Parity Objectives in Countable MDPs

Citation for published version:

Link:
Link to publication record in Edinburgh Research Explorer

Document Version:
Peer reviewed version

Published In:
Proceedings of Thirty-Second Annual ACM/IEEE Symposium on Logic in Computer Science (LICS 2017)

General rights
Copyright for the publications made accessible via the Edinburgh Research Explorer is retained by the author(s) and / or other copyright owners and it is a condition of accessing these publications that users recognise and abide by the legal requirements associated with these rights.

Take down policy
The University of Edinburgh has made every reasonable effort to ensure that Edinburgh Research Explorer content complies with UK legislation. If you believe that the public display of this file breaches copyright please contact openaccess@ed.ac.uk providing details, and we will remove access to the work immediately and investigate your claim.
Parity Objectives in Countable MDPs

Stefan Kiefer*, Richard Mayr†, Mahsa Shirmohammadi*, Dominik Wojtczak‡

*University of Oxford, UK
†University of Edinburgh, UK
‡University of Liverpool, UK

Abstract—We study countably infinite MDPs with parity objectives, and special cases with a bounded number of colors in the Mostowski hierarchy (including reachability, safety, Büchi and co-Büchi).

In finite MDPs there always exist optimal memoryless deterministic (MD) strategies for parity objectives, but this does not generally hold for countably infinite MDPs. In particular, optimal strategies need not exist.

For countable infinite MDPs, we provide a complete picture of the memory requirements of optimal (resp., $\epsilon$-optimal) strategies for all objectives in the Mostowski hierarchy.

In particular, there is a strong dichotomy between two different types of objectives. For the first type, optimal strategies, if they exist, can be chosen MD, while for the second type optimal strategies require infinite memory. (I.e., for all objectives in the Mostowski hierarchy, if finite-memory randomized strategies suffice then also MD-strategies suffice.) Similarly, some objectives admit $\epsilon$-optimal MD-strategies, while for others $\epsilon$-optimal strategies require infinite memory. Such a dichotomy also holds for the subclass of countably infinite MDPs that are finitely branching, though more objectives admit MD-strategies here.

Index Terms—countable MDPs, parity objectives, strategies, memory requirements

I. INTRODUCTION

Markov decision processes (MDPs) are a standard model for dynamic systems that exhibit both stochastic and controlled behavior [23]. The system starts in the initial state and makes a sequence of transitions between states. Depending on the type of the current state, either the controller gets to choose an enabled transition (or a distribution over transitions), or the next transition is chosen randomly according to a defined distribution. By fixing a strategy for the controller, one obtains a probability space of plays of the MDP. The goal of the controller is to optimize the expected value of some objective function on the plays of the MDP. The fundamental questions are “what is the optimal value that the controller can achieve?”, “does there exist an optimal strategy, or only $\epsilon$-optimal approximations?”, and “which types of strategies are optimal or $\epsilon$-optimal?”.

Such questions have been studied extensively for finite MDPs (see e.g. [10] for a survey) and also for certain types of countably infinite MDPs [23], [21]. However, the literature on countable MDPs is mainly focused on objective functions defined w.r.t. numeric costs (or rewards) that are assigned to transitions, e.g. (discounted) expected total reward or limit-average reward. In contrast, we study qualitative objectives that are expressed by Parity conditions and which are motivated by formal verification questions.

There are works that studied particular classes of countably infinite, but finitely branching, MDPs that arise from models in automata theory [14], [2], [8], [6], [1]. In each of these papers, a crucial part of the analysis is establishing the existence of optimal strategies of particular structure and memory requirements, but none of them looked at proving such properties for general countable MDPs. Countable MDPs also naturally occur in the analysis of queueing systems [17], gambling [4], and branching processes [22], which have multiple applications. They also show up in the analysis of finite-state models, e.g. in two-player stochastic games [24], [13] when reasoning about an optimal strategy against a fixed (randomised and memory-full) strategy of the opponent.

Finite MDPs vs. Infinite MDPs: It should be noted that many standard properties (and proof techniques) of finite MDPs do not carry over to infinite MDPs.

E.g., given some objective, consider the set of all states in an MDP that have nonzero value. If the MDP is finite then this set is finite and thus there exists some minimal nonzero value. This property does not carry over to infinite MDPs. Here the set of states is infinite and the infimum over the nonzero values can be zero. As a consequence, even for a reachability objective, it is possible that all states have value $> 0$, but still the value of some states is $< 1$. Such phenomena appear already in infinite-state Markov chains like the classic Gambler’s ruin problem with unfair coin tosses in the player’s favor (0.6 win, 0.4 lose). The value, i.e., the probability of ruin, is always $> 0$, but still $< 1$ in every state except the ruin state itself; cf. [15] (Chapt. 14). Another difference is that optimal strategies need not exist, even for qualitative objectives like reachability or parity. Even if some state has value 1, there might not be any single strategy that attains the value 1, but only an infinite family of $\epsilon$-optimal strategies for every $\epsilon > 0$.

Parity objectives: We study general countably infinite MDPs with parity objectives. Parity conditions are widely used in temporal logic and formal verification, e.g., they can express $\omega$-regular languages and modal $\mu$-calculus [16]. Every state has a color, out of a finite set of colors encoded as natural numbers. An infinite play is winning iff the highest color that is seen infinitely often in the play is even. The controller wants to maximize the probability of winning plays. Subclasses of parity objectives are defined by restricting the set of used colors; these are classified in the Mostowski hierarchy [20] which includes, e.g., Büchi and co-Büchi objectives. Such prefix-independent infinitary objectives cannot generally be
encoded by numeric transition rewards as in [23], though both
types subsume the simpler reachability and safety objectives.

There are different types of strategies, depending on whether
one can take the whole history of the play into account
(history-dependent; (H)), or whether one is limited to a finite
amount of memory (finite memory; (F)) or whether deci-
sions are based only on the current state (memoryless; (M)).
Moreover, the strategy type depends on whether the controller
can randomize (R) or is limited to deterministic choices (D).
The simplest type MD refers to memoryless deterministic
strategies.

The type of strategy needed for an optimal (resp. \(\epsilon\)-optimal)
strategy for some objective is also called the strategy com-
plexity of the objective. For finite MDPs, MD-strategies are
sufficient for all types of qualitative and quantitative parity
objectives [9], [11], but the picture is more complex for
countably infinite MDPs.

Since optimal strategies need not exist in general, we
consider both the strategy complexity of \(\epsilon\)-optimal strategies,
and the strategy complexity of optimal strategies under the
assumption that they exist. E.g., if an optimal strategy exists,
can it be chosen MD?

We provide a complete picture of the memory requirements
for objectives in the Mostowski hierarchy, which is summa-
rized in Figure 1.

In particular, our results show that there is a strong
dichotomy between two different classes of objectives. For
objectives of the first class, optimal strategies, where they exist,
can be chosen MD. For objectives of the second class, optimal
strategies require infinite memory in general, in the sense that
all FR-strategies achieve the objective only with probability
zero. A similar dichotomy applies to \(\epsilon\)-optimal strategies. For
certain objectives, \(\epsilon\)-optimal MD-strategies exist, while for
all others even \(\epsilon\)-optimal strategies require infinite memory
in general. This is a strong dichotomy because there are no
objectives in the Mostowski hierarchy for which other types of
strategies (MR, FD, or FR) are both necessary and sufficient.
Put differently, for all objectives in the Mostowski hierarchy,
if FR-strategies suffice then MD-strategies suffice as well.

We also consider the subclass of countable MDPs that are
finitely branching. (Note that these generally still have an
infinite number of states.) The above mentioned dichotomies
apply here as well, though the classes of objectives where
optimal (resp. \(\epsilon\)-optimal) strategies can be chosen MD are
larger than for general countable MDPs.

Outline of the results: In Section II we define countably
infinite MDPs, strategies and parity objectives. In Section III
we show examples that demonstrate that certain objectives
require infinite memory. For some of these we refer to previous
work. The main new result in this section is Theorem 1
that shows that even almost-sure \{1, 2, 3\}-Parity on finitely
branching MDPs requires infinite memory. These negative re-
results highlight the questions which other objectives still allow
MD-strategies. Apart from the case of reachability objectives,
these questions were open. We provide complete answers in
several steps. First, in Section IV, we prove a general result
(Theorem 5) that relates the strategy complexity of almost-
sure winning strategies and optimal strategies. The complexity
of the proof is due to the fact that we consider infinite
MDPs (which do not satisfy basic properties of finite MDPs
in general; see above). We then use this theorem to estab-
lish MD-strategies for Büchi, co-Büchi and \{0, 1, 2\}-Parity

Fig. 1: For countable MDPs, these diagrams show the memory requirements of optimal and \(\epsilon\)-optimal strategies for objectives in
the Mostowski hierarchy. An objective in a level of the hierarchy subsumes all objectives in lower levels, e.g., \{0, 1\}-Parity
subsumes \{1, 2\}-Parity. We have extended the Mostowski hierarchy to include reachability and safety. The magenta (resp.,
blue) regions enclose objectives where memoryless deterministic (MD) strategies are sufficient for optimal (resp., \(\epsilon\)-optimal)
strategies; for objectives outside the regions, infinite-memory strategies are necessary. The left diagram is for infinitely branching
MDPs; e.g., \(\epsilon\)-optimal strategies for all but reachability objectives require infinite memory, whereas MD-strategies are sufficient
for reachability. The right diagram is for finitely branching MDPs; e.g., optimal strategies (if they exist) can be chosen MD
for all objectives subsumed by \{0, 1, 2\}-Parity.
We say that (partial) play transducers a strategy. Strategies can be implemented by probabilistic σ subscript and write Σ and that a random state set (the memory of the strategy), all σ: since this paper focuses on the memory requirements of strategies. We assume familiarity with the syntax and semantics of the temporal logic LTL [12]. Formulas are interpreted on the structure (S,→). We use [ϕ]s ⊆ Sω to denote the set of plays starting from s that satisfy the LTL formula ϕ. This set is measurable [25], and we just write P_M,ω,σ([ϕ]) instead of P_M,ω,σ([ϕ]). We also write [ϕ] for ∪s∈S [ϕ]s.

Given a target set T ⊆ S, the reachability objective is defined by Reach(T) = [F T], i.e., s0s1⋯ ∈ Reach(T) ⇔ ∃i, s_i ∈ T. The safety objective is defined by Safety(T) = [G¬T], i.e., s0s1⋯ ∈ Safety(T) ⇔ ∀i, s_i /∈ T. Given a reachability or a safety objective, we can assume without loss of generality that T is a sink in M.

Let C ⊆ N be a finite set of colors. A color function Col: S → C assigns to each state s its color Col(s). For n ∈ N, ⊢ ∈ {<, ≤, =, ≥, >} and Q ⊆ S, let [Q]Col⊢n := {s ∈ Q |
Col(s) \triangleright n} be the set of states in Q with color \triangleright n. The parity objective is defined by
\[
\text{Parity}(Col) := \left( \bigvee_{i \in C} (GF[s]^{Col=2 \cdot i} \land FG[s]^{Col \leq 2 \cdot i}) \right),
\]
i.e., Parity(Col) is the set of infinite plays such that the largest color that occurs infinitely often along the play is even.

The Mostowski hierarchy [20] classifies parity objectives by restricting the range of the function Col to a set of colors \(C \subseteq \mathbb{N}\). We write \(C\text{-Parity}\) for such restricted parity objectives. In particular, Büchi objectives correspond to \(\{1,2\}\text{-Parity}\), and co-Büchi objectives correspond to \(\{0,1\}\text{-Parity}\). The objectives \(\{0,1,2\}\text{-Parity}\) and \(\{1,2,3\}\text{-Parity}\) are incomparable, but they both subsume (modulo renaming of colors) Büchi and co-Büchi objectives. Moreover, both \(\{0,1\}\text{-Parity}\) and \(\{1,2\}\text{-Parity}\) subsume the reachability objective Reach(T) (for MDPs with a sink T), by defining the color function so that \(Col(s) = 1 \implies s \not\in T\). Similarly, both \(\{0,1\}\text{-Parity}\) and \(\{1,2\}\text{-Parity}\) subsume Safety(T), by defining \(Col(s) = 1 \implies s \in T\).

### Optimal and \(\epsilon\)-Optimal Strategies

Given an objective \(\varphi\), the value of state \(s\) in an MDP \(M\), denoted by \(\text{val}_M(s)\), is the supremum of probability of achieving \(\varphi\), i.e.,
\[
\text{val}_M(s) := \sup_{\pi \in \mathcal{P}} \mathcal{P}_M(s,\pi)(\varphi).
\]
For \(\epsilon \geq 0\) and \(s \in S\), we say that a strategy \(\sigma\) is \(\epsilon\)-optimal if \(\mathcal{P}_M(s,\sigma)(\varphi) \geq \text{val}_M(s) - \epsilon\). A 0-optimal strategy is called optimal. An optimal strategy is almost-surely winning if \(\text{val}_M(s) = 1\). Unlike in finite-state MDPs, optimal strategies need not exist in countable MDPs, not even for reachability objectives in finitely branching MDPs. However, by the definition of the value, for all \(\epsilon > 0\), an \(\epsilon\)-optimal strategy exists.

