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On the Performance of Multi-Stream Receive-Spatial Modulation in the MIMO Broadcast Channel

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Abstract—In this paper, a novel architecture for the Multiple-Input Multiple-Output (MIMO) broadcast channel is proposed and studied. The new architecture is based on the concept of Multi-Stream Receive-Spatial Modulation (MSR-SM). MSR-SM is a closed-loop transmission scheme, which applies the concept of multi-stream space modulation at the receiver side. A new and accurate framework for computing the Average Bit Error Probability (ABEP) of the new architecture is proposed. In addition, the new architecture is compared against the state-of-the-art MIMO transmission in the broadcast channel and it is shown to: i) provide superior Bit Error Rate (BER) performance in the high Signal-to-Noise-Ratio (SNR) regime and ii) reduce the signal processing complexity at the transmitter.

I. INTRODUCTION

A Multiple-Input Multiple-Output (MIMO) communication scheme that has the potential to provide a low complexity system implementation is Spatial Modulation (SM) [1–7]. The operating principle of SM is designed in such a way that the transmitter requires only one Radio Frequency (RF) chain. Therefore, significant Energy Efficiency (EE) and complexity gains are obtained in comparison to conventional MIMO techniques [8]. Due to the potential of SM, several variants of SM have been published [9–11]. A comprehensive overview of the existing literature on SM is given in [6].

As demonstrated in [6], the study of SM and its variants in point-to-point communication is extensive. In addition, the formation of a Multi-User (MU) system based on the concept of SM can be undertaken with the aid of a multiple access scheme, such as Time Division Multiple Access (TDMA), Frequency Division Multiple Access (FDMA), or Orthogonal Frequency-Division Multiple Access (OFDMA). However, a new trend in wireless communication promotes the aggressive allocation of multiple users in the same time and frequency resources. In such systems, usually, the inherent interference is eliminated or mitigated via the deployment of Space Division Multiple Access (SDMA) techniques. In particular, the Multiple Access Channel (MAC) is formed in the uplink [12, 13] and the broadcast channel is established in the downlink [14].

Relevant to this field, [15–18] extend the concept of SM in a MU setup. In particular, the schemes proposed in [15, 16] are applicable to the uplink, whereas the scheme in [17, 18] is suitable for the downlink. Unfortunately, in [15, 16], either there is an increase in the the complexity of the receiver, or there is an error saturation in the Bit Error Rate (BER) performance. In addition, the scheme in [17, 18] has a BER performance degradation with respect to the Single-User (SU) communication. Note that the incorporation of SM in the broadcast channel is a challenging task since the deployment of interference elimination or reduction techniques is difficult. This challenge originates from the way in which information is transmitted in SM. Thus, the design of a SM-based architecture for the broadcast channel becomes a challenging and important task.

Against this background and based on the concept of Multi-Stream Receive-Spatial Modulation (MSR-SM), this paper proposes and studies a new MU architecture for the downlink. MSR-SM is a closed loop and point-to-point modulation scheme which applies the concept of Multi-Stream Spatial Modulation (MS-SM) at the receiver side [19–24]. In more detail, this paper incorporates MSR-SM in the MIMO broadcast channel. This is achieved using Zero Forcing (ZF) forcing precoding.

It is demonstrated that the new architectures outperforms the corresponding conventional spatially multiplexed MIMO broadcast channel, in terms of BER, in high Signal-to-Noise Ratio (SNR). In addition, this paper derives tight upper bounds for the theoretical Average Bit Error Probability (ABEP) of a typical user when: i) the wireless channel follows a Rayleigh distribution and ii) Perfect-Channel State Information at the Transmitter (P-CSIT) is available. This is undertaken by considering an accurate statistical framework for the received signal. The theoretical ABEP of point-to-point MSR-SM has been recently studied in [19, 20]. However, it is emphasized that the mathematical framework derived in this paper exhibits major differences with respect to those presented in [19, 20]. In [19], the statistical description of the received signal is not taken into account. Furthermore, the ABEP computed in [20] is derived for a scenario where a suboptimal detector, which decouples the detection process, is deployed. In contrast, the analysis given in this paper is different for the following reasons: i) multiple users are considered; ii) the statistical description of the received signal of a typical user is considered; and iii) the detection process is based on the Maximum Likelihood (ML) principle, which, as shown in Section IV, imposes some additional mathematical difficulty. Note that the new framework is directly applicable to a point-to-point single
user scenario.

