |R| \). By the UIND, there must be \( \bar{t}' \in R' \) such that \( \bar{t}'[i + k] = \bar{t}[j + k] \). Hence \( \delta(\bar{t}'[i + k]) \) is not true, in contradiction of \( I \models \forall \varphi. R(\bar{x}, \bar{y}) \not\rightarrow \delta(\bar{y}) \). \( \square \)

It turns out that when all of the CDCs that mention a variable \( y \) corresponding to an interpreted position affected by some Y-UIND have the form \( \mathbb{T} \rightarrow \delta(y) \), the domain propagation rule fully captures the interaction between such CDCs and Y-UINDs w.r.t. losslessness.

**Definition 9.** We say that a set \( \Delta \) of CDCs and Y-UINDs is dp-controllable if, for every Y-UIND \( R[y]_i \subseteq R[y]_j \) in \( \Delta \) with \( i \not= j \), all of the CDCs in \( \Delta \) mentioning the variable \( y \), where \( y \) is \( y_i \) or \( y_j \), are of the form \( \mathbb{T} \rightarrow \delta(y) \).

**Theorem 7.** Let \( \Delta \) be a dp-controllable set of CDCs and Y-UINDs. Then, the Y-UINDs are \( \{((dp)\}_i\} \)-separable in \( \Delta \) from the CDCs.

The above theorem is a special case of a more general result (Theorem 10) given later on.

Even though in general dp-controllability is not a necessary condition for the \( \{((dp)\}_i\} \)-separability of Y-UINDs from CDCs, the following examples show two different situations where, in the absence of dp-controllability, the Y-UINDs are not \( \{((dp)\}_i\} \)-separable from the CDCs.

**Example 7.** Let \( R \) be a ternary relation symbol, whose last two positions are interpreted over the integers. Let \( \Delta \) consist of the Y-UIND \( R[y]_1 \subseteq R[y]_2 \) and of the CDCs \( x_1 = a \rightarrow y_2 > 2, x_1 \not= a \rightarrow y_1 > 0, x_1 \not= a \rightarrow y_1 < 0 \), and consider the view symbol \( V : y_1 > 1 \). For \( x_1 \not= a \) there is no suitable value for \( y_1 \) to satisfy the above CDCs, thus every model \( I \) of \( \Delta \) is such that, for every \( \bar{t} \in R'^{a} \), \( \bar{t}[1] = a \) and \( \bar{t}[3] > 2 \). Moreover, by the Y-UIND \( R[y]_1 \subseteq R[y]_2 \), we also have that \( \bar{t}[2] > 2 \), and therefore every tuple in \( R'^{a} \) is...
also in $V'$, which means that $V$ is lossless under $\Delta$. Clearly, this is not the case in the absence of the Y-UIND, that is, under $\text{cdc}(\Delta)$. Let $\Delta'$ be the $\{(dp)\}$-closure of $\Delta$. Then, as $\Delta' = \Delta$, we have that $V$ is lossy under $\text{cdc}(\Delta')$ and, therefore, the Y-UIND is not $\{(dp)\}$-separable in $\Delta$ from the CDCs.

Example 8. Let $R$ be a relation symbol of arity 4 and with all of its positions interpreted over the integers. Consider the view symbol $V: y_3 < 3 \land y_1 > 4$, and let $\Delta$ consist of the Y-UIND $R[y_1] \subseteq R[y_2]$ and the CDCs $T \rightarrow y_1 + y_3 > 0$, $T \rightarrow y_2 + y_4 < 0$, $T \rightarrow y_3 - y_4 \leq 0$. The above CDCs entail $T \rightarrow y_1 - y_2 > 1$, thus in every model $I$ of $\Delta$ each tuple $t \in R' \setminus I$ must satisfy $t[1] - t[2] > 1$.

By the Y-UIND $R[y_1] \subseteq R[y_2]$, for each $d$ in $\pi_2(R')^{11}$ there exists $d' \in \pi_2(R')$ with $d' \geq d + 1$. Then, as $d' \neq d$, the instance $I$ is either infinite or empty. Hence, every horizontal decomposition is lossless under $\Delta$.