For an objective \(\varphi\) and a \(\triangleright \in \{\geq, >\}\) and \(c \in [0,1]\), we define \([\varphi]^c\triangleright\) as the set of states \(s\) for which there exists a strategy \(\sigma\) with \(\mathcal{P}_M(s,\sigma)(\varphi) \triangleright c\). We call a state \(s\) almost-surely winning if \(s \in [\varphi]^{\geq 1}\), and we call \(s\) limit-surely winning if \(s \in [\varphi]^{\triangleright c}\) for every constant \(c < 1\) (which is iff \(\text{val}_M(s) = 1\)). On infinite arenas, limit-surely winning states are not necessarily almost-surely winning.

### III. OBJECTIVES THAT REQUIRE INFINITE MEMORY

In this section we consider those objectives in the Mostowski hierarchy where optimal (resp., \(\epsilon\)-optimal) strategies require infinite memory. In each such case we construct an MDP that witnesses this requirement. In these MDPs, all FR-strategies achieve the objective only with probability 0, while some HD-strategy achieves the objective almost-surely (resp., with arbitrarily high probability).

**Theorem 1.** Let \(\varphi = \{1,2,3\}\text{-Parity}\). There exists a finitely branching MDP \(M\) with initial state \(s_0\) such that
- for all FR-strategies \(\sigma\), we have \(\mathcal{P}_M(s_0,\sigma)(\varphi) = 0\),
- there exists an HD-strategy \(\sigma\) such that \(\mathcal{P}_M(s_0,\sigma)(\varphi) = 1\).

Hence, optimal (and even almost-surely winning) and \(\epsilon\)-optimal strategies require infinite memory for \(\{1,2,3\}\text{-Parity}\), even in finitely branching MDPs.

The MDP in Theorem 1 is depicted in Figure 2 (left), where \(Col(s_i) = 1\) and \(Col(r_j) = 2\) for all \(i \in \mathbb{N}\), and \(Col(t) = 3\). For every FR-strategy there is a uniform lower bound on the probability of visiting \(t\) between consecutive visits to \(s_0\). Hence, unless the strategy with positive probability eventually always stays in states \(s_i\) (and thus also loses the almost-sure parity objective), in the long-run, the probability of visiting \(t\) (with color three) tends to 1, and the parity condition is satisfied with probability 0. Although the player cannot win by any FR-strategy, we construct an HD-strategy \(\sigma\) such that \(\mathcal{P}_M(s_0,\sigma)(\varphi) = 1\). This strategy is such that upon the \(i\)-th visit to \(s_0\), the ladder \(s_0s_1s_2\cdots\) is traversed and the transition \(s_i \rightarrow r_j\) is chosen. Moving further along the ladder \(s_0s_1s_2\cdots\) decreases the probability of visiting \(t\) between the previous and successive visits to \(s_0\). Hence, the probability of visiting color three infinitely often is 0.

**Remark 1.** A strict subclass of finitely branching MDPs are 1-counter MDPs, where a finite-state MDP is augmented with an integer counter [6]. The MDP in Theorem 1 (plus some auxiliary states) is implementable by a 1-counter MDP.

**Remark 2.** The classical Rabin and Streett conditions can encode \(\{1,2,3\}\text{-Parity}\). Thus, optimal and \(\epsilon\)-optimal strategies for Rabin/Streett require infinite memory, even in finitely branching countable MDPs.

On finite MDPs, optimal strategies can be chosen MD for parity and Rabin objectives, but not for Streett objectives. Optimal strategies for Streett objectives can be chosen MR or FD [9].

**Proof.** For an infinite play \(\pi\), let \(\text{Inf}(\pi)\) be the set of states that \(\pi\) visits infinitely often. Let us recall the Rabin and Streett conditions.

Given a Rabin condition \(\{(E_1,F_1), (E_2,F_2), \cdots, (E_n,F_n)\}\) with \(n\) pairs (or \(n\) disjunctions), an infinite play \(\pi\) satisfies the Rabin condition if there exists a pair \((E_i,F_i)\) such that \(\text{Inf}(\pi) \cap E_i = \emptyset\) and \(\text{Inf}(\pi) \cap F_i \neq \emptyset\). The Rabin condition
\[
\{(\{s\}^{\text{Col}=3}, \{s\}^{\text{Col}=2})\}
\]
encodes \(\{1,2,3\}\text{-Parity}\), since all satisfying runs must visit states with color 2 infinitely often and states with color 3 only finitely often. Note that \(\{1,2,3\}\text{-Parity}\) is encoded in a Rabin condition with only one disjunction.

Given a Streett condition \(\{(E_1,F_1), (E_2,F_2), \cdots, (E_n,F_n)\}\) with \(n\) pairs (or \(n\) conjunctions), an infinite play \(\pi\) satisfies the Streett condition if \(\text{Inf}(\pi) \cap E_i = \emptyset\) implies \(\text{Inf}(\pi) \cap F_i = \emptyset\) for all pairs \((E_i,F_i)\). The Streett condition
\[
\{(\{s\}^{\text{Col}=2}, S), (\emptyset, \{s\}^{\text{Col}=3})\}
\]
encodes \(\{1,2,3\}\text{-Parity}\), since all satisfying runs must visit states with color 2 infinitely often and states with color 3 only finitely often.

Note that a conjunction of two Streett pairs are needed to encode \(\{1,2,3\}\text{-Parity}\). A single Streett pair \((\{X,Y\})\) means “infinitely often X or only finitely often Y”, which can be encoded as a \(\{0,1,2\}\text{-Parity}\) condition by assigning
color 2 to $X$ and color 1 to $Y$. Unlike for $\{1,2,3\}$-Parity, optimal strategies for $\{0,1,2\}$-Parity (and thus also for a single Streett pair) can be chosen MD in finitely branching MDPs (Theorem 16).

It was known that quantitative Büchi objectives require infinite memory [18], [2]. For the sake of completeness, we present an example MDP for Proposition 2 in Figure 2 (right).

**Proposition 2 ([18]).** Let $\varphi = \{1,2\}$-Parity be the Büchi objective. There exists a finitely branching MDP $M$ with initial state $s_0$ such that
- for all FR-strategies $\sigma$, we have $P_{M,s_0,\sigma}(\varphi) = 0$,
- for every $c \in [0,1)$, there exists an HD-strategy $\sigma$ such that $P_{M,s_0,\sigma}(\varphi) \geq c$.

Hence, $c$-optimal strategies for Büchi objectives require infinite memory.

**Theorem 3.** Let $\varphi = \text{Safety}(T)$. There exists an infinitely branching MDP $M$ with initial state $s$ such that
- for all FR-strategies $\sigma$, we have $P_{M,s,\sigma}(\varphi) = 0$,
- for every $c \in [0,1)$, there exists an HD-strategy $\sigma$ such that $P_{M,s,\sigma}(\varphi) \geq c$.

Hence, $c$-optimal strategies for safety require infinite memory.

The MDP in Theorem 3, depicted in Figure 3 (left), was first introduced in [19]. Since our notion of finite-memory strategies allows for randomized memory updates (in contrast to [19]), our proof is somewhat more general. The target is $T = \{t\}$. For every FR-strategy there is a uniform lower bound on the probability of reaching $t$ between consecutive visits to $s_0$. Since $t$ is absorbing, it will be reached with probability 1. Thus every FR-strategy satisfies the safety objective with probability 0. However, for all $n \in \mathbb{N}$, we construct an HD-strategy $\sigma_n$ such that $P_{M,s,\sigma_n}(\text{Safety}\{\{t\}\}) \geq 1 - \frac{1}{2^n}$. This strategy is such that upon the $n$th visit to $s$, the transition $s \rightarrow r_{t+n}$ is chosen. Hence, the probability of visiting $t$ between two successive visits to $s$ decreases. A more detailed analysis shows that the probability of ever visiting $t$ is bounded by $\frac{1}{2^n}$.

**Theorem 4.** Let $\varphi = \{0,1\}$-Parity be the co-Büchi objective. There exists an infinitely branching MDP $M$ with initial state $s$ such that
- for all FR-strategies $\sigma$, we have $P_{M,s,\sigma}(\varphi) = 0$,
- there exists an HD-strategy $\sigma$ such that $P_{M,s,\sigma}(\varphi) = 1$.

Hence, optimal (and even almost-surely winning) strategies and $c$-optimal strategies for co-Büchi require infinite memory.

The MDP in Theorem 4 is depicted in Figure 3 (right). By a similar argument as in Theorem 3, every FR-strategy achieves co-Büchi with probability 0. However, the HD-strategy $\sigma$ that chooses the transition $s \rightarrow r_i$ upon the $i$th visit to $s$ is such that $P_{M,s,\sigma}(\varphi) = 1$.

**IV. FROM ALMOST-SURE WINNING TO OPTIMAL STRATEGIES**

In this section we prove Theorem 5. It says that, for certain objectives, if almost-surely winning strategies (where they exist) can be chosen MD, then optimal strategies (where they exist) can also be chosen MD.

We call a class $C$ of MDPs downward-closed if every MDP whose transition relation is a subset of the transition relation of some MDP in $C$ is also in $C$. The class of finitely branching MDPs is downward-closed, and so is the class of MDPs with a fixed sink $T$.

We call an objective $\varphi$ prefix-independent in $C$ (where $C$ is a class of MDPs) if for all $w_1, w_2 \in S^*$ and all $w \in S^\omega$ such that $w_1 w$ and $w_2 w$ are infinite plays in an MDP in $C$ we have $w_1 w \in [\varphi] \iff w_2 w \in [\varphi]$. Parity objectives are prefix-independent in the class of all MDPs. Both objectives Reach($T$) and Safety($T$) are prefix-independent in the class of MDPs with sink $T$.

The following theorem provides, under certain conditions, an optimal MD-strategy for all states that have an optimal strategy. In fact, a single MD-strategy is optimal for all states that have an optimal strategy:

**Theorem 5.** Let $\varphi$ be an objective that is prefix-independent in a downward-closed class $C$ of MDPs. Suppose that for
any \( M = (S, S_0, S_\omega, \rightarrow, P) \in C \) and any \( s \in S \) and any strategy \( \sigma \) with \( P_{M,s,\sigma}(\varphi) = 1 \) there exists an MD-strategy \( \sigma' \) with \( P_{M,s,\sigma'}(\varphi) = 1 \).

Under this condition, for each \( M \in C \) there is an MD-strategy \( \sigma' \) such that for all \( s \in S \):

\[
(\exists \sigma \in \Sigma. P_{M,s,\sigma}(\varphi) = \text{val}_M(s)) \implies P_{M,s,\sigma'}(\varphi) = \text{val}_M(s)
\]

The remainder of the section is devoted to the proof of Theorem 5.

For prefix-independent winning conditions, whenever an optimal strategy visits some state, it achieves the value of this state; see Lemma 20 in the appendix. We use this to show that the MDP constructed! the following lemma provides, under certain conditions, a uniform almost-surely winning MD-strategy, i.e., one that works for all initial states at the same time:

**Lemma 7.** Let \( M = (S, S_0, S_\omega, \rightarrow, P) \) be an MDP. Let \( \varphi \) be an objective that is prefix-independent in \( \{ M \} \). Suppose that for any \( s \in S \) and any strategy \( \sigma \) with \( P_{M,s,\sigma}(\varphi) = 1 \) there exists an MD-strategy \( \sigma' \) with \( P_{M,s,\sigma'}(\varphi) = 1 \). Then there is an MD-strategy \( \sigma' \) such that for all \( s \in S \):

\[
(\exists \sigma \in \Sigma. P_{M,s,\sigma}(\varphi) = 1) \implies P_{M,s,\sigma'}(\varphi) = 1
\]

**Proof.** We can assume that all states are almost-surely winning, since in order to achieve an almost-sure winning objective, the player must forever remain in almost-surely winning states. So we need to define an MD-strategy \( \sigma' \) so that for all \( s \in S \) we have \( P_{M,s,\sigma'}(\varphi) = 1 \).

Fix an arbitrary state \( s_1 \in S \). By assumption there is an MD-strategy \( \sigma_1 \) with \( P_{M,s_1,\sigma_1}(\varphi) = 1 \). Let \( U_1 \subseteq S \) be the set of states that occur in plays that both start from \( s_1 \) and are induced by \( \sigma_1 \). We have \( P_{M,s_1,\sigma_1}(\varphi^{s_1} \cap U_1) = 1 \). In fact, for any \( s \in U_1 \) and any strategy \( \sigma \) that agrees with \( \sigma_1 \) on \( U_1 \) we have \( P_{M,s,\sigma}(\varphi^{s_1} \cap U_1) = 1 \).