The rest of the paper is organized as follows: The system model of the new architecture is introduced in Section II. The computational complexity of the new architecture is discussed in Section III. The theoretical analysis of the ABEP of a typical user is presented in Section IV. The proposed MU MSR-SM architecture is compared against the corresponding benchmark system in Section V. In addition, numerical results that validate the new theoretical findings are presented in the same section. Finally, the concluding comments of this paper are given in Section VI.

Notation: In the following, lowercase bold letters denote vectors and uppercase bold letters denote matrices. \((\cdot)^T\), \((\cdot)^H\), \(\text{tr}(\cdot)\) and \(A^{1/2}\) denote transpose, Hermitian transpose, matrix trace and the square root of \(A\), respectively. \(\|\cdot\|_2\) represents the Euclidean norm. \(\text{diag}(a_1, \ldots, a_n)\) represents a diagonal matrix whose main diagonal includes the elements \(a_1, \ldots, a_n\). \(E[\cdot]\) is the mean value of a RV. A complex Gaussian distribution with mean \(m\) and variance \(\sigma^2\) is represented as \(\mathcal{CN}(m, \sigma^2)\), where its real and imaginary part are independent and identically distributed (i.i.d.) Gaussian RV with distribution \(\mathcal{N}(m, \sigma^2_\Re)\). \(\Re\{\cdot\}\) denotes the real part of a complex number or matrix.

II. SYSTEM MODEL

In this section, a single cell uncoded downlink transmission is considered. More specifically, a Base Station (BS) with \(N_t\) antennas aims to establish the MIMO broadcast channel in order to serve \(N_u\) users. Each user is equipped with \(N_r\) antennas. Since the transmitter is a BS, the assumption that \(N_t \geq N_u N_r\) is realistic. However, if this assumption is not satisfied, user scheduling can be applied in order to overcome this constraint. Furthermore, the MIMO wireless channel between the transmitter and the remote users is assumed to be frequency flat and quasi-static. Finally, in this paper, a scenario where the transmitter possesses P-CSIT, using either the channel reciprocity or fast and error free links from the transmitter to the remote users is assumed to be satisfied, user scheduling can be applied in order to overcome this constraint. Furthermore, ZF precoding is justified by its low complexity and its ability to totally eliminate the interference between different users and between different antennas of the same user. Therefore, ZF precoding can be considered as an efficient method for the formation of a MU architecture based on MSR-SM.

Given that the ZF precoder is the pseudo-inverse of the channel matrix \(H\), its matrix form is written as:

\[
P = H^H (HH^H)^{-1}.
\]

(10)

In this case, the \(i\)-th element of the main diagonal of \(D\) equals to:

\[
d_i = \sqrt{\frac{1}{\|p_i\|_2^2}}, \quad i = 1, \ldots, N_u N_r.
\]

(11)

The incorporation of (10) in (1) gives:

\[
y = Dx + w.
\]

(12)

In addition, the received signal at the \(i\)-th is given as:

\[
y_i = D_i x_i + w_i, \quad i = 1, \ldots, N_u.
\]

(13)
In (13), $D_i$, $i = 1, \ldots, N_u$, is the $N_r \times N_r$ diagonal normalization matrix of the corresponding precoding matrix $P_i$. In particular, the normalization matrix $D$ can be interpreted as the following block diagonal matrix, $D = \text{diag}(D_1, \ldots, D_{N_u})$.

