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Compiler Fuzzing through Deep Learning

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1 INTRODUCTION

Compilers should produce correct code for valid inputs, and meaningful errors for invalid inputs. Failure to do so can hinder software development or even cause catastrophic runtime errors. Still, properly testing compilers is hard. Modern optimizing compilers are large and complex programs, and their input space is huge. Hand designed suites of test programs, while important, are inadequate for covering such a large space and will not touch all parts of the compiler.

Random test case generation — fuzzing — is a well established and effective method for identifying compiler bugs [6, 7, 16]. When fuzzing, randomly generated valid or semi-valid inputs are fed to the compiler. Any kind of unexpected behavior, including crashes, freezes, or wrong binaries, indicates a compiler bug. While crashes and freezes in the compiler are easy to detect, determining that binaries are correctly compiled is not generally possible without either developer provided validation for the particular program’s behavior or a gold standard compiler from which to create reference outputs. In the absence of those, Differential Testing [22] can be used. The generated code and a set of inputs form a test case which is compiled and executed on multiple testbeds. If the test case should have deterministic behavior, but the output differs between testbeds, then a bug has been discovered.

Compiler fuzzing requires efficiently generating test cases that trigger compiler bugs. The state-of-the-art approach, CSmith [32], generates large random programs by defining and sampling a probabilistic grammar which covers a subset of the C programming language. Through this grammar, CSmith ensures that the generated code easily passes the compiler front-end and stresses the most complex part of the compiler, the middle-end. Complex static and dynamic analyses make sure that programs are free from undefined behavior. The programs are then differentially tested.

While CSmith has been successfully used to identify hundreds of bugs in compilers, it and similar approaches have a significant drawback. They represent a huge undertaking and require a thorough understanding of the target programming language. CSmith was developed over the course of years, and consists of over 41k lines of handwritten C++ code. By tightly coupling the generation logic with the target programming language, each feature of the grammar must be painstakingly and expertly engineered for each new target language. For example, lifting CSmith from C to OpenCL [20] — a superficially simple task — took 9 months and an additional 8k lines of code. Given the difficulty of defining a new grammar, typically only a subset of the language is implemented.

What we propose is a fast, effective, and low effort approach to the generation of random programs for compiler fuzzing. Our methodology uses recent advances in deep learning to automatically
construct probabilistic models of how humans write code, instead of painstakingly defining a grammar to the same end. By training a deep neural network on a corpus of handwritten code, it is able to infer both the syntax and semantics of the programming language and the common constructs and patterns. Our approach essentially frames the generation of random programs as a language modeling problem. This greatly simplifies and accelerates the process. The expressiveness of the generated programs is limited only by what is contained in the corpus, not the developer’s expertise or available time. Such a corpus can readily be assembled from open source repositories.

In this work we primarily target OpenCL, an open standard for programming heterogeneous systems, though our approach is largely language agnostic. We chose OpenCL for three reasons: it is an emerging standard with the challenging promise of functional portability across a diverse range of heterogeneous hardware; OpenCL is compiled “online”, meaning that even compiler crashes and freezes may not be discovered until a product is deployed to customers; and there is already a hand written random program generator for the language to compare against. We provide preliminary results supporting DeepSmith’s potential for multi-lingual compiler fuzzing.

We make the following contributions:

- a novel, automatic, and fast approach for the generation of expressive random programs for compiler fuzzing. We infer programming language syntax, structure, and use from real-world examples, not through an expert-defined grammar. Our system needs two orders of magnitude less code than the state-of-the-art, and takes less than a day to train;
- we discover a similar number of bugs as the state-of-the-art, but also find bugs which prior work cannot, covering more components of the compiler;
- in modeling real handwritten code, our test cases are more interpretable than other approaches. Average test case size is two orders of magnitude smaller than state-of-the-art, without any expensive reduction process.