If \( U_1 = S \) we are done. Otherwise, consider the MDP \( M_1 \) obtained from \( M \) by fixing \( \sigma_1 \) on \( U_1 \) (i.e., in \( M_1 \) we can view the states in \( U_1 \) as random states). We argue that, in \( M_1 \), for any state \( s \) there is an MD-strategy \( \sigma'_1 \) with \( P_{M_1,s,\sigma'_1}(\varphi) = 1 \). Indeed, let \( s \in S \) be any state. Recall that there is an MD-strategy \( \sigma' \) with \( P_{M,s,\sigma'}(\varphi) = 1 \). Let \( \sigma'_1 \) be the MD-strategy obtained by restricting \( \sigma' \) to the non-\( U_1 \) states (recall that the \( U_1 \) states are random states in \( M_1 \)). This strategy \( \sigma'_1 \) almost surely generates a play that either satisfies \( \varphi \) without ever entering \( U_1 \) or at some point enters \( U_1 \). In the latter case, \( \varphi \) is satisfied almost surely: this follows from prefix-independence and the fact that \( \sigma'_1 \) agrees with \( \sigma_1 \) on \( U_1 \). We conclude that \( P_{M_1,s,\sigma'_1}(\varphi) = 1 \).

Let \( s_2 \in S \setminus U_1 \). We repeat the argument from above, with \( s_2 \) instead of \( s_1 \), and with \( M_1 \) instead of \( M \). This yields an MD-strategy \( \sigma_2 \) and a set \( U_2 \ni s_2 \) with \( P_{M_1,s_2,\sigma_2}(\varphi^{s_2} \cap U_2) = 1 \). In fact, for any \( s \in U_2 \) and any strategy \( \sigma \) that agrees with \( \sigma_2 \) on \( U_2 \) and with \( \sigma_1 \) on \( U_1 \) we have \( P_{M,s,\sigma}(\varphi^{s_2} \cap U_2) = 1 \).

If \( U_1 \cup U_2 = S \) we are done. Otherwise we continue in the same manner, and so forth. Since \( S \) is countable, we can pick
\[ P_{M,s,\sigma}(\varphi) = \sup_{\sigma' \in \Sigma} P_{M,s,\sigma'}(\varphi) = val_M(s) \]
Definition 1. Let $\mathcal{M} = (S, S_0, S_\Omega, \rightarrow, P)$ be a finitely branching MDP, $\text{Col} : S \rightarrow \mathbb{N}$ a color function, $\varphi = \text{Safety}(\{S\}^{\text{Col} \neq 0})$, $\sigma_{\text{opt-av}}$ the strategy from Proposition 13 and $\tau \in [0, 1]$. We define

$$\text{Safe}_{\mathcal{M}}(\tau) := \{ s \in S \mid \mathcal{P}_{M, s, \sigma_{\text{opt-av}}} (\varphi) \geq \tau \},$$

i.e., $\text{Safe}_{\mathcal{M}}(\tau)$ is the set of states from which the player can remain within color-0 states forever with probability $\geq \tau$. We drop the subscript $\mathcal{M}$ when the MDP $\mathcal{M}$ is understood.

Loosely speaking, the following lemma gives a lower bound on the probability that, starting from a “safe” state, “unsafe” states are forever avoided by $\sigma_{\text{opt-av}}$:

Lemma 14. Let $\mathcal{M} = (S, S_0, S_\Omega, \rightarrow, P)$ be a finitely branching MDP, $\text{Col} : S \rightarrow \mathbb{N}$ a color function and $\sigma_{\text{opt-av}}$ the strategy from Proposition 13. Let $0 < \tau_1 \leq \tau_2 \leq 1$, and $s \in \text{Safe}(\tau_1)$. Then $\mathcal{P}_{M, s, \sigma_{\text{opt-av}}} (G\text{Safe}(\tau_1)) \geq \frac{\tau_2 - \tau_1}{1 - \tau_1}$.

Proof. We compute probabilities conditioned under the event $G\text{Safe}(\tau_1)$. Since $\text{Safe}(\tau_1) \subseteq \{S\}^{\text{Col} = 0}$, we have $\mathcal{P}_{M, s, \sigma_{\text{opt-av}}} (G\{S\}^{\text{Col} = 0} \mid G\text{Safe}(\tau_1)) = 1$. From the definition of $\text{Safe}(\tau_1)$ and the Markov property we get $\mathcal{P}_{M, s, \sigma_{\text{opt-av}}} (G\{S\}^{\text{Col} = 0} \mid -G\text{Safe}(\tau_1)) \leq \tau_1$. Applying the law of total probability and writing $x$ for $\mathcal{P}_{M, s, \sigma_{\text{opt-av}}} (G\text{Safe}(\tau_1))$ we obtain:

$$\tau_2 \leq \mathcal{P}_{M, s, \sigma_{\text{opt-av}}} (G\{S\}^{\text{Col} = 0}) \quad \text{Def. 1}$$

$$= \mathcal{P}_{M, s, \sigma_{\text{opt-av}}} (G\{S\}^{\text{Col} = 0} \mid G\text{Safe}(\tau_1)) \cdot x + \mathcal{P}_{M, s, \sigma_{\text{opt-av}}} (G\{S\}^{\text{Col} = 0} \mid -G\text{Safe}(\tau_1)) \cdot (1 - x)$$

$$\leq x + \tau_1 \cdot (1 - x)$$

It follows $x \geq \frac{\tau_2 - \tau_1}{1 - \tau_1}$.

The following lemma states for all $\tau < 1$ that eventually remaining in color-0 states but outside $\text{Safe}(\tau)$ has probability zero.

Lemma 15. Let $\mathcal{M} = (S, S_0, S_\Omega, \rightarrow, P)$ be a finitely branching MDP, and $\text{Col} : S \rightarrow \mathbb{N}$ a color function. Let $s$ be a state, and $\sigma$ a strategy, and $\tau < 1$. Then $\mathcal{P}_{M, s, \sigma}(FG-\text{Safe}(\tau) \land FG\{S\}^{\text{Col} = 0}) = 0$.

B. Optimal MD-strategies for $\{0, 1, 2\}$-Parity

Theorem 16. Let $\mathcal{M}$ be a finitely branching MDP, $\text{Col} : S \rightarrow \{0, 1, 2\}$, and $\varphi = \text{Parity}(\text{Col})$. Then there exists an MD-strategy $\sigma'$ that is optimal for all states that have an optimal strategy:

$$\exists \sigma \in \Sigma \mathcal{P}_{M, s, \sigma}(\varphi) = \text{val}_\mathcal{M}(s) \implies \mathcal{P}_{M, s, \sigma'}(\varphi) = \text{val}_\mathcal{M}(s)$$

By appealing to Theorem 5 it suffices to show:

Proposition 17. Let $\mathcal{M} = (S, S_0, S_\Omega, \rightarrow, P)$ be a finitely branching MDP, and $s_0 \in S$, and $\sigma$ a strategy, and $\text{Col} : S \rightarrow \{0, 1, 2\}$, and $\varphi = \text{Parity}(\text{Col})$. Suppose $\mathcal{P}_{M, s_0, \sigma}(\varphi) = 1$. Then there is an MD-strategy $\sigma'$ with $\mathcal{P}_{M, s_0, \sigma'}(\varphi) = 1$.

The following simple lemma provides a scheme for proving almost-sure properties.

Lemma 18. Let $\mathcal{P}$ be a probability measure over the sample space $\Omega$. Let $(\mathcal{R}_i)_{i \in I}$ be a countable partition of $\Omega$ in measurable events. Let $E \subseteq \Omega$ be a measurable event. Suppose $\mathcal{P}(\mathcal{R}_i \cap E) = \mathcal{P}(\mathcal{R}_i)$ holds for all $i \in I$. Then $\mathcal{P}(E) = 1$.

We are ready to prove Proposition 17.

Proof of Proposition 17. To achieve an almost-sure winning objective, the player must forever remain in states from which the objective can be achieved almost surely. So we can assume without loss of generality that all states are almost-sure winning, i.e., for all $s \in S$ we have $\mathcal{P}_{M, s, \sigma}(\varphi) = 1$ for some $\sigma$.

We will define an MD-strategy $\sigma'$ with $\mathcal{P}_{M, s, \sigma'}(\varphi) = 1$ for all $s \in S$. We first define the MD-strategy $\sigma'$ partially for the states in $\text{Safe}_{\mathcal{M}}(\frac{1}{2})$ and then extend the definition of $\sigma'$ to all states. For the states in $\text{Safe}_{\mathcal{M}}(\frac{1}{2})$ define $\sigma' := \sigma_{\text{opt-av}}$ as in Proposition 13, see Figure 4. Let $M'$ be the MDP obtained from $M$ by restricting the transition relation as prescribed by the partial MD-strategy $\sigma'$.

For any $\tau \in [0, 1]$, we have $\text{Safe}_{\mathcal{M}}(\tau) = \text{Safe}_{M'}(\tau)$. Indeed, since $M'$ restricts the options of the player, we have $\text{Safe}_{\mathcal{M}}(\tau) \supseteq \text{Safe}_{M'}(\tau)$. Conversely, let $s \in \text{Safe}_{M'}(\tau)$. The strategy $\sigma_{\text{opt-av}}$ from Proposition 13 achieves $\mathcal{P}_{M, s, \sigma_{\text{opt-av}}}(G\{S\}^{\text{Col} = 0}) \geq \tau$. Since $\sigma_{\text{opt-av}}$ can be applied in $M'$, and results in the same Markov chain as applying it in $M$, we conclude $s \in \text{Safe}_{\mathcal{M}}(\tau)$. This justifies to write $\text{Safe}_{\mathcal{M}}(\tau) = \text{Safe}_{M'}(\tau)$ in the remainder of the proof.

Next we show that, also in $M'$, for all states $s \in S$ there exists a strategy $\sigma_1$ with $\mathcal{P}_{M', s, \sigma_1}(\varphi) = 1$. This strategy $\sigma_1$ is defined as follows. First play according to a strategy $\sigma$ from the statement of the theorem. If and when the play visits $\text{Safe}_{M'}(\frac{1}{2})$, switch to the MD-strategy $\sigma_{\text{opt-av}}$ from Proposition 13. If and when the play then visits $\{S\}^{\text{Col} \neq 0}$, switch back to a strategy $\sigma$ from the statement of the theorem, and so forth.

We show that $\sigma_1$ achieves $\mathcal{P}_{M', s, \sigma_1}(\varphi) = 1$. To this end we will use Lemma 18. We partition the runs of $sS^n$ into three events $\mathcal{R}_0, \mathcal{R}_1, \mathcal{R}_2$ as follows:

- $\mathcal{R}_0$ contains the runs where $\sigma_1$ switches between $\sigma_{\text{opt-av}}$ and $\sigma$ infinitely often.
- $\mathcal{R}_1$ contains the runs where $\sigma_1$ eventually only plays according to $\sigma_{\text{opt-av}}$.
- $\mathcal{R}_2$ contains the runs where $\sigma_1$ eventually only plays according to $\sigma$.

Each time $\sigma_1$ switches to $\sigma_{\text{opt-av}}$, there is, by Proposition 13, a probability of at least $\frac{1}{3}$ of never visiting a color-$\{1, 2\}$ state again and thus of never again switching to $\sigma$. It follows that $\mathcal{P}_{M', s, \sigma_1}(\mathcal{R}_0) = 0$. By the definition of $\sigma_{\text{opt-av}}$ we have $\mathcal{R}_1 \subseteq \{FG\{S\}^{\text{Col} = 0} \subseteq \{\varphi\}$, and hence $\mathcal{P}_{M', s, \sigma_1}(\mathcal{R}_1 \cap \{\varphi\}) = \mathcal{P}_{M', s, \sigma_1}(\mathcal{R}_1)$. Since $\mathcal{P}_{M, s, \sigma}(\varphi) = 1$ and $\varphi$ is prefix-independent, we have $\mathcal{P}_{M', s, \sigma_1}(\mathcal{R}_2 \cap \{\varphi\}) = \mathcal{P}_{M', s, \sigma_1}(\mathcal{R}_2)$. Using Lemma 18, we obtain $\mathcal{P}_{M', s, \sigma_1}(\varphi) = 1$. 

8
Next we show that for all $s \in S$ the strategy $\sigma_1$ defined above achieves $P_{M',s,\sigma_1}(FSafe(\frac{3}{3}) \lor [S]_{Col=2} = 1$. To this end we will use Lemma 18 again. We partition the runs of $SS^o$ into three events $\mathcal{R}_1', \mathcal{R}_2', \mathcal{R}_0'$ as follows:

- $\mathcal{R}_1' = \{FG[S]_{Col=0} \}$
- $\mathcal{R}_2' = \{GF[S]_{Col=2} \}$
- $\mathcal{R}_0' = sS^o \setminus \{\varphi\}$

We have previously shown that $P_{M',s,\sigma_1}(\varphi) = 1$, hence $P_{M',s,\sigma_1}(\mathcal{R}_0') = 0$. By Lemma 15, almost all runs in $\mathcal{R}_1'$ satisfy $GF Safe(\frac{3}{3})$. Since $\{GF Safe(\frac{3}{3})\} \subseteq \{Safe(\frac{3}{3})\}$, we have $P_{M',s,\sigma_1}(\mathcal{R}_1' \cap \{F Safe(\frac{3}{3}) \lor [S]_{Col=2} = 1\}) = P_{M',s,\sigma_1}(\mathcal{R}_1')$. Since $\mathcal{R}_2' = \{FG[S]_{Col=2}\}$, we also have $P_{M',s,\sigma_1}(\mathcal{R}_2' \cap \{F Safe(\frac{3}{3}) \lor [S]_{Col=2} = 1\}) = P_{M',s,\sigma_1}(\mathcal{R}_2')$. Using Lemma 18 we obtain $P_{M',s,\sigma_1}(F Safe(\frac{3}{3}) \lor [S]_{Col=2} = 1) = 1$.

