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BEAS: Bounded Evaluation of SQL Queries

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Abstract
We demonstrate BEAS, a prototype system for querying relations with bounded resources. BEAS advocates an unconventional query evaluation paradigm under an access schema $\mathcal{A}$, which is a combination of cardinality constraints and associated indices. Given an SQL query $Q$ and a dataset $D$, BEAS computes $Q(D)$ by accessing a bounded fraction $D_Q$ of $D$, such that $Q(D_Q) = Q(D)$ and $D_Q$ is determined by $\mathcal{A}$ and $Q$ only, no matter how big $D$ grows. It identifies $D_Q$ by reasoning about the cardinality constraints of $\mathcal{A}$, and fetches $D_Q$ using the indices of $\mathcal{A}$. We demonstrate the feasibility of bounded evaluation by walking through each functional component of BEAS. As a proof of concept, we demonstrate how BEAS conducts CDR analyses in telecommunication industry, compared with commercial database systems.

1. Introduction
Querying big relations is often beyond reach for small companies. It may take hours to join tables of millions of tuples. Given a query $Q$ and a dataset $D$, it is NP-complete to decide whether a tuple is in $Q(D)$ when $Q$ is in PSPACE (selection, projection, Cartesian product). It is PSPACE-complete for $Q$ in relational algebra [1]. One might think that parallelism would solve the problem by using more processors. However, small companies can often afford only limited resources.

Is it feasible to query big data with bounded resources? One approach to addressing this challenge is based on bounded evaluation [5–9]. The idea is to use an access schema $\mathcal{A}$ over a database schema $\mathcal{R}$, which is a combination of cardinality constraints and associated indices. A query $Q$ is boundedly evaluable under $\mathcal{A}$ if for each instance $D$ of $\mathcal{R}$ that conforms to $\mathcal{A}$, there exists a small $D_Q \subseteq D$ such that:

- $Q(D_Q) = Q(D)$, and
- the time for identifying $D_Q$ is decided by $Q$ and $\mathcal{A}$.

Intuitively, $D_Q$ consists of only data needed for answering $Q$. Its size $|D_Q|$ is determined by $Q$ and $\mathcal{A}$ only, not by $|D|$.

BEAS. As a proof of concept of [5–9], we have developed BEAS [4], a prototype system for Bounded Evaluation of SQL. BEAS has the following unique features that differ from conventional query evaluation paradigm and DBMS.

1. Quantified data access. BEAS identifies $D_Q$ by reasoning about the cardinality constraints in $\mathcal{A}$, and fetches $D_Q$ by using the indices in $\mathcal{A}$. In the process it deduces a bound $M$ on the amount of data to be accessed, and can hence decide whether $Q$ is boundedly evaluable before $Q$ is executed.

2. Reduced redundancy. Using $\mathcal{A}$, BEAS fetches only (distinct) partial tuples needed for answering $Q$. This reduces duplicated and unnecessary attributes in tuples fetched by traditional DBMS. It also reduces joins, in which the redundancies get inflated rapidly (see an example shortly).

3. Scalability. Putting these together, BEAS computes $Q(D)$ by accessing a bounded fraction $D_Q$ of $D$, no matter how big $D$ grows. Hence to an extent, it makes big data analysis possible for small businesses with bounded resources.

4. Ease of use. BEAS can be built on top of any conventional DBMS, and make seamless use of existing optimization techniques of DBMS. This makes it easy to extend DBMS with the functionality of bounded evaluation.

One of our industry collaborators has deployed and tested BEAS, and found that BEAS outperforms commercial DBMS by orders of magnitude for more than 90% of their queries.

Demo overview. We demonstrated the bounded evaluation functionality of BEAS in two parts. (1) To illustrate how bounded evaluation works, we walk through each functional component of BEAS, from access schema discovery and maintenance to bounded query plan generation and execution. (2) To demonstrate the performance of BEAS, we adopt a real-life scenario from telecommunication industry for CDR (call detail record) analyses, and visualize how different query plans perform compared with commercial DBMS.

Below we first present the foundation (Section 2) and the functional components (Section 3) of BEAS. We then propose a more detailed demonstration plan (Section 4).

2. Foundations of BEAS
We start with a review of access schema and bounded evaluability [5, 8, 9], which are the foundations of BEAS.