2 DEEPSMITH

DeepSmith\(^1\) is our open source framework for compiler fuzzing. Figure 1 provides a high-level overview. In this work we target OpenCL, though the approach is language agnostic. This section describes the three key components: a generative model for random programs, a test harness, and voting heuristics for differential testing.

2.1 Generative Model

Generating test cases for compilers is hard because their inputs are highly structured. Producing text with the right structure requires expert knowledge and a significant engineering effort, which has to be repeated from scratch for each new language. Instead, we treat the problem as an unsupervised machine learning task, employing state-of-the-art deep learning techniques to build models for how humans write programs. Our approach is inspired by breakthrough results in modeling challenging and high dimensional datasets through unsupervised learning [4, 27, 28]. Contrary to existing tools, our approach does not require expert knowledge of the target language and is only a few hundred lines of code.

*Handwritten Programs.* The generative model needs to be trained on a *seed corpus* of example programs. We automated the assembly of this corpus by mining 10k OpenCL kernels from open source repositories on GitHub. We used an *oracle compiler* (LLVM 3.9) to statically check the source files, discarding files that are not well-formed. The main purpose of this step is to remove the need to manually check that each file selected from GitHub does indeed contain OpenCL. A downside is that any training candidate which triggers a bug in the LLVM 3.9’s front end will not be included. However, this did not prevent our system from uncovering errors in that compiler (Section 4.4).

This corpus, exceeding one million lines of code, is used as a representative sample of OpenCL code from which a generative model can be derived.

*Encoder.* The textual representation of program codes must be encoded as numeric sequences for feeding as input to the machine learning model. Prior machine learning works have used character-level encodings, token-level encodings, or fixed length feature vectors. We extend the hybrid character/token-level encoding of [9], in which a programming language’s keywords and common names are treated as individual tokens while the rest of the text is encoded on a character-level basis. This approach hits a balance between compressing the input text and keeping the number of tokens in the vocabulary low.

\(^1\)DeepSmith available at: https://chriscummins.cc/deepsmith
We additionally employed semantic-preserving transformations to simplify the training programs. First, each source file is preprocessed to expand macros and remove conditional compilation and comments. Then, all user-declared identifiers are renamed using an arbitrary, but consistent pattern based on their order of declaration: \{a, b, c, \ldots, aa, ab, ac, \ldots\} for variables and \{A, B, C, \ldots, AA, AB, AC, \ldots\} for functions. This ensures a consistent naming convention, without modifying program behavior. Finally, a uniform code style is enforced to ensure consistent use of braces, parentheses, and white space. These rewriting simplifications give more opportunities for the model to learn the structure and deeper aspects of the language and speed up the learning. On the other hand, some bugs in the preprocessor or front-end might no longer be discoverable. We reason that this is an acceptable trade-off. For languages where the corpus can be many orders of magnitude larger, for example, C or Java, models may be generated without these modifications.

Neural Network. We use the Long Short-Term Memory (LSTM) architecture of Recurrent Neural Network to model program code [12]. In the LSTM architecture activations are learned with respect not just to their current inputs but to previous inputs in a sequence. In our case, this allows modeling the probability of a token appearing in the text given a history of previously seen tokens. Unlike previous recurrent networks, LSTMs employ a forget gate with a linear activation function, allowing them to avoid the vanishing gradients problem [24]. This makes them effective at learning complex relationships over long sequences [21] which is important for modeling program code. Our LSTM networks model the vocabulary distribution over the encoded corpus. After initial experiments using different model parameters, we found that a two layer LSTM network of 512 nodes per layer provided a good trade-off between the fidelity of the learned distribution and the size of the network, which limits the rate of training and inference. The network is trained using Stochastic Gradient Descent for 50 epochs, with an initial learning rate of 0.002 and decaying by 5% every epoch. Training the model on the OpenCL corpus took 12 hours using a single NVIDIA Tesla P40. We provided the model with no prior knowledge of the structure or syntax of a programming language.