On the other hand, let $\Delta'$ be the $\{(dp)\}$-closure of $\Delta$ and observe that $\Delta' = \Delta$. Let $J = \{(R(1,0,0,0))\}$; then, since $J \models \text{cdc}(\Delta')$, $V$ is lossy under $\text{cdc}(\Delta')$. Therefore, the Y-UIND is not $\{(dp)\}$-separable in $\Delta$ from the CDCs.

5.2.2 UINDs on Non-Interpreted Attributes

We now turn our attention to combinations of CDCs and X-UINDs (i.e., UINDs at non-interpreted positions). First, we show that the syntactic restrictions introduced in [16] on the CDCs are not sufficient for the $\varnothing$-separability of X-UINDs. Indeed, the following is a counterexample to Theorem of [16].

Example 9. Let $R$ be a ternary relation symbol, with the third position interpreted over the integers. Let $\Delta$ consist of the CDC $x_2 = a \rightarrow y_1 \leq 0 \land y_1 > 0$ and the X-UIND $R[1] \subseteq R[2]$. The CDCs in $\Delta$ are trivially globally consistent. Then, the X-UINDs are $\varnothing$-separable in $\Delta$ from the CDCs.

Example 10. Let $R$ be a ternary relation symbol, whose third position is interpreted over the integers, and let $\psi$ be the X-UIND $R[2] \subseteq R[1]$. The set $\Delta_1$ consisting of the CDCs $x_1 = a \rightarrow y_1 < 0$ and $x_1 = a \rightarrow y_1 > 0$ is not globally consistent, as there is no suitable value for the

Theorem 8. Let $\Delta$ consist of CDCs and X-UINDs such that $\text{cdc}(\Delta)$ is globally consistent. Then, the X-UINDs are $\varnothing$-separable in $\Delta$ from the CDCs.

The above theorem, like Theorem 7, is a special case of a more general result (Theorem 10) given later on.

It is possible to check for the global consistency of a set $\Delta$ of CDCs in a way similar to the one described in Theorem 1 for consistency, by building the propositional theory $\Pi_\Delta$ associated with $\Delta$, along with the auxiliary theory $\Pi_{\text{aux}}$ for $\Pi_\Delta$, and then checking that the $\alpha$-filtering $\Pi_{\text{aux}}^\alpha$ of $\Pi_{\text{aux}}$ is satisfiable for every (rather than just for one) valuation $\alpha$ of $\vartheta_{\text{aux}}(\Pi_{\text{aux}})$ that satisfies $\Pi_{\text{aux}}$. Indeed, under the assumptions of Theorem 1, $\Delta$ is globally consistent if and only if $\Pi_{\text{aux}}^\alpha$ is satisfiable for every valuation $\alpha$ of $\vartheta_{\text{aux}}(\Pi_{\text{aux}})$ satisfying $\Pi_{\text{aux}}$. Checking for global consistency is expensive, because it requires an exponential number of satisfiability checks in $\epsilon$; the associated decision problem is in PSPACE (the space used for one satisfiability check can be reused for the next) for UTVPVs as well as for BUTVPVs.

Devising purely syntactic restrictions that guarantee the global consistency of CDCs depends on the specific constraint language $\epsilon$ in use, which is indeed what we overlooked in [16]. As it turns out, the non-overlapping and partition-free restrictions of [16] ensure global consistency (and so also the $\varnothing$-separability of X-UINDs) only for sets of CDCs whose consequents are UTVPVs. This is not the case anymore for CDCs whose consequents are BUTVPVs, which indeed allow to express Example 9.

We provide a condition that, although not guaranteeing global consistency, ensures the $\varnothing$-separability of the X-UINDs. Moreover, this restriction can be checked more efficiently than global consistency, as it requires only a polynomial number of $\epsilon$-satisfiability checks.

Definition 11. Let $\Delta$ be a set of CDCs. We say that the CDCs in $\Delta$ are disjoint w.r.t. an X-UIND $R[x] \subseteq R[x]$ if for any two CDCs $\phi_1(\overline{x}_1, \overline{y}_1)$ and $\phi_2(\overline{x}_2, \overline{y}_2)$ in $\Delta$, with $x_i$ in $\overline{x}_i$, the consequent of $\phi_1$ is satisfiable and has no variables in common with the consequent of $\phi_2$.