Writing $T = Safe(\frac{3}{3}) \cup [S]_{Col=2}$ we have just shown that for all $s \in S$ there is a strategy $\sigma_1$ with $P_{M,s,\sigma_1}(FT) = 1$. By Lemma 7 there is an MD-strategy $\sigma$ for $M'$ with $P_{M,s,\sigma}(FT) = 1$ for all $s \in S$. We extend the (so far partially defined) strategy $\sigma'$ by $\sigma$. Thus we obtain a (fully defined) strategy $\sigma$ for $M$ such that for all $s \in S$ we have $P_{M,s,\sigma_1}(FT) = 1$.

It remains to show that for all $s \in S$ we have $P_{M,s,\sigma_1}(\varphi) = 1$. To this end we will use Lemma 18 again. We partition the runs of $sS^o$ in two events $\mathcal{R}_1'', \mathcal{R}_2''$.

- $\mathcal{R}_1'' = \{FG Safe(\frac{3}{3})\}^*$, i.e., $\mathcal{R}_1''$ contains the runs that visit $Safe(\frac{3}{3})$ infinitely often.
- $\mathcal{R}_2'' = \{FG Safe(\frac{3}{3})\}^*$, i.e., $\mathcal{R}_2''$ contains the runs that start from some point on never visit $Safe(\frac{3}{3})$.

Every time a run enters $Safe(\frac{3}{3})$, by Lemma 14, the probability is at least $\frac{1}{2}$ that the run remains in $Safe(\frac{3}{3})$ forever. It follows that almost all runs in $\mathcal{R}_1''$ eventually remain in $Safe(\frac{3}{3})$ forever, i.e., $P_{M,s,\sigma_1}(\mathcal{R}_1'' \cap \{FG Safe(\frac{3}{3})\} = P_{M,s,\sigma_1}(\mathcal{R}_1')$. Since $Safe(\frac{1}{3}) \subseteq [S]_{Col=0}$, we have $\{FG Safe(\frac{3}{3})\} \subseteq [S]_{Col=0} \subseteq [\varphi]$. Hence also $P_{M,s,\sigma_1}(\mathcal{R}_1'' \cap [\varphi]) = P_{M,s,\sigma_1}(\mathcal{R}_1')$.

We have previously shown that $P_{M,s,\sigma_1}(FT) = 1$ holds for all $s \in S$. Hence also $P_{M,s,\sigma}(GFT) = 1$ holds for all $s \in S$. In particular, almost all runs in $\mathcal{R}_0''$ satisfy $GF$. By comparing the definitions of $\mathcal{R}_0''$ and $T$ we see that almost all runs in $\mathcal{R}_0''$ even satisfy $GF[S]_{Col=2}$. Since $\{GF[S]_{Col=2}\} \subseteq [\varphi]$, we obtain $P_{M,s,\sigma}(\mathcal{R}_0'' \cap [\varphi]) = P_{M,s,\sigma_1}(\mathcal{R}_0'')$.

A final application of Lemma 18 yields $P_{M,s,\sigma_1}(\varphi) = 1$ for all $s \in S$.

C. $\epsilon$-Optimal MD-strategies for Co-Büchi

Theorem 19. Let $M = (S, S_0, S, \rightarrow, P)$ be a finitely branching MDP, $Col : S \rightarrow \{0,1\}$, and $\varphi = Parity(Col)$ be the co-Büchi objective. Then there exist uniform $\epsilon$-optimal MD-strategies. I.e., for every $\epsilon > 0$ there is an MD-strategy $\sigma_\epsilon$ with $P_{M,s,\sigma_\epsilon}(\varphi) \geq \text{val}_{M,s_0}(s_0) - \epsilon$ for every $s_0 \in S$.

Proof. Let $\epsilon_1 > 0$ be a suitably small number (to be determined later), $\tau_1 := 1 - \epsilon_1$ and $Safe_{M}(\tau_1)$ defined as in Definition 1. Let $\sigma_{opt-av}$ be the MD-strategy from Proposition 13. From $M$ we obtain a modified MDP $M'$ by fixing all player choices from states in $Safe_{M}(\tau_1)$ according to $\sigma_{opt-av}$.

We show that $\text{val}_{M,s_0}(s_0) \geq \delta$. By definition of the value $\text{val}_{M,s_0}(s_0)$, for every $\delta > 0$ there exists a strategy $\delta$ in $M$ s.t. $P_{M,s_0,\sigma_1}(\varphi) \geq \text{val}_{M,s_0}(s_0) - \delta$. We define a strategy $\sigma_{\epsilon}'$ in $M'$ from state $s_0$ as follows. First play like $\delta$. If and when a state in $Safe_{M}(\tau_1)$ is reached, play like $\sigma_{opt-av}$. This is possible, since no moves from states outside $Safe_{M}(\tau_1)$ have been fixed in $M'$, and all moves from states inside $Safe_{M}(\tau_1)$ have been fixed according to $\sigma_{opt-av}$. Then we have:

$$P_{M,s_0,\sigma_{\epsilon}'}(\varphi)$$

$$= P_{M,s_0,\sigma_{\epsilon}'}(\varphi)$$

$$- P_{M,s_0,\sigma_{\epsilon}'}(F Safe_{M}(\tau_1)) \cdot P_{M,s_0,\sigma_{\epsilon}'}(\varphi | F Safe_{M}(\tau_1))$$

$$+ P_{M,s_0,\sigma_{\epsilon}'}(F Safe_{M}(\tau_1)) \cdot P_{M,s_0,\sigma_{\epsilon}'}(\varphi | F Safe_{M}(\tau_1))$$

$$\geq \text{val}_{M,s_0}(s_0) - \delta - P_{M,s_0,\sigma_{\epsilon}'}(F Safe_{M}(\tau_1)) \cdot (1 - \tau_1)$$

$$\geq \text{val}_{M,s_0}(s_0) - \delta - \epsilon_1$$
Since this holds for every $\delta > 0$ we obtain $\text{val}_M(s_0) \geq \text{val}_M(s_0) - \epsilon_1$.

Now let $\tau_2 := 1 - \epsilon_1/k$ for a suitably large $k \geq 1$ (to be determined later) and $\text{Safe}_M(\tau_2)$ be defined as in Definition 1. In particular, $\text{Safe}_M(\tau_2) = \text{Safe}_M(\tau_2)$ (by the same argument as in the proof of Proposition 17).

By definition of the value, for every $\epsilon_2 > 0$ there exists a strategy $\sigma_{\epsilon_2}$ in $M'$ with $\text{P}_{M',s_0,\sigma_{\epsilon_2}}(\varphi) \geq \text{val}_M(s_0) - \epsilon_2$. Moreover, by Lemma 15 and $\tau_2 < 1$, $\text{P}_{M',s_0,\sigma}(\text{Safe}_M(\tau_2)) \geq \text{P}_{M',s_0,\sigma}(\varphi)$ for every strategy $\sigma$ and thus in particular for $\sigma_{\epsilon_2}$. Therefore, $\text{P}_{M',s_0,\sigma_{\epsilon_2}}(\text{Safe}_M(\tau_2)) \geq \text{val}_M(s_0) - \epsilon_2$. By Theorem 9, for every $\epsilon_3 > 0$ there exists an MD-strategy $\sigma'$ in $M'$ with $\text{P}_{M',s_0,\sigma'}(\text{Safe}_M(\tau_2)) \geq \text{val}_M(s_0) - \epsilon_2 - \epsilon_3$. In particular, $\sigma'$ must coincide with $\sigma_{\text{opt-av}}$ at all states in $\text{Safe}_M(\tau_1)$, since in $M'$ these choices are already fixed.

We obtain the MD-strategy $\sigma_\epsilon$ in $M$ by combining the MD-strategies $\sigma'$ and $\sigma_{\text{opt-av}}$. The strategy $\sigma_\epsilon$ plays like $\sigma'$ at all states outside $\text{Safe}_M(\tau_1)$ and like $\sigma_{\text{opt-av}}$ at all states inside $\text{Safe}_M(\tau_1)$.

In order to show that $\sigma_\epsilon$ has the required property $\text{P}_{M,s_0,\sigma_\epsilon}(\varphi) \geq \text{val}_M(s_0) - \epsilon$, we first estimate the probability that a play according to $\sigma_\epsilon$ will never leave the set $\text{Safe}_M(\tau_1)$ after having visited a state in $\text{Safe}_M(\tau_2)$.

Let $s \in \text{Safe}_M(\tau_2)$. Then, by Lemma 14,

$$\text{P}_{M,s,\sigma_{\text{opt-av}}}(\text{GSafe}_M(\tau_1)) \geq \frac{\tau_2 - \tau_1}{1 - \tau_1} = \frac{(1 - \epsilon_1/k) - (1 - \epsilon_1)}{\epsilon_1} = 1 - \frac{1}{k}.$$ 

In particular we also have $\text{P}_{M,\sigma_\epsilon}(\text{GSafe}_M(\tau_1)) \geq 1 - \frac{1}{k}$, since $\sigma_\epsilon$ coincides with $\sigma_{\text{opt-av}}$ inside the set $\text{Safe}_M(\tau_1)$. Finally we obtain:

$$\text{P}_{M,s_0,\sigma_\epsilon}(\varphi) \geq \text{P}_{M,s_0,\sigma}(\text{FSafe}_M(\tau_2)) \cdot \text{P}_{M,s_0,\sigma}(\text{GFSafe}_M(\tau_1) \mid \text{FSafe}_M(\tau_2)) \geq \text{P}_{M,s_0,\sigma}(\varphi) \cdot (1 - 1/k) \geq (\text{val}_M(s_0) - \epsilon_2 - \epsilon_3) \cdot (1 - 1/k) \geq (\text{val}_M(s_0) - \epsilon_2 - \epsilon_3) \cdot (1 - 1/k)$$

This holds for every $\epsilon_1, \epsilon_2, \epsilon_3 > 0$ and every $k \geq 1$, and moreover $\text{val}_M(s_0) \leq 1$. Thus we can set $\epsilon_1 = \epsilon_2 = \epsilon_3 := \epsilon/6$ and $k := \frac{3}{\epsilon}$ and obtain $\text{P}_{M,s_0,\sigma}(\varphi) \geq \text{val}_M(s_0) - \epsilon$ for every $s_0 \in S$ as required.

VII. DISCUSSION

Our results on the memory requirements of ($\epsilon$)-optimal strategies (Figure 1) directly imply how much memory is needed to win quantitative objectives of type $[\varphi]^{>\text{opt}}$ (considered, e.g., in [7]). For $c < 1$ the assumed winning strategy might have to be an $\epsilon$-optimal one, since optimal strategies do not always exist. Thus MD-strategies are only sufficient for reachability objectives in countable MDPs (resp., for {0,1}-Parity, safety and reachability objectives in finitely branching MDPs). In the special case of $[\varphi]^{>\text{opt}}$ objectives (i.e., winning almost-surely), the winning strategy (assuming it exists) must be optimal. Thus MD-strategies are only sufficient for safety, reachability and {1,2}-Parity in countable MDPs (resp., for all objectives subsumed by {0,1,2}-Parity in finitely branching MDPs).

In this paper we have studied countable MDPs. Not all our results carry over to uncountable MDPs. The first issue is measurability. The probabilities are only well-defined if the strategies are measurable functions, which might not exist without further conditions on the MDP; cf. Section 2.3 in [23]. Another issue is that strategies cannot generally be chosen uniform, i.e., independent of the initial state. E.g., in countable MDPs $\epsilon$-optimal strategies for reachability can be chosen uniform MD (Theorem 9), but this does not carry over to uncountable MDPs (Thm. A in [21]). However, optimal strategies for reachability, if they exist, can be chosen uniform MD (Proposition B in [21]).

Acknowledgements. This work was partially supported by the EPSRC through grants EP/M027287/1, EP/M027651/1, EP/P020909/1 and EP/M003795/1 and by St. John’s College, Oxford.

REFERENCES


A. Proofs of Section III

We recall two results that are used throughout the proofs in this section:

a) Strong fairness of probabilistic choices in Markov chains.: Given a Markov chain, let $P_q(\cdot)$ denotes the probability of events starting in state $q$ of the chain. Let $p, q$ be two states and $\pi$ a finite path starting in $p$ with strictly positive probability. Strong fairness of probabilistic choices states that $P_q(GF\pi) = P_q(GFp)$. Intuitively, it means that under the condition that state $p$ is visited infinitely often, any finite path starting in $p$ will be taken infinitely often, almost surely [3, Theorem 10.25].

b) The Borel-Cantelli lemma.: Suppose that $(E_n)_{n\in\mathbb{N}}$ is a sequence of events in a probability space. Denote by $E^\infty$ the event

$$\bigcap_{k=1}^{\infty} \bigcup_{n=k}^{\infty} E_n,$$

that intuitively is the event "$E_n$ occurs for infinitely many $n$". The Borel-Cantelli lemma asserts that if $\sum_{n=1}^{\infty} P(E_n) < \infty$ then $P(E^\infty) = 0$. Informally speaking, if the sum of probabilities of the events $E_n$ is bounded then the probability that infinitely many of them occur is zero [5].