Program Generation. The trained network is sampled to generate new programs. The model is seeded with the start of a kernel (identified in OpenCL using the keywords kernel void), and sampled token-by-token. A “bracket depth” counter is incremented or decremented upon production of \{ or \} tokens respectively, so that the end of the kernel can be detected and halted. The generated sequence of tokens is then decoded back to text and used for compiler testing.

2.2 Test Harness
OpenCL is an embedded compute kernel language, requiring host code to compile, execute, and transfer data between the host and device. For the purpose of compiler fuzzing, this requires a test harness to run the generated OpenCL programs. At first, we used the test harness of CLSmith. The harness assumes a kernel with no input and a ulong buffer as its single argument where the result is written. Only 0.2% of the GitHub kernels share this structure.

We desired a more flexible harness so as to test a more expressive range of programs, capable of supporting multi-argument kernels and generating data to use as inputs.

We developed a harness which first determines the expected arguments from the function prototype and generates host data for them. At the moment, we support scalars and arrays of all OpenCL primitive and vector types. For a kernel execution across \(n\) threads, buffers of size \(n\) are allocated for pointer arguments and populated with values \([1\ldots n]\); scalar inputs are given value \(n\), since we observe that most kernels use these for specifying buffer sizes.

The training programs from which the generative model is created are real programs, and as such do not share the argument type restrictions. The model, therefore, may generate correct programs for which our driver cannot create example inputs. In this case, a “compile-only” stub is used, which only compiles the kernel, without generating input data or executing the compiled kernel.

Unlike the generative model, this test harness is language-specific and the design stems from domain knowledge. Still, it is a relatively simple procedure, consisting of a few hundred lines of Python.

Test Harness Output Classes. Executing a test case on a testbed leads to one of seven possible outcomes, illustrated in Figure 2. A build failure occurs when online compilation of the OpenCL kernel fails, usually accompanied by an error diagnostic. A build crash or build timeout outcome occurs if the compiler crashes or fails to produce a binary within 60 seconds, respectively. For compile-only test cases, a pass is achieved if the compiler produces a binary. For test cases in which the kernel is executed, kernel execution leads to one of three potential outcomes: runtime crash if the program crashes, timeout if the kernel fails to terminate within 60 seconds, or pass if the kernel terminates gracefully and computes an output.

![Test Harness Output Classes](image-url)
2.3 Voting Heuristics for Differential Testing

We employ established Differential Testing methodologies to expose compiler defects. As in prior work, voting on the output of programs across compilers has been used to circumvent the oracle problem and detect miscompilations [22]. However, we extend this approach to describe not only miscompilations, but also anomalous build failures and crashes.

When evaluating the outcomes of test cases, build crash (bc) and build timeout (bto) outcomes are of immediate interest, indicative of erroneous compiler behavior (examples may be found in Section 4.1). For all other outcomes, differential tests are required to confirm anomalous behavior. We look for test cases where there is a majority outcome – i.e. for which some fraction of the testbeds behave the same – but some testbed deviates. We use the presence of the majority increasing the likelihood that there is a ‘correct’ behavior for the test case. In this work, we choose the majority fraction to be $\frac{\lceil \frac{2}{3} n \rceil}{n}$, where $n$ is the number of testbeds.

An anomalous build failure (abf) or anomalous runtime crash (arc) occurs if, for a given test case, the majority of testbeds execute successfully, and a testbed yields a compilation error or runtime crash. An anomalous wrong-output (awo) occurs if, for a given test case, the majority of testbeds execute successfully, producing the same output values, and a testbed yields a result which differs from this majority output. Anomalous wrong-output results are indicative of miscompilations, a particularly hard to detect class of bug in which the compiler silently emits wrong code. CSmith is designed specifically to target this class of bug.