Intuitively, the above requires that all of the variables appearing in the consequent of any CDC $\phi$ whose antecedent mentions the variable $x_i$, affected by an X-UIND $R[x_i] \subseteq R[x_i]$, do not occur in the consequent of any other CDC; moreover, the consequent of each such $\phi$ must be satisfiable.

Theorem 9. Let $\Delta$ be a set of CDCs and X-UINDs, where the CDCs are disjoint w.r.t. each X-UIND in $\Delta$. Then, the X-UINDs are $\varnothing$-separable in $\Delta$ from the CDCs.

The above theorem is a special case of a more general result (Theorem 11) given later on in Section 5.3. Clearly, as global consistency is a property of the CDCs in isolation, whereas disjointness is relative to a X-UIND, these two notions are incomparable, in the sense that one does not imply the other and vice versa, as shown below.

Example 11. Let $R$ be a ternary relation symbol, whose third position is interpreted over the integers, and let $\psi$ be the X-UIND $R[2] \subseteq R[1]$. The set $\Delta_1$ consisting of the CDCs $x_1 = a \rightarrow y_1 < 0$ and $x_1 = a \rightarrow y_1 > 0$ is not globally consistent, as there is no suitable value for the
third position (associated with \(y_1\)) whenever the value of the first (associated with \(x_1\)) is \(a\), but the CDCs in \(\Delta_1\) are disjoint w.r.t. \(\psi\). On the other hand, the set \(\Delta_2\) consisting of \(x_1 = a \rightarrow y_1 > 0\) and \(x_2 = a \rightarrow y_1 > 1\) is not disjoint w.r.t. \(\psi\) (because both CDCs mention a variable affected by \(\psi\) and their consequences have \(y_1\) in common) but \(\Delta_2\) is globally consistent.

5.2.3 UIINDs on All Attributes

We now study the separability of UIINDs (i.e., X-UIINDs and Y-UIINDs together)\(^{12}\) from CDCs. The following is a generalisation of both Theorem 7 and Theorem 8.

**Theorem 10.** Let \(\Delta\) be a set of globally consistent CDCs, X-UIINDs and Y-UIINDs, such that the CDCs and Y-UIINDs are dp-controllable. Then, the UIINDs are \(\{(dp)\}\)-separable in \(\Delta\) from the CDCs.

To give the proof of the above theorem, we will need to prove several lemmas, showing how any given model of the (saturated set of) CDCs can be extended in order to satisfy the UIINDs as well.

**Lemma 6.** Let \(\Delta\) be a dp-controllable set of CDCs and Y-UIINDs, and let \(\bar{\tau}\) be a tuple such that \(\{R(\bar{\tau})\} \models \text{cdc}(\Delta^*)\), where \(\Delta^*\) is the \(\{(dp)\}\)-closure of \(\Delta\). Let \(\psi = R^i[j] \subseteq R^j[j]\) be a Y-UIIND in \(\Delta\), and let \(\bar{\tau}'\) be identical to \(\bar{\tau}\) except for \(\bar{\tau}'[j] = \bar{\tau}[i]\). Then, \(\{R(\bar{\tau}')\} \models \text{cdc}(\Delta^*) \land \{\psi\}\).

**Proof.** Let \(\Delta' = \text{cdc}(\Delta^*)\). Since all of the UIINDs in \(\Delta\) are Y-UIINDs, \(i, j > k\) with \(k = |R^i|\). As \(\bar{\tau}\) satisfies \(\Delta'\) and \(\bar{\tau}'\) differs from \(\bar{\tau}\) only on the \(j\)th element, \(\bar{\tau}'\) satisfies every CDC in \(\Delta'\) not mentioning the variable \(y_{j-k}\). The only CDCs in \(\Delta'\) which are allowed to mention \(y_{j-k}\) have the form \(\top \rightarrow \delta(y_{j-k})\). For each such CDC, since \(R^i[j] \subseteq R^j[j]\) is in \(\Delta'\), by (dp) also \(\top \rightarrow \delta(y_{j-k})\) is in \(\Delta'\). Hence \(\delta(\bar{\tau}[i])\) holds true, and in turn \(\delta(\bar{\tau}'[j])\) is true as well, because \(\bar{\tau}'[j] = \bar{\tau}[i]\). Therefore, \(\bar{\tau}'\) satisfies all the CDCs of the form \(\top \rightarrow \delta(y_{j-k})\). Moreover, \(\bar{\tau}'\) trivially satisfies the UIIND \(\psi\), as \(\bar{\tau}'[i] = \bar{\tau}'[j] = \bar{\tau}[i]\).