From (13), it can be inferred that the structure of $x_i$, $i = 1, \ldots, N_u$, determines the way that information is transmitted to each user. For example, conventional Spatial Multiple Xing (SMX) MIMO transmission is established if all of the elements of $x_i$ are drawn from a conventional $M$-ary constellation diagram $\mathcal{M}$.

In order to establish a MSR-SM transmission mechanism between the transmitter and the $i$-th user, every symbol period, the following two requirements have to be fulfilled. Firstly, $N_s \leq N_r$ conventional symbols have to be conveyed from the transmitter to the $i$-th user per symbol period. In this way, only a subset of $N_r$ antennas at the $i$-th user receive a non-zero signal. All of the other antennas face only thermal noise. And secondly, additional binary information has to be encoded via the indices of the $N_s$ (out of $N_r$) receiving antennas. Note that when $N_s = N_r$, MSR-SM reduces to conventional SMX transmission and no additional information is conveyed via the indices of the receiving antennas.

Provided that the deployed precoding method is ZF, the received signal at each user is given in (13). Hence, via the appropriate selection of the structure of $x_i$, the transmitter is able to impose that the noise-free received signal $D_i x_i$ has exactly $N_s$ non-zero elements and $N_r - N_s$ zero elements. Therefore, the non-zero elements of $D_i x_i$ is a scaled version of the corresponding non-zero elements of $x_i$. Also, the positions of the zero elements of $D_i x_i$ are the same as the positions of the zero elements of $x_i$. Thus, a number of bits can be encoded on the positions of the non-zero elements of $x_i$.

The structure of the transmission alphabet of MSR-SM, $B_i$, is given in (14) at the top of this page. As shown in (14), $x_i \in B_i$ has exactly $N_s$ non-zero elements which belong to a conventional constellation, $\{s_1, \ldots, s_{N_r}\} \in \mathcal{M}$. Here, $\mathcal{M}$ denotes the deployed constellation. The positions of the non-zero elements of $x_i$ correspond to the indices of the signal receiving antennas, while the zero elements of $x_i$ correspond to the indices of the non-receiving antennas. Given that the length of $x_i$ is $N_r$ and there are $N_r$ non-zero elements, the number of total combinations of $N_r$ non-zero elements (receiving antennas) out of $N_r$ is $\binom{N_r}{N_s}$. Here, $\binom{\cdot}{\cdot}$ denotes the binomial coefficient. However, only $k_c = 2^{\log_2(\binom{N_r}{N_s})}$ combinations are used in order to encode $k^{\log_2(\binom{N_r}{N_s})}$ bits. Also, here, $\lfloor \cdot \rfloor$ denotes the floor function which corresponds a real number to its largest previous integer number. This is done by assigning a unique binary index of length of $k^{\log_2(\binom{N_r}{N_s})}$ bits to each one of the $k_c$ used (legal) combinations. The selection of these (legal) combinations of receiving antennas can be done intelligently in order to minimize the instantaneous BER or it can be done randomly. In the first case, the system complexity is increased and the selection is undertaken adaptively. Given that MSR-SM is a closed loop scheme, the latter case is expected to offer a good performance with no further complexity overhead. In this paper, the focus is on the latter case.

In this way, every symbol period, the bit-stream to be transmitted is divided in two parts. The first part of $k^{\log_2(\binom{N_r}{N_s})}$ bits, is encoded and transmitted via $N_s$ symbols drawn from the $M$-ary constellation $\mathcal{M}$. The second part of $k^{\log_2(\binom{N_r}{N_s})}$ bits is encoded on the combination of the signal receiving antennas. Therefore, the spectral efficiency of MSR-SM is:

$$k_{\text{MSR-SM}} = N_s \log_2(M) + \log_2\left(\binom{N_r}{N_s}\right)$$

bits per channel use (bpcu) per user.