False Positives for Anomalous Runtime Behavior. Generated programs may contain undefined or non-deterministic behavior which will incorrectly be labeled as anomalous. CSmith circumvents this problem by performing complex analyses during generation so as to minimize the chance of producing programs with undefined behavior. Although similar analyses could be created as filters for our system, we take a simpler approach, filtering only the few types of non-deterministic behavior we have actually observed to happen in practice.

We filter data races, out-of-bounds and uninitialized accesses with GPUverify [2] and Oclgrind [26]. Some compiler warnings provide strong indication of non-deterministic behavior (e.g. comparison between pointer and integer) – we check for these warnings and filter accordingly.

Floating point operations in OpenCL can be imprecise, so code can produce different output on different testbeds. For this reason, CSmith and CLSmith do not support floating point operations. DeepSmith allows floating point operations but since it cannot apply differential testing on the outputs, it can detect all results except for the anomalous wrong-output results.

The last type of undefined behavior we observed comes from division by zero and related mathematical functions which require non-zero values. We apply a simple detection and filtering heuristic – we change the input values and check to see if the output remains anomalous. While theoretically insufficient, in practice we found that no false positives remained.

3 EXPERIMENTAL SETUP

In this section we describe the experimental parameters used.

3.1 OpenCL Systems

We conducted testing of 10 OpenCL systems, summarized in Table 1. We covered a broad range of hardware: 3 GPUs, 4 CPUs, a co-processor, and an emulator. 7 of the compilers tested are commercial products, 3 of them are open source. Our suite of systems includes both combinations of different drivers for the same device, and different devices using the same driver.

3.2 Testbeds

For each OpenCL system, we create two testbeds. In the first, the compiler is run with optimizations disabled. In the second, optimizations are enabled. Each testbed is then a triple, consisting of <device, driver, is_optimized> settings. This mechanism gives 20 testbeds to evaluate.

3.3 Test Cases

For each generated program we create inputs as described in Section 2.2. In addition, we need to choose the number of threads to use. We generate two test cases, one using one thread, the other using 2048 threads. A test case is then a triple, consisting of <program, inputs, threads> settings.

3.4 Bug Search Time Allowance

We compare both our fuzzer and CLSmith. We allow both to run for 48 hours on each of the 20 testbeds. CLSmith used its default configuration. The total runtime for a test case consists of the generation and execution time.

Table 1: OpenCL systems and the number of bug reports submitted to date (22% of which have been fixed, the remainder are pending). For each system, two testbeds are created, one with compiler optimizations, the other without.
We report on the results of DeepSmith testing of the 10 OpenCL
verbatim, with only minor formatting changes applied to save space.

1. **Testbeds**

   Denote optimizations on, off, or either, respectively. A

   OpenCL system number from Table 1, suffixed with

   +

   No test case reduction, either manual or automatic, was needed.

   +

   A moving target.

   Features and regressions ensures that compiler validation remains

   compilers are becoming more robust, yet the introduction of new

   as a metric of compiler robustness. We find that progress is good,

   quantitative analysis of compiler robustness over time, using the

   for only a fraction of the engineering effort. Finally, we provide a

   identify a broad range of defects, many of which CLSmith cannot,

   DeepSmith is able to

   compute the wrong result. To date, we have submitted 67 bug re-

   every compiler generated programs which either crash or silently

   bugs in all the compilers we tested — every compiler crashed, and

   systems from Table 1, in which each ran for 48 hours. We found

   4 EVALUATION

   We report on the results of DeepSmith testing of the 10 OpenCL

   systems from Table 1, in which each ran for 48 hours. We found

   bugs in all the compilers we tested — every compiler crashed, and

   every compiler generated programs which either crash or silently

   compute the wrong result. To date, we have submitted 67 bug re-

   ports to compiler vendors. We first provide a qualitative analysis

   of compile-time and runtime defects found, followed by a quanti-

   tative comparison of our approach against the state-of-the-art in

   OpenCL compiler fuzzing — CLSmith [20]. DeepSmith is able to

   identify a broad range of defects, many of which CLSmith cannot,

   for only a fraction of the engineering effort. Finally, we provide a

   quantitative analysis of compiler robustness over time, using the

   compiler crash rate of every LLVM release in the past two years as

   a metric of compiler robustness. We find that progress is good,

   compilers are becoming more robust, yet the introduction of new

   features and regressions ensures that compiler validation remains

   a moving target.