**Lemma 7.** Let \(\Delta\) be a dp-controllable set of CDCs and Y-UIINDs, and let \(I\) be a model of \(\text{cdc}(\Delta^*)\), where \(\Delta^*\) is the \(\{(dp)\}\)-closure of \(\Delta\). Then, there exists an instance \(J \supseteq I\) such that \(J \models \Delta^*\).

**Proof sketch.** Let \(J_0 = I\). We will iteratively add tuples to \(J_0\) so as to obtain a model of \(\Delta^*\). At each iteration \(k\), proceed as follows:

1. Find a violation of some X-UIIND \(R^i[j] \subseteq R^j[j]\) in \(\Delta^*\), that is, a value \(d \in \pi_i(R^k)\) which is not in \(\pi_j(R^k)\).
2. Take \(\bar{\tau} \in R^k\) such that \(\bar{\tau}[i] = d\) and \(\bar{\tau}[j] \neq d\).
3. Let \(J_{k+1} = J_k \cup \{R(\bar{\tau})\}\), with \(\bar{\tau}'\) identical to \(\bar{\tau}\) except for \(\bar{\tau}'[j] = d\).

By construction, \(J\) satisfies all of the Y-UIINDs in \(\Delta\) and, since at each step \(\{R(\bar{\tau})\} \models \text{cdc}(\Delta)\) by Lemma 6, it can be shown that \(J \models \text{cdc}(\Delta)\) by an easy induction.

12. Recall that UIINDs between non-interpreted and interpreted positions are not allowed, as they make little sense.
model of \( \text{cdc}(\Delta) \). Then, there exists an instance \( J \) such that \( J \models I \) and \( J \models \Delta \).

**Proof sketch.** The construction of \( J \) is the same as in Lemma 8, with the difference that in step 3 suitable values for \( \overline{t} \) at interpreted positions exist by the disjointness of the CDCs w.r.t. the X-UINDs.

5.3 FDs and UINDs Together

Unfortunately, the separability results presented above for combinations of CDCs and UINDs do not automatically carry over to the case in which FDs are also present. In fact, although FDs do not directly interact with CDCs, they do in general interact with UINDs,\(^{13}\) which in turn interact with CDCs.

We write FDs over \( R \) as implications between sets of positions of \( R \) (e.g., \( \{1,3\} \rightarrow \{4\} \)). We call X-FD (resp., Y-FD) an FD whose l.h.s. and r.h.s. both consist of non-interpreted (resp., interpreted) positions; and we call XY-FD (resp., YX-FD) an FD where the l.h.s. consists of non-interpreted (resp., interpreted) positions and the r.h.s. of interpreted (resp., non-interpreted) ones.

The following generalises Theorem 7 in the presence of X-FDs and YX-FDs.

**Theorem 12.** Let \( \Delta \) be a set of CDCs, X-UINDs, X-FDs and YX-FDs, where the CDCs and Y-UINDs are dp-controllable. Then, the X-FDs, YX-FDs and Y-UINDs are \((\text{dp})\)-separable in \( \Delta \) from the CDCs.

**Proof sketch.** The proof of Theorem 10 can be modified as follows: in the “only if” direction take \( J' = I \), which contains only the tuple \( \overline{t} \), and construct \( J'' \) as in Lemma 7 (with \( \Delta' = \Delta \)), by extending \( J'' \) with tuples that have the same values as \( \overline{t} \) at non-interpreted positions. Therefore, \( J'' \) satisfies any FD whose r.h.s. is a set of non-interpreted positions.

Theorem 8 does not hold anymore in the presence of Y-FDs, that is, X-UINDs and Y-FDs are not \( \otimes \)-separable in general from globally consistent CDCs, as shown below.

**Theorem 13.** There is a set of X-UINDs, Y-FDs and globally consistent CDCs, in which the X-UINDs and Y-FDs are not \( \otimes \)-separable from the CDCs.