Assuming that the $i$-th user is aware of $D_i$, the reconstruction of the transmitted bit-streams is conducted by detecting the transmitted vectors $x_i$, $i = 1, \ldots, N_u$. This is undertaken at every user independently by using the following ML detector:

$$\hat{x}_i = \arg \min_{x_i} \|y_i - D_i x_i\|^2_2, \quad i = 1, \ldots, N_u.$$

III. COMPUTATION COMPLEXITY ANALYSIS

The observation of (14) shows that the sparse structure of the transmission alphabet of MSR-SM, $B_i$, can be deployed in order to decrease the computational complexity of the transmitter. Assuming that the ZF precoder in (10) is precomputed before the transmission of each block of symbols, the transmitted signal, $s = PDx$, in (1) can be calculated with:

$$C_t = N_r(8N_uN_s - 2) + 2N_uN_s,$$

real operations (additions or multiplications). Hence, lower values of $N_r$ result in lower computational complexity at the transmitter. Note that when it holds that $N_s = N_r$, MSR-SM is transformed into a conventional SMX transmission. In this way, it can be inferred that, at the transmitter side, the computational complexity of MSR-SM is less than the corresponding complexity of conventional SMX transmission.

At the users’ side, the ML detection process of the conventional MIMO broadcast channel is decoupled in per single stream detection. Therefore, its detection complexity is lower compared with the corresponding complexity of MU MSR-SM. However, in [20] a suboptimal detector for MSR-SM is proposed. This detector achieves almost the same computational complexity as the previous decoupled detector. The study of the detector of [20] is, however, outside of the scope of this paper.

$$x_i = \begin{bmatrix} 0, \ldots, 0, s_{i_1} \overline{i_1}-\text{th position}, 0, \ldots, 0, s_{i_2} \overline{i_2}-\text{th position}, \ldots, 0, \ldots, 0, \overline{i_N_u-}\text{th position} \end{bmatrix}_T$$ (14)
IV. Theoretical Evaluation of the Average Bit Error Probability

In this section, the ABEP of a typical user is derived. Based on the union bound technique, the ABEP of the $i$-th user, $P_{\text{bit}}^i(\gamma)$, for a given transmit SNR $\gamma$, is bounded as:

$$P_{\text{bit}}^i(\gamma) \leq \frac{1}{|B_i|_{\text{MSR-SM}}} \sum_{x_i} \sum_{x_i \neq \hat{x}_i} d(x_i \rightarrow \hat{x}_i) P_e^i(x_i \rightarrow \hat{x}_i, \gamma).$$

In (18), $P_e^i(x_i \rightarrow \hat{x}_i, \gamma)$ denotes the Pairwise Error Probability (PEP) of transmitting $x_i$ to the $i$-th user while its detector erroneously decides in favor of $\hat{x}_i$. The Hamming distance between the bit-words represented by $x_i$ and $\hat{x}_i$ is denoted as $d(x_i \rightarrow \hat{x}_i)$. Also, the size of the transmission alphabet of MSR-SM to the $i$-th user is given as $|B_i| = M^{N_2} 2^{L|\log_2 (N_s')|}$.

The evaluation of (18) requires the knowledge of $P_e^i(x_i \rightarrow \hat{x}_i, \gamma)$, which is the expectation of the instantaneous PEP over all channel realizations. The instantaneous PEP of the $i$-th user is expressed as:

$$P_e^i(x_i \rightarrow \hat{x}_i, \gamma|D^i_\gamma) = Q\left(\sqrt{\frac{c_i^H D^i_\gamma c_i}{2}} \gamma\right) = Q\left(\sqrt{\frac{z_i \gamma}{2}}\right),$$

where, $c_i = x_i - \hat{x}_i$; $z_i = e_i^H D^i_\gamma e_i$; and $\gamma = 1/\sigma_w^2$ is the transmit SNR. The proof of (19) relies on the fact that a symbol error occurs at the $i$-th user when, $E_i(x_i, \hat{x}_i) = \{||y_i - D_i x_i||_2^2 > ||y_i - D_i \hat{x}_i||_2^2\}$. It is not difficult to show that, $E_i(x_i, \hat{x}_i) = \{\text{Re} \{e_i^H D_i w_i\} > e_i^H D^i_\gamma e\}$. Therefore, using the statistical distribution of $w_i$, the PEP of the $i$-th user is expressed as in (19).