**Theorem 1.** Let $\varphi = \{1,2,3\}$-Parity. There exists a finitely branching MDP $M$ with initial state $s_0$ such that

- for all FR-strategies $\sigma$, we have $P_{\mathcal{M},s_0,\sigma}(\varphi) = 0$,
- there exists an HD-strategy $\sigma$ such that $P_{\mathcal{M},s_0,\sigma}(\varphi) = 1$.

Hence, optimal (and even almost-surely winning) and $\epsilon$-optimal strategies require infinite memory for $\{1,2,3\}$-Parity, even in finitely branching MDPs.

**Proof.** Consider the MDP $M$ shown (on the left side) in Figure 2 where $S_\circ = \{r_i\}_{i\in\mathbb{N}}$ and $S_D = \{t\} \cup \{s_i\}_{i\in\mathbb{N}}$. For all $i \in \mathbb{N}$, there are transitions $s_i \rightarrow r_i$ and $s_i \rightarrow s_{i+1}$ in $s_i$ states, whereas $P(r_i,t) = 2^{-i}$ and $P(r_i,s_0) = 1 - 2^{-i}$ in random states.

Let $\sigma$ be an arbitrary FR-strategy. We prove that $P_{\mathcal{M},s_0,\sigma}(\varphi) = 0$. By definition there is a transducer $T$ with finite memory $M$ and initial mode $m_0$ such that $\sigma_T = \sigma$. Let $\mathcal{M}^T$ be the Markov chain obtained by the product of the MDP $\mathcal{M}$ and the transducer $T$. The set of states in $\mathcal{M}^T$ is $M \times S$. We define a coloring function for $\mathcal{M}^T$ such that it ignores the memory mode and assigns to $(m,s)$ the same color as state $s$ in $\mathcal{M}$. By a slight abuse of notation, we use the same notation $Col$ for the coloring functions of both $\mathcal{M}$ and $\mathcal{M}^T$. We also denote by $P_{\mathcal{M}^T,q}(\mathcal{R})$ the probability of a measurable set $\mathcal{R}$ of infinite paths (i.e., infinite plays), starting in the state $q$ of $\mathcal{M}^T$.

We prove that $P_{\mathcal{M},s_0,\sigma}(\varphi^T) = 0$ where $\varphi^T = GF[M \times S]^{Col=2} \land GF[M \times S]^{Col\neq3}$. Equivalently, we show that $P_{\mathcal{M}^T,(m_0,s_0)}(\varphi^T) = 0$. We proceed in three steps: we will show, using strong fairness of probabilistic choices in Markov chains, that for all modes $m \in M$,

$$P_{\mathcal{M}^T,(m_0,s_0)}(GF(m,s_0) \land GF[M \times S]^{Col\neq3}) = 0. \quad (1)$$

Moreover,

$$\|GF[M \times S]^{Col=2}\| \subseteq \|GF(M \times \{s_0\})\| \quad (2)$$

Since the memory of strategy $\sigma$ is finite ($|M| < \infty$), we will show that

$$\|GF(M \times \{s_0\})\| = \left\| \bigvee_{m \in M} GF(m,s_0) \right\| \quad (3)$$

Using (1)-(3) we complete the proof as follows:

$$P_{\mathcal{M}^T,(m_0,s_0)}(GF[M \times S]^{Col=2} \land GF[M \times S]^{Col\neq3}) \leq P_{\mathcal{M}^T,(m_0,s_0)}\left( \bigvee_{m \in M} (GF(m,s_0) \land GF[M \times S]^{Col\neq3}) \right) \quad \text{by (2), (3)}$$

$$\leq \sum_{m \in M} P_{\mathcal{M}^T,(m_0,s_0)}(GF(m,s_0) \land GF[M \times S]^{Col\neq3}) \quad \text{union bound}$$

$$= 0. \quad \text{by (1)}$$

As a result $P_{\mathcal{M},s_0,\sigma}(\varphi) = 0$. Below, we prove (1), (2) and (3).

We first highlight two properties of the Markov chain $\mathcal{M}^T$. Consider the MDP $\mathcal{M}$: (i) the only states with color 2 are random states $r_i \in S_\circ$, wherein the only successors are $s_0$ and $t$. Hence, from all states $(m,r_i) \in M \times S_\circ$ in $\mathcal{M}^T$, there are
successors $q_1, q_2$ such that $q_1 \in M \times \{s_0\}$ and $q_2 \in M \times \{t\}$; moreover, all successors are in $M \times \{s_0, t\}$. (ii) The state $t$ has the unique successor $s_0$. Hence, in $M^T$, from all states $(m, t) \in M \times S_O$, all successors $q$ are such that $q \in M \times \{s_0\}$.

To establish (1), let $(m, s_0) \in M \times \{s_0\}$ be some state in the Markov chain $M^T$. For the case $\mathcal{P}_{M^T,(m_0,s_0)}(GF(m, s_0)) = 0$, we trivially have (1). Therefore, we assume that $\mathcal{P}_{M^T,(m_0,s_0)}(GF(m, s_0)) > 0$. So, there exists an infinite path satisfying $GF(m, s_0)$. Hence, there exists a finite path $\pi$ from $(m, s_0)$ to itself. By the structure of the chain, $\pi$ has a prefix $(m_s, s_0)(m_s, s_1) \cdots (m_s, s_i) \in (M \times S_O)^*$ traversing player’s states, with $m_s = m$, and next visiting some state $(m_r, r_i)$. By property (i), we know that $(m_r, r_i)$ has a successor $(m_r, t) \in M \times \{t\}$. It implies that

$$\pi_t = (m_s, s_0)(m_s, s_1) \cdots (m_s, s_i)(m_r, r_i)(m_r, t)$$

is a finite path in $M^T$, starting in $(m, s_0)$ with positive probability. By strong fairness of probabilistic choices, $\mathcal{P}_{M^T,(m_0,s_0)}(GF\pi_t) = \mathcal{P}_{M^T,(m_0,s_0)}(GF(m, s_0))$. In other words, under the condition that state $(m, s_0)$ is visited infinitely often, the finite path $\pi_t$ will be taken infinitely often, almost-surely. As an immediate result of this and the fact that $(m_t, t)$ has color 3, we conclude (1).

To establish (2), we need to show that, for all infinite plays $\pi^\infty$ of $M^T$, if $\pi^\infty \in \{GF[M \times S]^{Col=2}\}$ then $\pi^\infty \in \{GF[M \times \{s_0\}]\}$. By the properties (i) and (ii) of $M^T$, whenever $\pi^\infty$ visits some state from $M \times S_O$, which are the only states with color 2, then $\pi^\infty$ must visit some state from $M \times \{s_0\}$ within two steps. This results in (2).

To establish (3), we observe that the inclusion $\supseteq$ is trivial. To show $\subseteq$, let $\pi^\infty \in \{GF[M \times \{s_0\}]\}$ be an infinite path in the chain. As $\pi^\infty$ visits infinitely many elements from the finite set $M \times \{s_0\}$, there must exist some element $(m, s_0)$ that is visited infinitely often. Hence, $\pi^\infty \in \{GF[M \times \{s_0\}]\}$, which gives the inclusion and thus (3).

Now, we construct an HD-strategy $\sigma_h$ such that $\mathcal{P}_{M,s_0,\sigma_h}(\varphi) = 1$. The strategy $\sigma_h$ is defined, for all partial plays $\rho$, as follows:

$$\sigma_h(\rho) = \begin{cases} r_0 & \text{if } \rho = s_0 \\ r_k & \text{if there exists } k > 0 \text{ such that } \rho = (s_0(S \setminus \{s_0\})^*)^{k-1}s_0s_1\cdots s_k \\ s_j & \text{otherwise, where the last state visited by } \rho \text{ is } s_j-1. \end{cases}$$

Intuitively, $\sigma_h$ is such that upon the $k$-th visit to state $s_0$, the path $s_0s_1\cdots s_k$ is traversed and then the transition $s_k \rightarrow r_k$ is chosen. Observe that $\mathcal{P}_{M,s_0,\sigma_h}(GF[S]^{Col=2}) = 1$. Below, we argue that $\mathcal{P}_{M,s_0,\sigma_h}(GF[S]^{Col=3}) = 0$, which proves that $\mathcal{P}_{M,s_0,\sigma_h}(\varphi) = 1$.

We define the sequence of events $E_k$ of visiting $t$ between the $k$th and $k+1$st visits of $s_0$. For $k \geq 1$, let

$$E_k = (s_0(S \setminus \{s_0\})^*)^{k-1}s_0(S \setminus \{s_0, t\})^*ts_0S^\omega.$$

Observe that

$$\bigcap_{n=1}^{\infty} \bigcup_{k \geq n} \pi(E_k) = \{GF(t)\} = \{GF[S]^{Col=3}\}.$$

We use the Borel-Cantelli lemma to prove that infinitely many of $E_k$’s occur with zero probability, that is the probability of $GF[S]^{Col=3}$. In fact, by construction of $\sigma_h$, observe that $\mathcal{P}_{M,s_0,\sigma_h}(E_k) = 2^{-k}$. Consequently, we have

$$\sum_{k=1}^{\infty} \mathcal{P}_{M,s_0,\sigma_h}(E_k) = 1 + \frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \cdots = 2 < \infty.$$

By the Borel-Cantelli lemma, we then have $\mathcal{P}_{M,s_0,\sigma_h}(GF[S]^{Col=3}) = 0$, and thus $\mathcal{P}_{M,s_0,\sigma_h}(\varphi) = 1$. The proof is complete. □

Remark 1. A strict subclass of finitely branching MDPs are 1-counter MDPs, where a finite-state MDP is augmented with an integer counter [6]. The MDP in Theorem 1 (plus some auxiliary states) is implementable by a 1-counter MDP.

Proof. Figure 5 shows a 1-counter MDP with control-states $\{s, r, r', t\}$ that is functionally equivalent to the one used in Theorem 1 (Figure 2 (left)). It just uses some auxiliary states that have no influence on the parity objective. Starting in $s$, the player can choose whether to increase the counter by 1 or to go to $r$. In the random state $r$ the behavior depends on the counter value. If the counter is non-zero then the successors $r, r'$ are chosen with equal probability and the counter is decreased by 1. If the counter is zero then $t$ is the unique successor. In state $r'$ the counter is deterministically decreased until it becomes zero, and then one goes to state $s$. The color function is $Col(s) = 1$, $Col(r) = Col(r') = 2$ and $Col(t) = 3$. If one is in state $r$ with counter value $n$, then the probability of seeing state $t$ before returning to state $s$ is $2^{-n}$. □
Fig. 5: A 1-counter MDP implementing an MDP similar to the one used in Theorem 1. The dashed transitions (labeled with zero test = 0?) are taken when the counter is zero.

**Theorem 3.** Let $\varphi = \text{Safety}(T)$. There exists an infinitely branching MDP $\mathcal{M}$ with initial state $s$ such that

- for all FR-strategies $\sigma$, we have $P_{\mathcal{M},s,\sigma}(\varphi) = 0$,
- for every $c \in [0, 1]$, there exists an HD-strategy $\sigma$ such that $P_{\mathcal{M},s,\sigma}(\varphi) \geq c$.

Hence, $c$-optimal strategies for safety require infinite memory.

**Proof.** Consider the MDP $\mathcal{M}$ shown (on the left side) in Figure 3 where $S_T = \{s, t\}$ and $S_0 = \{r_i\}_{i \geq 1}$. The state $s$ is infinitely branching: $s \rightarrow r_i$ for all $i \geq 1$.

For all random states $r_i$, there are two successors $P(r_i, t) = 2^{-i}$ and $P(r_i, s_0) = 1 - 2^{-i}$. The state $t$ is a sink state.

Let $\sigma$ be an arbitrary FR-strategy. By definition there is a transducer $T$ with finite memory $M$ and initial mode $m_0$ such that $\sigma_T = \sigma$. Let $\mathcal{M}^T$ be the Markov chain obtained by the product of the MDP $\mathcal{M}$ and the transducer $T$. The set of states in $\mathcal{M}^T$ is $M \times S$. We denote by $P_{\mathcal{M}^T, q}(\mathcal{R})$ the probability of a measurable set $\mathcal{R}$ of infinite paths (i.e., infinite plays), starting in the state $q$ of $\mathcal{M}^T$.