Access schema. Over a database schema $\mathcal{R}$, an access constraint $\psi$ is of the form $R(X \rightarrow Y, N)$, where $R$ is a relation in $\mathcal{R}$, $X$, $Y$ are sets of attributes of $R$, and $N$ is a natural number [5, 8, 9]. A relation instance $D$ of $\mathcal{R}$ conforms to $\psi$ if:

- for any $X$-value $a$ in $D$, $|D_Y(X = a)| \leq N$, where $D_Y(X = a) = \{(t[Y]) | t \in D, t[X] = a\}$; and
- there exists an index on $X$ for $Y$ that given an $X$-value $a$, retrieves $D_Y(X = a)$ by accessing at most $N$ tuples.

That is, for any given $X$-value, there exist at most $N$ distinct corresponding $Y$ values in $D$ (cardinality constraint), and the $Y$ values can be fetched by using the index for $X$ (index).

An access schema $\mathcal{A}$ over $\mathcal{R}$ is a set of access constraints over $\mathcal{R}$. A database instance $D$ of $\mathcal{R}$ conforms to $\mathcal{A}$, denoted by $D \models \mathcal{A}$, if $D$ conforms to each constraint in $\mathcal{A}$.

Example 1: Consider a commercial benchmark of schema $\mathcal{R}_0$ from a telecommunication company (name withheld). It includes three (simplified) relations: (a) call($\text{pnum}$, $\text{recnum}$, $\text{date}$, $\text{region}$), recording that number $\text{pnum}$ called $\text{recnum}$ in region on date; (b) package($\text{pnum}$, $\text{pid}$, $\text{start}$, $\text{end}$, $\text{year}$), saying that package is in service package pid from month start to end in year; and (c) business($\text{pnum}$, $\text{type}$, $\text{region}$) says that business number $\text{pnum}$ in region is of type, e.g., bank, hospital.

An access schema $\mathcal{A}_0$ over $\mathcal{R}_0$ includes:

- $\psi_1: \text{call}(\{\text{pnum}, \text{date}\}) \rightarrow \{\text{recnum}, \text{region}\}$, 500,
- $\psi_2: \text{package}(\{\text{pnum}, \text{year}\}) \rightarrow \{\text{pid}, \text{start}, \text{end}\}$, 12, and
Finding regions to access constraint of \( \psi \) that still accesses over 12 million tuples. One of the reasons is that \( \psi \)'s from relation \( \text{call} \) has at most 2000 \( \text{pnum} \)s, which is 1953 times faster, although it takes 187.8s, 96.13ms using the bounded plan, while a commercial DBMS can find exact answers to \( Q \) with bounded resources; (3) otherwise, it generates partially bounded plans and uses DBMS to compute exact answers (see Section 3).

Figure 1: Architecture of BEAS

### 3. The Architecture of BEAS

As shown in Fig. 1, BEAS consists of three major components: (1) offline service \text{AS Catalog} to manage access schema for different applications; and (2) online service \text{BE Query Planner} and \text{BE Plan Executor} to process SQL queries. It can be built on top of any commercial DBMS. \text{AS Catalog} consists of three modules itself.

(1) \text{Metadata module} maintains (a) access schema, and (b) statistics including the index size in a system table \text{as catalog}, for query plan generation and optimization.

Partial distinct tuples. It reduces redundancies introduced by irrelevant and duplicated attributes, and their inflation by joins. Hence bounded evaluation may substantially improve the scalability and efficiency of query evaluation.

While desirable, it is undecidable to determine whether an SQL query is boundedly evaluable under an access schema \[8\]. Nonetheless, we can still make practical use of bounded evaluation due to the existence of an effective syntax.

Feasibility Theorem \[5\]: Under an access schema \( \mathcal{A} \), there exists a class of queries covered by \( \mathcal{A} \) such that (1) an RA query \( Q \) is boundedly evaluable under \( \mathcal{A} \) if and only if there exists an RA query \( Q' \) that is covered by \( \mathcal{A} \) such that \( Q(D) = Q'(D) \) for all \( D \models \mathcal{A} \); and (2) \( Q \) is in \( \mathcal{PTIME} \) to decide whether \( Q \) is covered by \( \mathcal{A} \).

That is, \( Q \) is boundedly evaluable if and only if it can be rewritten into an equivalent \( Q' \) covered by \( \mathcal{A} \). In other words, covered queries make the core subclass of boundedly evaluable queries in relational algebra, without sacrificing their expressive power. This is along the same lines as the study of (undecidable) safe relational calculus queries \[1\].