The observation of the instantaneous PEP of the $i$-th user in (19) shows that it is conditioned on the Random Variable (RV) $z_i$. Thus, the Probability Density Function (PDF) of $z_i$ has to be derived. However, it holds that:

$$z_i = \sum_{k=1}^{N_r} |x_k - \hat{x}_k|^2 d_k^i = \sum_{x_k - \hat{x}_k \neq 0} |x_k - \hat{x}_k|^2 d_k^i.$$  \hspace{1cm} (20)

In [25], it is explicitly shown that $d_k$ is a gamma RV with $d_k^i \sim \Gamma(M_{\text{MU}}, 1)$, where, $M_{\text{MU}} = N_i - N_N N_s + 1$. Therefore, the RV $Z_k = |x_k - \hat{x}_k|^2 d_k^i$ is also a gamma RV, when it holds that $x_k - \hat{x}_k \neq 0$. Here, for notational convenience, the following variable is introduced, $b_k = |x_k - \hat{x}_k|^2$. Based on the previous arguments, the PDF of $Z_k$ is given as:

$$f_{Z_k}(x) = \frac{1}{b_k^{L+1}} x^{L-1} e^{-\frac{x}{b_k}} H_0(x),$$

where, $H_0(x)$ is the Heaviside step function, for which it holds that: $H_0(x) = 0$ for $x < 0$ and $H_0(x) = 1$ for $x \geq 0$.

Usually, in published research, the RVs $d_k^i$ are assumed to be statistically independent in order to simplify the whole analysis [26]. This assumption is in contradiction with the structure of $d_k^i = 1/\left(\{H, H^T\}^{-1}\right)_{k,k}$, since the realization of every RV $d_k^i$ occurs using the same mathematical operation on the same random matrix $H_i$.

In this paper, the statistical dependence of the RVs $d_k^i$ is considered. For this reason, $z_i$ is the result of the sum of correlated gamma RVs. Thus, using the result from [27, Corollary 1] and (20), the PDF of $z_i$ is directly expressed as:

$$f_{z_i}(x) = \prod_{l=1}^{N_i} \left(\frac{\alpha_l}{\alpha_l^2}\right)^L \sum_{k=0}^{x} \frac{\delta_{k,x}^N N_{L+1-k} e^{-x}}{\alpha_l N_{L+k+1}^N k} H_0(x),$$

where, $N_i$ is the number of non zero elements of $c_i$ for a given pair of $x_i$ and $\hat{x}_i$. In (22), $\alpha_l, l = 1, \ldots, N_i$, denote the eigenvalues of $A = BR$ in ascending order. Here, $B$ is defined as the following diagonal matrix, $B = \text{diag} (b_1, \ldots, b_{N_i})$, where, $b_1, l = 1, \ldots, N_i$, is the absolute value of the $l$-th non zero element of $c_i$. Also, $R$ is a $N_i \times N_i$ matrix defined as:

$$R = \begin{bmatrix} 1 & \sqrt{\rho_c} & \cdots & \sqrt{\rho_c} \\ \sqrt{\rho_c} & \cdots & \cdots & \cdots \\ \vdots & \cdots & \cdots & \cdots \\ \sqrt{\rho_c} & \cdots & \cdots & 1 \end{bmatrix},$$

where, $\rho_c$ is the Pearson product-moment correlation coefficient between any pair of two different RVs of the main diagonal of $D^i_\gamma$. Finally, $\delta_k, k = 0, 1, 2, \ldots$, are given in (24) at the top of this page.