Below, we prove that $P_{\mathcal{M},s,\sigma}(\text{Safety}(\{t\})) = 0$. Equivalently, we show that $P_{\mathcal{M}^T,(m_0,s)}(\text{Safety}(M \times \{t\})) = 0$. We proceed in the following three steps: We will show, using strong fairness of probabilistic choices in Markov chains, that for all modes $m \in M$,

\[ P_{\mathcal{M}^T,(m_0,s)}(\text{GF}(m, s) \land \neg \text{F}(M \times \{t\})) = 0. \quad (4) \]

Since the memory of strategy $\sigma$ is finite ($|M| < \infty$), we will show that

\[ \| \text{GF}(M \times \{s\}) \| = \bigvee_{m \in M} \text{GF}(m, s) \]  

Moreover, we will prove that

\[ \| \neg \text{F}(M \times \{t\}) \| \subseteq \| \text{GF}(M \times \{s\}) \|. \quad (5) \]

Using (4)-(6), we complete the proof as follows:

\[ P_{\mathcal{M}^T,(m_0,s)}(\neg \text{F}(M \times \{t\})) = P_{\mathcal{M}^T,(m_0,s)}(\text{GF}(M \times \{s\}) \land \neg \text{F}(M \times \{t\})) \]  

by (6)

\[ = P_{\mathcal{M}^T,(m_0,s)}(\bigvee_{m \in M} (\text{GF}(m, s_0) \land \neg \text{F}(M \times \{t\}))) \]  

by (5)

\[ \leq \sum_{m \in M} P_{\mathcal{M}^T,(m_0,s)}(\text{GF}(m, s_0) \land \neg \text{F}(M \times \{t\})) \]  

union bound

\[ = 0. \quad \text{by (4)} \]

As a result, $P_{\mathcal{M},s_0,\sigma}(\text{Safety}(\{t\})) = 0$. Below we prove (4), (5) and (6).

To establish (4), let $(m, s) \in M \times \{s\}$ be some state in the Markov chain $\mathcal{M}^T$. For the case $P_{\mathcal{M}^T,(m_0,s)}(\text{GF}(m, s)) = 0$, we trivially have (4). Therefore, we assume that $P_{\mathcal{M}^T,(m_0,s)}(\text{GF}(m, s)) > 0$. So, there exists an infinite path satisfying $\text{GF}(m, s)$. Hence, there exists a finite path $\pi$ from $(m, s)$ to itself. By the structure of the chain, $\pi$ visits some state $(m_{r_i}, r_i)$. In $\mathcal{M}$, for all random states $r_i \in S_0$, the only successors are $s$ and $t$. Hence, from all states $(m, r_i) \in M \times S_0$ in $\mathcal{M}^T$, there is some successor $q \in M \times \{t\}$. It implies that $\pi_t = (m, s)(m_{r_i}, r_i)(m_{r_i}, t)$ is a finite path in $\mathcal{M}^T$, that starts in $(m, s)$ with positive probability. By strong fairness of probabilistic choices, $P_{\mathcal{M}^T,(m_0,s)}(F\pi_t) = P_{\mathcal{M}^T,(m_0,s)}(\text{GF}(m, s))$. As a result, we conclude (4).
To establish (5), we observe that the inclusion \( \subseteq \) is trivial. To show \( \subseteq \), let \( \pi^\infty \in [\text{GF}(M \times \{s_0\})] \) be an infinite path in the chain. As \( \pi^\infty \) visits infinitely many elements from the finite set \( M \times \{s_0\} \), there must exist some element \((m, s_0)\) that is visited infinitely often. Hence, \( \pi^\infty \in \bigvee_{m \in M} \text{GF}(m, s_0) \), which gives the inclusion and thus (5).

To establish (6), note that in \( \mathcal{M} \), all successors of \( s \) are from random states, from which the only successors are \( s \) and \( t \). Hence, for all infinite plays \( \pi^\infty \not\in [\text{F}(M \times \{t\})] \) in \( \mathcal{M}^T \), \( \pi^\infty \) alternates between some states from \( M \times \{s\} \) and next some state from \( M \times S_C \). It implies that \( \pi \in [\text{GF}(M \times \{s\})] \).

Now, let \( c \in [0, 1) \). Let \( n \) be such that \( c \leq 1 - \frac{1}{2^n} \). We construct an HD-strategy \( \sigma_n \) such that

\[
\mathcal{P}_{M, s, \sigma_n}(\varphi) \geq 1 - \frac{1}{2^n} \geq c.
\]

For all partial plays \( \rho \), we define \( \sigma_n(\rho) = r_{n+k} \) where \( k \) is the number of times \( \rho \) has visited \( s \). Intuitively, upon the \( k \)-th visit to \( s \), the strategy \( \sigma_n \) chooses the transition \( s \rightarrow r_{n+k} \).

For all \( k \geq 1 \), let \( E_k \) be the event of visiting \( t \) after the \( k \)-th visit of \( s \), defined as follows

\[
E_k = (sS_C)^{k-1}sS_Ct^\omega.
\]

Observe that \([\text{F}(t)]\) is the disjoint union of all \( E_k \) events. Hence, we have

\[
\mathcal{P}_{M, s, \sigma_n}(\text{F}(t)) = \sum_{k=1}^{\infty} \mathcal{P}_{M, s, \sigma_n}(E_k) = \sum_{k=1}^{\infty} \mathcal{P}_{M, s, \sigma_n}([sS_C]^{k-1}sS_C] \cdot \frac{1}{2^{n+k}} \leq \sum_{k=1}^{\infty} \frac{1}{2^{n+k}} = \frac{1}{2^n}\]

This proves that \( \mathcal{P}_{M, s, \sigma_n}(\text{Safety}([t])) = 1 - \mathcal{P}_{M, s, \sigma_n}(\text{F}(t)) \geq 1 - \frac{1}{2^n} \geq c \).

**Theorem 4.** Let \( \varphi = \{0, 1\}\)-Parity be the co-Büchi objective. There exists an infinitely branching MDP \( \mathcal{M} \) with initial state \( s \) such that

- for all FR-strategies \( \sigma \), we have \( \mathcal{P}_{M, s, \sigma}(\varphi) = 0 \),
- there exists an HD-strategy \( \sigma \) such that \( \mathcal{P}_{M, s, \sigma}(\varphi) = 1 \).

Hence, optimal (and even almost-surely winning) strategies and \( \epsilon \)-optimal strategies for co-Büchi require infinite memory.

**Proof.** Consider the MDP \( \mathcal{M} \) shown (on the right side) in Figure 3 where \( S_0 = \{s, t\} \) and \( S_C = \{r_i\}_{i \geq 1} \). The state \( s \) is infinitely branching: \( s \rightarrow r_i \) for all \( i \geq 1 \). For all random states \( r_i \), there are two successors \( P(r_i, t) = 2^{-i} \) and \( P(r_i, s) = 1 - 2^{-i} \). The state \( t \) has the unique successor \( s \).

Let \( \sigma \) be an arbitrary FR-strategy. By definition there is a transducer \( T \) with finite memory \( M \) and initial mode \( m_0 \) such that \( \sigma_T = \sigma \). Let \( \mathcal{M}^T \) be the Markov chain obtained by the product of the MDP \( M \) and the transducer \( T \). The set of states in \( \mathcal{M}^T \) is \( M \times S \). We define a coloring function for \( \mathcal{M}^T \) such that it ignores the memory mode and assigns to \((m, s)\) the same color as state \( s \) in \( M \). In particular, all states \( q \in M \times \{t\} \) have color 1. We use the same notation \( \text{Col} \) for the coloring functions of both \( M \) and \( \mathcal{M}^T \). We denote by \( \mathcal{P}_{\mathcal{M}, s, \sigma}(\text{Col}) \) the probability of a measurable set \( \text{Col} \) of infinite paths (i.e., infinite plays), starting in the state \( q \) of \( \mathcal{M}^T \).

We prove that \( \mathcal{P}_{M, s, \sigma}(\text{FG}[S]^\text{Col} \neq 1) = 0 \). Equivalently, we show that \( \mathcal{P}_{\mathcal{M}^T, (m_0, s)}(\text{FG}[M \times S]^\text{Col} \neq 1) = 0 \). We proceed in three steps: we will show, using strong fairness of probabilistic choices in Markov chains, that for all modes \( m \in M \),

\[
\mathcal{P}_{\mathcal{M}^T, (m_0, s)}(\text{GF}(m, s) \land \text{FG}[M \times S]^\text{Col} \neq 1) = 0.
\]

Since the memory of strategy \( \sigma \) is finite (\( |M| < \infty \)), we will show that

\[
[\text{GF}(M \times \{s\})] = \bigvee_{m \in M} \text{GF}(m, s)
\]
Moreover, we show that
\[ \left[ \text{FG}[M \times S]^{\text{Col} \neq 1} \right] \subseteq \left[ \text{GF}(M \times \{s\}) \right] \] (9)

Using (7), (8) and (9), we complete the proof as follows:
\[
P \mathcal{M},(m_0,s) \left( \text{FG}[M \times S]^{\text{Col} \neq 1} \right) =
\]
\[
= P \mathcal{M},(m_0,s)(\bigvee_{m \in M} (\text{GF}(m,s) \wedge \text{FG}(M \times S)^{\text{Col} \neq 1}))
\]
\[
\leq \sum_{m \in M} P \mathcal{M},(m_0,s)(\text{GF}(m,s) \wedge \text{FG}(M \times S)^{\text{Col} \neq 1})
\]
\[
= 0.
\]

As a result, \( P \mathcal{M},s,\sigma(\varphi) = 0 \). Below we prove (7), (8) and (9).

To establish (7), let \((m, s) \in M \times \{s\}\) be some state in the Markov chain \(\mathcal{M}^T\). For the case \(P \mathcal{M},(m_0,s)(\text{GF}(m,s)) = 0\), we trivially have (7). Therefore, we assume that \(P \mathcal{M},(m_0,s)(\text{GF}(m,s)) > 0\). So, there exists an infinite path satisfying \(\text{GF}(m,s)\). Hence, there exists a finite path \(\pi\) from \((m, s)\) to itself. By the structure of the chain, \(\pi\) visits some state \((m_{r_1}, r_1)\). Consider the MDP \(\mathcal{M}\), for all random states \(r_1 \in S_0\), the only successors are \(s\) and \(t\). Hence, from all states \((m, r_1) \in M \times S_0\) in \(\mathcal{M}^T\), there is some successor \(q\) such that \(q \in M \times \{t\}\). It implies that \(\pi_i = (m, s)(m_{r_1}, r_1)(m_t, t)\) is a finite path in \(\mathcal{M}^T\), that starts in \((m, s)\) with positive probability. By strong fairness of probabilistic choices, \(P \mathcal{M},(m_0,s)(\text{FG}_\pi) = P \mathcal{M},(m_0,s)(\text{GF}(m, s))\). As a result, we conclude (7).

To establish (8), we observe that the inclusion \(\supseteq\) is trivial. To show \(\subseteq\), let \(\pi^\infty \in \left[ \text{GF}(M \times \{s\}) \right]\) be an infinite path in the chain. As \(\pi^\infty\) visits infinitely many elements from the finite set \(M \times \{s\}\), there must exist some element \((m, s)\) that is visited infinitely often. Hence, \(\pi^\infty \in \left[ \bigvee_{m \in M} \text{GF}(m,s) \right]\), which gives the inclusion and thus (8).

To establish (9), note that in \(\mathcal{M}\), all infinite plays must visit \(s\) infinitely often. Thus all runs in \(\mathcal{M}^T\) must visit \(M \times \{s\}\) infinitely often. In particular, this holds for those infinite runs \(\pi^\infty \in \left[ \text{FG}(M \times S)^{\text{Col} \neq 1} \right]\) in \(\mathcal{M}^T\). Thus \(\pi^\infty \in \left[ \text{GF}(M \times \{s\}) \right]\).

Now we construct an HD-strategy \(\sigma_h\) such that \(P \mathcal{M},s,\sigma_h(\varphi) = 1\). For all partial plays \(\rho\), we define \(\sigma_h(\rho) = r_k\) where \(k\) is the number of times \(\rho\) has visited \(s\). Intuitively, upon the \(k\)-th visit to \(s\), the strategy \(\sigma_h\) chooses the transition \(s \rightarrow r_k\). Below we argue that \(P \mathcal{M},s,\sigma_h(\text{GF}[S]^{\text{Col} = 1}) = 0\), which proves that \(P \mathcal{M},s,\sigma_h(\varphi) = 1\).