\(\text{BEAS} \) extends Theorem 1 to SQL queries, and extends DBMS with a bounded evaluation functionality as follows:

(1) (2) it generates a bounded query plan and computes exact answers to \( Q \) within bounded resources; (3) otherwise, it generates partially bounded plans and uses DBMS to compute exact answers (see Section 3).

Algorithms for checking the bounded evaluability and for generating bounded query plans have been reported in \[5, 7\].

In fact, for queries that are not covered by \( \mathcal{A} \) (case (3) above), \( \text{BEAS} \) supports bounded approximation under access constraints, which guarantees a deterministic accuracy lower bound on approximate answers computed, and moreover, it accesses a bounded number of tuples in the entire process. These offer a solution to querying relations with bounded resources. To simplify the presentation, we focus on bounded evaluation for computing exact query answers in the demo.
We next present our demonstration plan, including settings and scenarios, as well as an industry application. The target audience of the demo includes anyone who is interested in query answering with bounded resources.
We have implemented BEAS on top of PostgreSQL 9.4.6. We have also created a demo portal as shown in Fig. 2, via which the audience will be able to interact with BEAS. It is deployed at a workstation with Xeon E3-1535M@2.9GHz CPU, 64GB of memory and 1.5TB of disk.

We will invite the audience to (1) examine the feasibility of bounded evaluation by interacting with BEAS. The audience will also (2) experience the performance of BEAS for exact SQL query answering, compared with commercial DBMS, using a telecommunication application as a testbed.

(1) A walk through. We visualize and demonstrate each major component of bounded evaluation underlying BEAS.

(a) Bounded evalability checking. As shown in Fig. 2(A), the audience will be invited to enter an SQL query \( Q \), select a dataset, pick an access schema \( \mathcal{A} \) discovered, and check whether \( Q \) is boundedly evaluable under \( \mathcal{A} \) using BE Checker. Users can also enter a budget on the amount of data to be accessed, and use BE Checker to find whether \( Q \) can be answered within the budget under \( \mathcal{A} \), without executing \( Q \).

(b) Bounded planning and optimization. As shown in Fig. 2(B), when \( Q \) is boundedly evaluable under \( \mathcal{A} \), the users will see a bounded query plan suggested by BE Plan Generator, in which each fetch operation is annotated with an upper bound on the amount of data to be fetched. The upper bound is deduced by reasoning about \( \mathcal{A} \).

If \( Q \) is not bound, BEAS picks a query plan \( \xi \) generated by PostgreSQL BE Plan Optimizer then makes \( \xi \) partially bounded by identifying bounded sub-queries of \( Q \) under \( \mathcal{A} \).

(c) Analysis. After a query plan is carried out by BE Plan Executor, a performance analysis is provided (Fig. 2(C)).

(d) Access schema management. As offline services, (i) the discovery module of BEAS takes as input a dataset \( D \), a set \( Q \) of query patterns, and a choice of the objective function; it discovers an access schema \( \mathcal{A} \) and register it by AS_Catalog (Fig. 2(D)). For instance, Figure 2(E) shows part of an access schema discovered. (ii) The maintenance module automatically updates \( \mathcal{A} \) in response to changes to \( D \) and queries. It also allows users to add or remove access constraints.

(2) Performance. We demonstrate how BEAS works in practice using a commercial benchmark from a telecommunication company (name withheld), and compare its performance with PostgreSQL, MySQL and MariaDB.

Efficiency. The users are invited to interact with BEAS, pick built-in queries or enter their own queries, and examine the effectiveness of bounded evaluation. For instance, for query \( Q \) of Example 2 on a TLC dataset \( D \) of 20GB, a snapshot of the BEAS performance analyzer is given in Fig. 3, which shows that BEAS is 1953, 6562 and 5135 times faster than PostgreSQL, MySQL and MariaDB, respectively. It details (a) the overall execution time, acceleration ratio compared to commercial DBMS, the total number of tuples fetched and the number of access constraints employed; and (b) a breakdown of the cost to each individual operation in the query plan, compared to its counterpart in plans generated by commercial DBMS. It illustrates why BEAS works better.

Scalability. The audience will also witness the scalability of BEAS by scaling up the datasets. Figure 4 shows the evaluation time of \( Q \) of Example 2 with BEAS, PostgreSQL, MySQL and MariaDB when varying TLC from 1GB to 200GB. One can see that BEAS consistently takes about 1s when \( D \) varies, and is hence “scale-independent”. In contrast, PostgreSQL, MySQL and MariaDB grow to 1932s, 6187s and 5243s, respectively, if we allow them to run to completion.

5. References