In order to further simplify (19), the Chernoff bound of the $Q$-function, $Q(x) \leq \frac{1}{2} e^{-\frac{x}{2}}$, is considered. In this way, the PEP of the $i$-th user is expressed as:

$$P_e^i(x_i \rightarrow \hat{x}_i, \gamma_i) \leq \frac{1}{2} E_{z_i} \left[ e^{-\frac{z_i \gamma_i}{2}} \right] = \frac{1}{2} \int_{-\infty}^{+\infty} e^{-\frac{y \gamma_i}{2}} f_{z_i}(y) dy.$$  \hspace{1cm} (25)

Provided that the PDF of $z_i$ is given from (22), after some algebraic manipulations which are omitted here due to space limitation, the evaluation of (25) gives:

$$P_e^i(x_i \rightarrow \hat{x}_i, \gamma_i) \leq \frac{\prod_{l=1}^{N_i} \left(\frac{\alpha_l}{\alpha_l^2}\right)^L_{\text{MU}} \left(\frac{\alpha_l}{4} \xi_{\gamma} + 1\right)^{N_{\text{MU}}}}{2} \times \sum_{k=0}^{+\infty} \delta_k \left(\frac{\alpha_l}{4} \xi_{\gamma} + 1\right)^{-k}.$$  \hspace{1cm} (26)

Therefore, the computation of the ABEP of the $i$-th user is conducted from (18), by using (26). Note that the derived bounds are directly applicable to the conventional MIMO broadcast channel when $N_s = N_s$. 

\begin{equation}
\delta_{k+1} = \begin{cases} 1, & k \neq -1 \\ \frac{k}{k+1} \sum_{i=1}^{k+1} \left[ \sum_{j=1}^{N} \left(1 - \frac{\alpha_j}{\alpha_i}\right)^j \right] \delta_{k+1-i}, & k = 0, 1, 2, \ldots
\end{cases}
\end{equation}
Empirical and Analytical PDF of $z_i$

**Fig. 1.** Illustration of empirical and analytical PDF of (20) by assuming that: i) the RVs $d_2^k$, $k = 1, \ldots, 2$ are statistically dependent and ii) they are independent. Setup: $H \sim \mathcal{C}\mathcal{N}(0_{2 \times 4}, I_{2 \times 4})$; and ii) $b_1 = 0.5$ and $b_1 = 1.2$, where, $b_i$, $i = 1, 2$, are parameters defined in (21).

**V. Simulation Results and Discussion**

The main objectives of this section are: i) the validation of the theoretical results of Section IV and ii) the comparison of the new MU architecture with the corresponding State-of-the-Art (SotA) benchmark scheme. In this section, the wireless channel is assumed to follow a Rayleigh distribution. Section IV provides the derivation of the ABEP of MU MSR-SM using the PDF of $z_i$ given in (22). The deployment of (22) relies on the fact that the RVs $d_2^k$, $k = 1, \ldots, N_r$, are statistically dependent. In order to confirm that the RVs $d_2^k$ are statistical dependent, Fig. 1 depicts the empirical PDF of (20) against its analytical form as given in (22). In addition, Fig. 1 presents the analytical PDF of (22) under the assumption that $d_2^k$, $k = 1, \ldots, N_r$, are statistically independent RVs. If this assumption was valid, the PDF of (20) could be directly obtained using the result from [27, Theorem 1]. From Fig. 1, it can be concluded that the theoretical PDF of (22) perfectly matches its empirical PDF. In contrast, when the RVs $d_2^k$, $k = 1, \ldots, N_r$, are assumed to be independent, the obtained PDF from [27, Theorem 1] deviates from the empirical results.

As a form of validation of the bounds derived in Section IV, in Fig. 2, these bounds are compared against the corresponding BER curves obtained via Monte Carlo simulations. The observation of Fig. 2 shows that the upper bounds are tight in high SNR regime. In low SNR, there is a gap between the theoretical and simulated curves. However, this gap is a well known effect of the union bound technique deployed in Section IV [28].