We define the sequence of events \(E_k\) of visiting \(t\) between the \(k\)-th and \(k + 1\)st visits of \(s\). For \(k \geq 1\), let
\[ E_k = (s(S_0 + S_0 t))^{k-1} s(S_0 t)sS^\omega. \]

Observe that
\[ \bigcap_{n=1}^{\infty} \bigcup_{k \geq n} \pi(E_k) = \left[ \text{GF}([t]) \right] = \left[ \text{GF}[S]^{\text{Col} = 1} \right]. \]

We use the Borel-Cantelli lemma to prove that infinitely many of \(E_k\)’s occur with zero probability, that is the probability of \(\text{GF}[S]^{\text{Col} = 1}\). In fact, by construction of \(\sigma_h\), observe that \(P \mathcal{M},s,\sigma_h(E_k) = 2^{-k}\). Consequently, we have
\[ \sum_{k=1}^{\infty} P \mathcal{M},s_0,\sigma_h(E_k) = \frac{1}{2} + \frac{1}{4} + \frac{1}{8} + \cdots = 1 < \infty. \]

By the Borel-Cantelli lemma, we then have \(P \mathcal{M},s,\sigma_h(\text{GF}[S]^{\text{Col} = 1}) = 0\), and thus \(P \mathcal{M},s,\sigma_h(\varphi) = 1\). The proof is complete. \(\square\)

B. Proofs of Section IV

Lemma 20. Let \(\varphi\) be an objective that is prefix-independent in a class C of MDPs. Let \(\mathcal{M} = (S, S_0, S_\omega, \rightarrow, P) \in C\), and \(s_0 \in S\), and \(s\) be an objective with a strategy \(\mathcal{M},s_0,\sigma(\varphi) = \text{val}_\mathcal{M}(s_0)\). Suppose that \(s_0 s_1 \cdots s_n\) for some \(n \geq 0\) is a partial play starting in \(s_0\) and induced by \(\sigma\). Then:

1) \(\text{val}_\mathcal{M}(s_n) = P \mathcal{M},s_0,\sigma(\varphi^{s_0} | s_0 s_1 \cdots s_n S^\omega)\).
2) If \(s_n \in S_0\) then \(\text{val}_\mathcal{M}(s_n) = \sum_{s_{n+1} \in S} P(s_{n+1})(s_{n+1}) \cdot \text{val}_\mathcal{M}(s_{n+1})\).
3) If \(s_n \in S_0\), then \(\text{val}_\mathcal{M}(s_n) = \text{val}_\mathcal{M}(s_n)\) for all \(s_n \in S\) and \(\text{supp}(\sigma(s_0 s_1 \cdots s_n))\).

Proof. First we show \(P \mathcal{M},s_0,\sigma(\varphi^{s_0} | s_0 s_1 \cdots s_n S^\omega) \leq \text{val}_\mathcal{M}(s_n)\). Define a strategy \(\sigma' : S^* S_0 \rightarrow D(S)\) by \(\sigma'(w) = \sigma(s_0 s_1 \cdots s_{n-1} w)\) for all \(w \in S^* S_0\). Then we have \(P \mathcal{M},s_0,\sigma(\varphi^{s_0} | s_0 s_1 \cdots s_n S^\omega) = P \mathcal{M},s_0,\sigma'(\varphi^{s_0}) \leq \text{val}_\mathcal{M}(s_n)\).

Next we show \(\text{val}_\mathcal{M}(s_n) \leq P \mathcal{M},s_0,\sigma(\varphi^{s_0} | s_0 s_1 \cdots s_n S^\omega)\). Towards a contradiction, suppose that \(\text{val}_\mathcal{M}(s_n) > P \mathcal{M},s_0,\sigma(\varphi^{s_0} | s_0 s_1 \cdots s_n S^\omega)\). Then, by the definition of \(\text{val}_\mathcal{M}(s_n)\), there is a strategy \(\sigma'\) with \(P \mathcal{M},s_0,\sigma'(\varphi^{s_0}) > \text{val}_\mathcal{M}(s_n)\).
Define a strategy $\sigma''$ that plays according to $\sigma$; if and when partial play $s_0 s_1 \cdots s_n$ is played, then $\sigma''$ acts like $\sigma'$ henceforth; otherwise $\sigma''$ continues with $\sigma$ forever. Using prefix-independence we get:

\[
P_{\mathcal{M},s_0,\sigma''}(\varphi_s^{\text{val}} | s_0 s_1 \cdots s_n S^\omega) = P_{\mathcal{M},s_0,\sigma''}(\varphi_s^{\text{val}} | s_0 s_1 \cdots s_n S^\omega) \cdot P_{\mathcal{M},s_0,\sigma'}(s_0 s_1 \cdots s_n S^\omega) + P_{\mathcal{M},s_0,\sigma''}(\varphi_s^{\text{val}} \setminus s_0 s_1 \cdots s_n S^\omega) = P_{\mathcal{M},s_0,\sigma}(\varphi_s^{\text{val}}) \cdot P_{\mathcal{M},s_0,\sigma}(s_0 s_1 \cdots s_n S^\omega) + \sum_{s_{n+1} \in S} P_{\mathcal{M},s_0,\sigma}(s_0 s_1 \cdots s_n s_{n+1} S^\omega)
\]

This contradicts the definition of $\text{val}_\mathcal{M}(s_0)$. Hence we have shown item 1.

Towards items 2 and 3, we extend $\sigma : S^* S \rightarrow \mathcal{D}(S)$ to $\sigma : S^* S \rightarrow \mathcal{D}(S)$ by defining $\sigma(w s) = P(s)$ for $w \in S^*$ and $s \in S_\square$. Then we have for all $s_{n+1} \in S$:

\[
P_{\mathcal{M},s_0,\sigma}(s_0 s_1 \cdots s_n s_{n+1} S^\omega) = P_{\mathcal{M},s_0,\sigma}(s_0 s_1 \cdots s_n S^\omega) \cdot \sigma(s_0 s_1 \cdots s_n)(s_{n+1})
\]

Further we have:

\[
\text{val}_\mathcal{M}(s_n)
\]

Thus we have shown item 2. Towards item 3, suppose $s_n \in S_\square$. Then prefix-independence implies $\text{val}_\mathcal{M}(s_n) \geq \text{val}_\mathcal{M}(s_{n+1})$ for all $s_{n+1}$ with $s_n \rightarrow s_{n+1}$. Since $\sigma(s_0 s_1 \cdots s_n)$ is a probability distribution, the equality chain above shows that $\text{val}_\mathcal{M}(s_n) = \text{val}_\mathcal{M}(s_{n+1})$ for all $s_{n+1} \in \text{supp}(\sigma(s_0 s_1 \cdots s_n))$. Thus we have shown item 3.

**Lemma 6.** Let $\varphi$ be an objective that is prefix-independent in a class $C$ of MDPs. Let $\mathcal{M} = (S, S_\square, S_\circ, \rightarrow, P) \in \mathcal{C}$. Construct an MDP $\mathcal{M}_s = (S_s, S_s \cap S_\circ, S_s \cap S_\square, \rightarrow_s, P_s)$ by setting

\[
S_s = \{ s \in S | \exists \sigma. P_{\mathcal{M},s,\sigma}(\varphi) = \text{val}_\mathcal{M}(s) > 0 \}
\]

and $S_s \cap S_\square$ and $S_s \cap S_\circ$ and

\[
\rightarrow_s = \{ (s,t) \in S_s \times S_s | s \rightarrow t \text{ and if } s \in S_s \rightarrow_s \text{ then } \text{val}_\mathcal{M}(s) = \text{val}_\mathcal{M}(t) \}
\]

and $P_s : S_s \rightarrow \mathcal{D}(S_s)$ so that

\[
P_s(s)(t) = \frac{\text{val}_\mathcal{M}(t)}{\text{val}_\mathcal{M}(s)}
\]

for all $s \in S_s$ and $t \in S_s$ with $s \rightarrow_s t$. Then:
1) For all $\sigma \in \Sigma_M$ and all $n \geq 0$ and all $s_0, \ldots, s_n \in S_*$ with $s_0 \rightarrow^*_s s_1 \rightarrow^*_s \cdots \rightarrow^*_s s_n$:

$$P_{M,s_0,\sigma}(s_0s_1 \cdots s_n S^\omega) = \frac{\operatorname{val}_M(s_n)}{\operatorname{val}_M(s_0)}$$

2) For all $s_0 \in S_*$ and all $\sigma \in \Sigma_M$ with $P_{M,s_0,\sigma} (\varphi) = \operatorname{val}_M(s_0) > 0$ and all measurable $\Re \subseteq s_0 S^\omega$ we have $P_{M,s_0,\sigma}(\Re) = P_{M,s_0,\sigma}(\Re \mid \varphi)^n$.

**Proof.** Note that by Lemma 20.2 we have that $P_*(s)$ is a probability distribution for all $s \in S_* \Re$; hence the MDP $M_*$ is well-defined.

We prove item 1 by induction on $n$. For $n = 0$ it is trivial. For the step, suppose that the equality in item 1 holds for some $n$. If $s_n \in S_* \Re$ then we have:

$$P_{M,s_0,\sigma}(s_0s_1 \cdots s_n s_{n+1} S^\omega) = \frac{\operatorname{val}_M(s_n)}{\operatorname{val}_M(s_0)} P_*(s_{n+1})$$

This completes the inductive step, and we have proved item 1.

Towards item 2, let $s_0 \in S_*$ and $\sigma \in \Sigma_M$ such that $P_{M,s_0,\sigma}(\varphi) = \operatorname{val}_M(s_0) > 0$. Observe that $\sigma$ can be applied also in the MDP $M_*$. Indeed, for any $s \in S_* \Re$, if $t$ is a possible successor state of $s$ under $\sigma$, then $\operatorname{val}_M(s) = \operatorname{val}_M(t)$ by Lemma 20.3 and thus $t \in S_* \Re$.

Let again $n \geq 0$ and $s_0, s_1, \ldots, s_n \in S$.

- Suppose $s_0 \cdots s_n$ is a partial play in $M_*$ induced by $\sigma$. Then we have:

$$P_{M,s_0,\sigma}(s_0s_1 \cdots s_n S^\omega) \cdot P_{M,s_0,\sigma}(\varphi) = \frac{\operatorname{val}_M(s_n)}{\operatorname{val}_M(s_0)} P_{M,s_0,\sigma}(\varphi) \quad \text{item 1}$$

- Suppose $s_0 \cdots s_n$ is not a partial play in $M_*$ induced by $\sigma$. Hence $P_{M,s_0,\sigma}(s_0s_1 \cdots s_n S^\omega) = 0$. If $s_0 \cdots s_n$ is not a partial play in $M_*$ induced by $\sigma$ then $P_{M,s_0,\sigma}(s_0s_1 \cdots s_n S^\omega) = 0$. Otherwise, since $\sigma$ is optimal, there is $i \leq n$ with $\operatorname{val}_M(s_i) = 0$, hence $P_{M,s_0,\sigma}(s_0 \cdots s_n S^\omega) = 0$. In either case we have $P_{M,s_0,\sigma}(s_0 \cdots s_n S^\omega) = 0 = P_{M,s_0,\sigma}(\Re | \varphi)^n \cap s_0 s_1 \cdots s_n S^\omega)$. In either case we have the equality $P_{M,s_0,\sigma}(\Re) = P_{M,s_0,\sigma}(\Re \mid \varphi)^n$ for cylinders $\Re = s_0 s_1 \cdots s_n S^\omega$. Since probability measures extend uniquely from cylinders [5], the equality holds for all measurable $\Re \subseteq s_0 S^\omega$. Thus we have shown item 2.

**Lemma 8.** Let $S$ be countable and $s \in S$. Call a set of the form $swS^\omega$ for $w \in S^*$ a cylinder. Let $P, P'$ be probability measures on $SS^\omega$ defined in the standard way, i.e., first on cylinders and then extended to all measurable sets $\Re \subseteq SS^\omega$.  

...
Suppose there is $x \geq 0$ such that $x \cdot P(\mathcal{C}) \leq P'(\mathcal{C})$ for all cylinders $\mathcal{C}$. Then $x \cdot P(\mathcal{R}) \leq P'(\mathcal{R})$ holds for all measurable $\mathcal{R} \subseteq sS^\omega$.

Proof. Let $\mathcal{C} = \{ \mathcal{C} \subseteq sS^\omega \mid \mathcal{C} \text{ cylinder} \}$ denote the class of cylinders. This class generates an algebra $\mathcal{C}_* \supseteq \mathcal{C}$, which is the closure of $\mathcal{C}$ under finite union and complement. The classes $\mathcal{C}$ and $\mathcal{C}_*$ generate the same $\sigma$-algebra $\sigma(\mathcal{C})$. The class $\mathcal{C}_*$ is the set of finite disjoint unions of cylinders [5, Section 2]. Hence $x \cdot P(\mathcal{R}) \leq P'(\mathcal{R})$ for all $\mathcal{R} \in \mathcal{C}_*$.

Define

$$Q = \{ \mathcal{R} \in \sigma(\mathcal{C}) \mid x \cdot P(\mathcal{R}) \leq P'(\mathcal{R}) \}.$$ 

We have $\mathcal{C} \subseteq \mathcal{C}_* \subseteq Q \subseteq \sigma(\mathcal{C})$. We show that $Q$ is a monotone class, i.e., if $\mathcal{R}_1, \mathcal{R}_2, \ldots \in Q$, then $\mathcal{R}_1 \subseteq \mathcal{R}_2 \subseteq \cdots$ implies $\bigcup_i \mathcal{R}_i \in Q$. Suppose $\mathcal{R}_1, \mathcal{R}_2, \ldots \in Q$ and $\mathcal{R}_1 \subseteq \mathcal{R}_2 \subseteq \cdots$. Then:

$$x \cdot P\left(\bigcup_i \mathcal{R}_i\right) = \sup_i x \cdot P(\mathcal{R}_i) \leq \sup_i P'(\mathcal{R}_i) \leq P'(\bigcup_i \mathcal{R}_i),$$

measures are continuous from below

$$\implies x \cdot P(\mathcal{R}_i) \leq P'(\mathcal{R}_i),$$

definition of $Q$

$$= \sup_i \left( x \cdot P(\mathcal{R}_i) \right) \leq \sup_i P'(\mathcal{R}_i),$$

measures are continuous from below

So $\bigcup_i \mathcal{R}_i \in Q$. Using the fact that measures are continuous from above, one can similarly show that if $\mathcal{R}_1, \mathcal{R}_2, \ldots \in Q$ and $\mathcal{R}_1 \supseteq \mathcal{R}_2 \supseteq \cdots$ then $\bigcap_i \mathcal{R}_i \in Q$. Hence $Q$ is a monotone class.