**Proof.** Let \( R \) be a relation symbol of arity 4, whose last two positions are interpreted over the integers. Let \( \Delta \) consist of the X-UIND \( R[1] \subseteq R[2] \), the Y-FD \( R : \{3\} \rightarrow \{4\} \), and the CDCs \( x_1 = a \land x_2 = b \rightarrow y_1 = 0 \land y_2 > 1 \) and \( x_2 = a \rightarrow y_1 = 0 \land y_2 < 1 \). These CDCs are globally consistent, because their consequents are satisfiable and their antecedents are never true at the same time (as \( x_2 \) cannot be simultaneously equal to \( b \) and \( a \)). Let \( \Sigma \) be the horizontal decomposition specified by \( V_1 : x_1 \neq a \) and \( V_2 : x_2 \neq b \). Clearly, \( \Sigma \) is lossy under \( \text{cdc}(\Delta) \) as the instance \( I = \{ R(a,b,0,2) \} \) satisfies \( \text{cdc}(\Delta) \) and \( \Sigma \); indeed, the tuple \((a,b,0,2)\) is in \( R^I \) but it is not selected by either \( V_1 \) or \( V_2 \). Suppose that \( \Sigma \) is lossy under \( \Delta \). Then, there exists a model \( J \) of \( \Delta \cup \Sigma \) and a tuple \( \overline{t} \in R^J \) such that \( \overline{t} \notin V_1^J \cup V_2^J \). By definition of \( V_1 \) and \( V_2 \), we have that \( t[1] = a \) and \( t[2] = b \) and, in turn, \( t[3] = 0 \) and \( t[4] > 1 \) by the first CDC. By the X-UIND, there must be \( \overline{t}' \in R^I \) such that \( \overline{t}'[2] = a \) and, in turn, \( \overline{t}'[3] = 0 \) and \( \overline{t}'[4] < 1 \) by the second CDC. But then, \( \overline{t} \) and \( \overline{t}' \) violate the Y-FD, since they agree on the third position but must differ on the fourth. So \( J \not\models \Delta \), which is a contradiction. Hence, \( \Sigma \) is lossless under \( \Delta \), and we conclude that the X-UIND and Y-FD are not \( \otimes \)-separable from the CDCs.

\(^{13}\) The interaction between FDs and UINDs can be fully captured, as there is a sound and complete axiomatization for the finite implication of FDs and UINDs [11].

6 DISCUSSION AND OUTLOOK

In this paper, we studied lossless horizontal decomposition under constraints in a setting where the values for some of the attributes in the schema are taken from an interpreted domain. Data values in such a domain can be compared in ways beyond equality, according to a first-order language \( \mathcal{C} \). We did not make any assumption on \( \mathcal{C} \), other than requiring it to be closed under negation.

In the above setting, we considered a class of integrity constraints, CDCs, based on those introduced in [16]. We have characterised the consistency of a set of CDCs in terms of satisfiability in \( \mathcal{C} \) and we have shown that the problem of deciding consistency is \( \text{NP} \)-complete when \( \mathcal{C} \) is the language of either UTVPIS or BUTVPIS. We considered a more general form of selections than in [16] and characterised, in terms of unsatisfiability in \( \mathcal{C} \), whether a horizontal decomposition specified by such selections is lossless under CDCs. We have shown that the problem of deciding losslessness is \( \text{co-NP} \)-complete when \( \mathcal{C} \) is the language of either UTVPIS or BUTVPIS.

We also considered losslessness under CDCs in combination with FDs and UINDs. We introduced and
studied the important notion of separability, which indicates whether constraints other than CDCs can be disregarded w.r.t. losslessness, after incorporating the effect of their interaction in terms of entailed CDCs. A summary of all the separability results presented in this paper is given in Table 1.

A promising direction for future research we are currently investigating is the generalisation of the separability results for UINDs to arbitrary inclusion dependencies (INDs). Observe that INDs, differently from UINDs, can affect both interpreted and non-interpreted attributes at the same time, e.g., in $R[x_1,y_1] \subseteq R[x_2,y_2]$. Some care is needed in allowing FDs in this setting as well, because logical implication for unrestricted combinations of FDs and INDs is undecidable and has no axiomatization [11].

Another interesting direction is that of allowing equalities between two variables in the antecedents of CDCs as well as in the selection conditions on non-interpreted attributes of view definitions. We believe our approach could be extended in this direction by representing such equalities by propositional variables and by adding suitable axioms to the auxiliary theory to handle transitivity and symmetry.