In Fig. 3, the BER performance of MU MSR-SM as a function of $N_s$. Setup: $N_t = 20$, $N_r = 4$, $N_u = 4$. The spectral efficiency is 8 bpcu.

In Fig. 4, the BER performance of MSR-SM versus the conventional MIMO broadcast channel. Setup: $N_t = 20$, $N_r = 4$, and $N_u = 4$. matches its empirical PDF. In contrast, when the RVs $d_2^k$, $k = 1, \ldots, N_r$, are assumed to be independent, the obtained PDF from [27, Theorem 1] deviates from the empirical results.

As a form of validation of the bounds derived in Section IV, in Fig. 2, these bounds are compared against the corresponding BER curves obtained via Monte Carlo simulations. The observation of Fig. 2 shows that the upper bounds are tight in high SNR regime. In low SNR, there is a gap between the theoretical and simulated curves. However, this gap is a well known effect of the union bound technique deployed in Section IV [28].

In Fig. 3, the BER performance of MU MSR-SM as a function of $N_s$ is shown for the same spectral efficiency. In particular, Fig. 3 demonstrates that the BER of a typical user improves as $N_s$ increases. This happens because higher values of $N_s$ require a lower modulation order of Quadrature Amplitude Modulation (QAM) in order to achieve the same spectral efficiency. However, as shown in Fig 4, the optimal value of $N_s$ is not to $N_r$, i.e., the conventional MIMO broadcast channel.
A comparison between the BER performance of MU MSR-SM and the corresponding benchmark system is presented in Fig. 4. In this paper, the selected benchmark system is the conventional spatially multiplexed MIMO broadcast channel, where $N_r$ symbol streams are established per user. The spectral efficiency of both schemes is set to be the same by selecting the appropriate constellation order. In addition, MU MSR-SM spatially modulates $N_s = 2$ and $N_t = 3$ parallel symbol streams in order to achieve 4 and 8 bpcu, respectively. As shown in Fig. 4, in low SNR the conventional MIMO broadcast channel offers a slightly better performance than MU MSR-SM. However, for practical values of BER approximately less than $10^{-2}$, MU MSR-SM outperforms the benchmark systems. The performance gap between the considered schemes is about 0.5 dB and 1 dB for 4 and 8 bpcu, respectively. Furthermore, Fig. 4 indicates that MU MSR-SM achieves the same diversity order as the benchmark system. The performance gap between the considered schemes is about 0.5 dB and 1 dB for 4 and 8 bpcu, respectively. Therefore, MU MSR-SM could be considered as an alternative transmission mechanism for the MIMO broadcast channel. It was shown that the new architecture outperforms the benchmark system in high SNR regime. Also, it offers lower computational complexity in the transmitter. Therefore, MSR-SM could be considered as an alternative transmission mechanism for the MIMO broadcast channel.

Finally, the derivation of the diversity order and coding of MU MSR-SM, as well as its theoretical comparison with the SotA is a subject of our future work.

VI. CONCLUSIONS

The incorporation of MSR-SM in the MIMO broadcast channel was introduced and studied in this paper. In particular, a novel upper bound for the theoretical BER of the new architecture was derived. This was undertaken by using a new and accurate statistical framework for the received signal of a typical user. It was shown that the new bounds are tight. For the purpose of comparison, MU MSR-SM was evaluated, in terms of BER, against the corresponding conventional MIMO broadcast channel. It was shown that the new architecture outperforms the benchmark system in high SNR regime. Also, it offers lower computational complexity in the transmitter. Therefore, MSR-SM could be considered as an alternative transmission mechanism for the MIMO broadcast channel.

Finally, the derivation of the diversity order and coding of MU MSR-SM, as well as its theoretical comparison with the SotA is a subject of our future work.

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