Now the monotone class theorem (see, e.g., [5, Theorem 3.4]) implies that $\sigma(\mathcal{C}) \subseteq Q$, thus $Q = \sigma(\mathcal{C})$. Hence $x \cdot P(\mathcal{R}) \leq P'(\mathcal{R})$ for all $\mathcal{R} \in \sigma(\mathcal{C})$. □

C. Proofs of Section V

Theorem 10. Let $\mathcal{M} = (S, S_0, S_\omega, \to, P)$ be an MDP, and $\varphi = \operatorname{Reach}(T)$. Let $s_0 \in S$ and $\sigma$ be a strategy with $P_{M,s_0}(\varphi) = 1$. Then there is an MD-strategy $\hat{\sigma}$ with $P_{M,s_0}(\varphi) = 1$.

Proof. We can assume that $T = \{ t \}$ for some $t \in S$. We can also assume that all states are almost-surely winning, since in order to achieve an almost-sure winning objective, the player must forever remain in almost-surely winning states.

Let $\epsilon_1 := 1/2$. By Theorem 9 there exists an MD-strategy $\sigma_1$ such that $P_{M,s_0,\sigma_1}(\varphi) \geq 1 - \epsilon_1$. In fact, by the proof of Theorem 9 there exists a finite subset $V_1 \subseteq S$ such that $P_{M,s_0,\sigma_1}(\varphi^{s_0} \cap V_1^\omega) \geq 1 - \epsilon_1$. Let $U_1$ denote the states that occur in those plays that are both contained in $\varphi^{s_0} \cap V_1^\omega$ and induced by $\sigma_1$. Then $P_{M,s_0,\sigma_1}(\varphi^{s_0} \cap U_1^\omega) = P_{M,s_0,\sigma_1}(\varphi^{s_0} \cap V_1^\omega) \geq 1 - \epsilon_1$. By the definition of $U_1$, for all $s \in U_1$ the MD-strategy $\sigma_1$ induces a play from $s_0$ to $t$ via $s$. Hence we have $P_{M,s_0,\sigma_1}(\varphi^{s_0} \cap U_1^\omega) > 0$. Since $U_1 \subseteq V_1$ is finite, there are $c > 0$ and $\ell \in \mathbb{N}$ such that for all $s \in U_1$ we have $P_{M,s_0,\sigma_1}(sU_1^{t=1}\{t\}^\omega) \geq c$, i.e., from any state in $U_1$ the probability that $t$ is reached in $\leq \ell$ steps is at least $c$.

Consider the MDP $M_1$ obtained from $M$ by fixing $\sigma_1$ on $U_1$ (i.e., in $M_1$ the states in $U_1$ are random states). We argue that all states are almost-surely winning in $M_1$. Indeed, let $s \in S$ be any state. Recall that $s$ is almost-surely winning in $M$. Define an HR-strategy $\sigma$ in $M_1$ as follows: first play a strategy that is almost-surely winning for $s$ in $M$; if and when $U_1$ is entered and then left again (entering some state $s' \in S \setminus U_1$) then forget the history and play again a strategy that is almost-surely winning for $s'$ in $M$; and so forth. This strategy $\sigma$ reaches $\{ t \}$ with probability 1 whenever the play stays outside of $U_1$. I.e., almost all plays that eventually always avoid $U_1$ reach $\{ t \}$. Moreover, whenever the play enters $U_1$, the probability that $t$ is reached in $\leq \ell$ steps is at least $c$, i.e., there is a uniform bound. Thus almost all plays that enter $U_1$ infinitely often reach $\{ t \}$. It follows that we have $P_{M_1,s,\sigma}(\varphi) = 1$.

Now we repeat the argument, but with $M_1$ instead of $M$ and with $\epsilon_2 = 1/4$ instead of $\epsilon_1$. This yields a set $U_2 \supseteq U_1$ and an MD-strategy $\sigma_2$ that agrees with $\sigma_1$ on $U_1$ so that $P_{M,s_0,\sigma_2}(\varphi^{s_0} \cap U_2^\omega) = P_{M,s_0,\sigma_2}(\varphi^{s_0} \cap U_1^\omega) = 1 - \epsilon_2$. Similarly as before, obtain an MDP $M_2$ from $M_1$ by fixing $\sigma_2$ on $U_2$. Then repeat again, and so forth, with $\epsilon_i = 1/2^i$ for $i = 1, 2, \ldots$

Define $U := \bigcup_{i \geq 1} U_i$. Observe that on all $s \in U$ almost all (i.e., all except finitely many) strategies $\sigma_i$ agree. Let $\hat{\sigma}$ be an MD-strategy that on all states in $U$ agrees with almost all MD-strategies $\sigma_i$. By our construction we have $P_{M,s_0,\hat{\sigma}}(\varphi^{s_0} \cap U^\omega) \geq 1 - \epsilon$ for all $\epsilon > 0$. Hence $P_{M,s_0,\hat{\sigma}}(\varphi^{s_0} \cap U^\omega) = 1$. □

Proposition 11. Let $\mathcal{M} = (S, S_0, S_\omega, \to, P)$ be an MDP, and $s_0 \in S$, and $\sigma$ a strategy, and $\operatorname{Col} : S \to \{ 1, 2 \}$, and $\varphi = \operatorname{Parity}(\operatorname{Col})$. Suppose $P_{M,s_0,\sigma}(\varphi) = 1$. Then there is an MD-strategy $\sigma'$ with $P_{M,s_0,\sigma'}(\varphi) = 1$.

Proof. We can assume that all states are almost-surely winning, since in order to achieve an almost-sure winning objective, the player must forever remain in almost-surely winning states. We provide an MD-strategy $\hat{\sigma}$ such that for all states $s \in S$ we have $P_{M,s,\hat{\sigma}}(\varphi) = 1$. □
Set $\varphi' = \text{Reach}(\{S\}^{\text{Col}=2})$. Note that $[\varphi] \subseteq [\varphi']$. Since all states are almost-surely winning for $\varphi$, all states are almost-surely winning for $\varphi'$. By Theorem 10 and Lemma 7 there is an MD-strategy $\hat{\sigma}$ such that for all states $s \in S$ we have $P_{M,s,\hat{\sigma}}(\varphi') = 1$. That is, $\hat{\sigma}$ reaches the set $[S]^{\text{Col}=2}$ with probability 1, regardless of the start state. It follows that it reaches, with probability 1, the set $[S]^{\text{Col}=2}$ infinitely often. Hence $P_{M,s,\hat{\sigma}}(\varphi) = 1$ holds for all $s \in S$. \hfill \Box

D. Proofs of Subsection VI-A

**Proposition 13.** Let $M = (S, S_{\mathcal{D}}, S_C, \rightarrow, P)$ be a finitely branching MDP, and $T \subseteq S$, and $\varphi = \text{Safety}(T)$. Define an MD-strategy $\sigma_{\text{opt-av}}$ (for “optimal avoiding”) that, in each state $s$, picks a successor state with the largest value $\text{val}(s) = \sup_{s' \in S} P_{M,s,s'}(\varphi)$. Then for all states $s \in S$ we have $P_{M,s,\sigma_{\text{opt-av}}}(\varphi) = \text{val}(s)$, i.e., $\sigma_{\text{opt-av}}$ is uniformly optimal.

**Proof.** We can assume that $T$ is a sink. Fix a state $s_0$. Write $s_0 s_1 s_2 \cdots \in sS^\omega$ for a random run, i.e., $s_1, s_2, \ldots$ denote random states. For any $n \in \mathbb{N}$ let $[X^n-T] : s_0 S^\omega \rightarrow \{0, 1\}$ be the random variable that indicates if $s_n \notin T$. Note that $[X^n-T] \geq \text{val}(s_n)$. Writing $E_{M,s_0,\sigma_{\text{opt-av}}}$ for the expectation with respect to $P_{M,s_0,\sigma_{\text{opt-av}}}$, we have:

$$P_{M,s_0,\sigma_{\text{opt-av}}}(\varphi) = P_{M,s_0,\sigma_{\text{opt-av}}} \left( \bigcap_{i=0}^{\infty} [X^i-T]^{s_0} \right) \quad \text{semantics of } \varphi$$

$$= \lim_{n \rightarrow \infty} P_{M,s_0,\sigma_{\text{opt-av}}} \left( \bigcap_{i=0}^{n} [X^i-T]^{s_0} \right) \quad \text{measures are continuous from above}$$

$$= \lim_{n \rightarrow \infty} P_{M,s_0,\sigma_{\text{opt-av}}}(X^n-T) \quad \text{T is a sink}$$

$$= \lim_{n \rightarrow \infty} E_{M,s_0,\sigma_{\text{opt-av}}}(X^n-T) \quad \text{definition of } [X^n-T]$$

$$\geq \lim\inf_{n \rightarrow \infty} E_{M,s_0,\sigma_{\text{opt-av}}}(\text{val}(s_n)) = [X^n-T] \geq \text{val}(s_n)$$

A straightforward induction on $n$, using the definition of $\sigma_{\text{opt-av}}$, shows that $E_{M,s_0,\sigma_{\text{opt-av}}}(\text{val}(s_n)) = \text{val}(s_n)$ holds for all $n \in \mathbb{N}$. Hence $P_{M,s_0,\sigma_{\text{opt-av}}}(\varphi) \geq \text{val}(s_0)$. The converse inequality holds by the definition of the value. Hence $P_{M,s_0,\sigma_{\text{opt-av}}}(\varphi) = \text{val}(s_0)$.

**Lemma 15.** Let $M = (S, S_{\mathcal{D}}, S_C, \rightarrow, P)$ be a finitely branching MDP, and $\text{Col} : S \rightarrow \mathbb{N}$ a color function. Let $s$ be a state, and $\sigma$ a strategy, and $\tau < 1$. Then $P_{M,s,\sigma}(\text{FG-Safe}(\tau) \land \text{FG}[S]^{\text{Col}=0}) = 0$.

**Proof.** For any $n \in \mathbb{N}$ define $Z_n = ([S]^{\text{Col}=0})^n$. That is, $Z_n S^\omega$ is the event that the first $n$ visited states have color 0. For any state $s \notin \text{Safe}(\tau)$, let $n(s) \in \mathbb{N}$ be the smallest number such that $P_{M,s,\sigma}(Z_{n(s)} S^\omega) \leq (1 + \tau)/2$. This is well-defined.

Let $L \subseteq S^*$ be the set of finite sequences $s_0 s_1 \cdots s_{n-1}$ such that $s_0 \notin \text{Safe}(\tau)$ and $n = n(s_0)$ and $\forall i < n, s_i \in [S]^{\text{Col}=0} \setminus \text{Safe}(\tau)$. We show for all $s \in S$, $\text{Safe}(\tau)$ and all $k \in \mathbb{N}$ that $P_{M,s,\sigma}(L^k S^\omega) \leq \left(\frac{1 + \tau}{2}\right)^k$. We proceed by induction on $k$. The case $k = 0$ is trivial. For the induction step, let $k \geq 0$. We have:

$$P_{M,s,\sigma}(L^{k+1} S^\omega) \leq P_{M,s,\sigma}(Z_{n(s)} L^k S^\omega) \leq P_{M,s,\sigma}(Z_{n(s)} S^\omega) \cdot \sup_{s' \in S \setminus \text{Safe}(\tau)} P_{M,s',\sigma}(L^k S^\omega) \leq P_{M,s,\sigma}(Z_{n(s)} S^\omega) \cdot \left(\frac{1 + \tau}{2}\right)^k$$

induction hypothesis

$$\leq \left(\frac{1 + \tau}{2}\right)^{k+1}$$

definition of $n(s)$

This completes the induction proof. Write $\varphi := \text{G-Safe}(\tau) \land \text{G}[S]^{\text{Col}=0}$. We have for all $s \in S$:

$$P_{M,s,\sigma}(\varphi) = P_{M,s,\sigma}(L^\omega) = \lim_{k \rightarrow \infty} P_{M,s,\sigma}(L^k S^\omega) \leq \lim_{k \rightarrow \infty} \left(\frac{1 + \tau}{2}\right)^k = 0$$

as shown above

$$\tau < 1$$

20
It follows: 
\[ P_{M,s,\sigma}(X^j \varphi) = 0 \] for all \( s \in S \) and all \( j \in \mathbb{N} \)

Thus we have:
\[
\begin{align*}
P_{M,s,\sigma}(FG\neg Safe(\tau) \land FG[S]^{Col=0}) & = P_{M,s,\sigma}(F \varphi) \\
& = P_{M,s,\sigma}(\bigcup_{j \in \mathbb{N}} X^j \varphi) \\
& \leq \sum_{j \in \mathbb{N}} P_{M,s,\sigma}(X^j \varphi) \\
& = 0
\end{align*}
\]