The main motivation for our study of lossless horizontal decomposition is that it provides the groundwork for the consistent and unambiguous propagation of updates in the context of selection views. By applying the general criterion of [6], given a lossless horizontal decomposition it is possible to determine whether an update issued on some (possibly all) of the fragments can be propagated to the underlying database without affecting the other fragments. Similarly, it is possible to partition the source relation by adding suitable conditions in the selections that define the fragments, so that each is disjoint with the others. In general, a lossy horizontal decomposition can always be turned into a lossless one by defining an additional fragment, called a complement, which selects the missing tuples. In particular, there is a unique minimal complement selecting all and only the rows of the source relation that are not selected by any of the other fragments. In follow-up work, we will show how to compute the definition of such a complement, in the scope of an in-depth study of partitioning and update propagation in the setting studied in this paper.

Most of the work in the field of horizontal decomposition has been carried out in the context of distributed databases systems, where one is mainly concerned with finding an optimal decomposition w.r.t. some parameters (e.g., workload, query-execution time, storage quotas), rather than determining whether a given horizontal decomposition is lossless.

De Bra ([12], [13]) developed a theory of horizontal decomposition to partition a relation into two sub-relations such that one satisfies certain FDs that the other does not. The approach is based on constraints that capture partial implications between sets of FDs and exceptions to sets of FDs, for which a sound and complete set of inference rules is provided. These constraints are $\emptyset$-separable from our CDCs (for the same reason FDs are).

Maier and Ullman [15] consider horizontal decomposition involving physical and virtual fragments over the same attributes. Fragments are defined in an arbitrary (first-order) language closed under Booleans, where entailment is decidable and consisting of formulae that, as in our case, can be evaluated by examining one tuple at a time, in isolation from the others. Differently from our case, the language allows to express equalities between variables associated with non-interpreted attributes. But, if such equalities are forbidden, the setting of [15] can be recast into ours: the union of the physical fragments is the single source relation $R$ we consider here, the definitions of the physical fragments can be taken as integrity constraints over $R$, and the definition of each virtual fragment (given in terms of the physical fragments and other virtual ones) can be expressed only in terms of $R$ by query unfolding. Then, the problem of determining whether the virtual fragments constitute a lossless horizontal decomposition of the physical fragments, which is not addressed in [15], can be solved by applying the techniques we described in this paper. Virtual fragments in [15] are defined by selection and union, that is, in our notation, by formulae of the form $\lambda(x) \land \sigma(y)$ or $\lambda(x) \lor \sigma(y)$. As we remarked in Section 2, in such a case losslessness can be checked by considering two views $\lambda(x)$ and $\sigma(y)$ in place of each view of the latter form.

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REFERENCES


TABLE 1

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Summary of $S$-Separability Results (unr = Unrestricted, dpc = dp-Controllable, gc = Globally Consistent, dis = Disjoint)


Ingo Feinerer received the PhD degree in computer science from the Vienna University of Technology and the PhD degree in business administration from the Vienna University of Economics and Business. He has been an assistant professor at the Vienna University of Technology since 2008. In 2007, he was a postdoctoral visiting scholar in the Computer Science Department, Carnegie Mellon University. His expertise and research interests comprise configuration, formal methods, text mining, and databases.

Enrico Franconi is an associate professor at the Free University of Bozen-Bolzano, Italy, where he is the director of the KRDB Research Centre for Knowledge and Data. His main research interest is on knowledge representation and reasoning technologies applied to databases, and in particular description logics, temporal representations, conceptual modeling, intelligent access to information, and natural language interfaces. Recently, he has been involved as a principal investigator in many European projects (SEWA-SIE, DWQ), networks of excellence (KnowledgeWeb, InterOp), and Marie Curie actions (NET2). He is a member of the Advisory Committee of the World Wide Web Consortium (W3C). He is in the editorial board of the Journal of Applied Logic (JAL), published by Elsevier. He is the author of more than 150 publications in international workshops, conferences, and journals.

Paolo Guagliardo completed the European Master in computational logic in 2009, received the MSc degree in computer science with honors from the Vienna University of Technology and the Free University of Bozen-Bolzano. He is currently working toward the PhD degree at the Free University of Bozen-Bolzano under the supervision of Enrico Franconi. His research interests include database theory, knowledge representation, and Semantic Web.

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