   Unless stated otherwise, DeepSmith code listings are presented

   verbatim, with only minor formatting changes applied to save space.

   No test case reduction, either manual or automatic, was needed.

   For the remainder of the paper we identify testbeds using the

   OpenCL system number from Table 1, suffixed with +, −, or ± to

   denote optimizations on, off, or either, respectively.

   ![Figure 3: Example kernels which crash compilers.](image)

   **Figure 4: Example codes which crash parsers.**

   4.1 Compile-time Defects

   OpenCL is typically compiled online, which amplifies the signifi-

   cance of detecting compile-time defects, as they may not be discov-

   ered until code has been shipped to customers. We found numerous

   cases where DeepSmith kernels trigger a crash in the compiler

   (and as a result, the host process), or cause the compiler to loop

   indefinitely. In the testing time allotted we have identified 199 test

   cases which trigger unreachable code failures, triggered 31 different

   compiler assertions, and produced 114 distinct stack traces from

   other compiler crashes.

   Semantic Analysis Failures. Compilers should produce meaningful

   diagnostics when inputs are invalid, yet we discovered dozens

   of compiler defects attributable to improper or missing error han-

   dling. Many generation and mutation based approaches to com-

   pile validation have focused solely on testing under valid inputs.

   As such, this class of bugs may go undiscovered. We believe that

   our approach contributes a significant improvement to generat-

   ing plausibly-erroneous code over prior random-enumeration ap-

   proaches.

   The use of undeclared identifiers is a core error diagnostic which

   one would expect to be robust in a mature compiler. DeepSmith

   discovered cases in which the presence of undeclared identifiers

   causes the Testbeds 10± compiler to crash. For example, the unde-

  clared identifier c in Figure 3a raises an assertion during semantic

   analysis of the AST when used as an array index.

   Type errors were an occasional cause of compile-time defect.

   Figure 3b induces a crash in NVIDIA compilers due to an implicit

   conversion between global to constant address qualifiers. Worse,

   we found that Testbeds 3± would loop indefinitely on some kernels

   containing implicit conversions from a pointer to an integer, as

   shown in Figure 5a. While spinning, the compiler would utilize 100%

   of the CPU and consume an increasing amount of host memory

   until the entire system memory is depleted and the process crashes.

   Occasionally, incorrect program semantics will remain unde-

   tected until late in the compilation process. Both Figures 3c and 3d

   pass the type checker and semantic analysis, but trigger compiler

   assertions during code generation.

   An interesting yet unintended byproduct of having trained Deep-

   Smith on thousands of real world examples is that the model learned

   to occasionally generate compiler-specific code, such as invoking

   compiler builtins. We found the quality of error handling on these

   builtins to vary wildly. For example, Figure 3e silently crashes 6

   of the 10 compilers, which, to the best of our knowledge, makes

   ![Figure 5: Example kernels which crash parsers.](image)
which we have reported to Intel. Each reduction took around 6
4
Testbeds
1
(b) Testbeds
4
1
(c) Testbeds
4
1
Testbed
3
1
(a) Testbeds
3
1
Testbed
2
1
Testbed
1

Testbeds
3± loop indefinitely, leaking memory until the entire system mem-
ory is depleted and the process crashes.

1 kernel void A(global int*: a) {
2 int b = get_global_id(0);
3 a[b] = (6 * 32) + 4 * (32 / 32) + a;
4 }

(a) Testbeds 3± loop indefinitely, leaking memory until the entire system memory is depleted and the process crashes.

1 kernel void A(global float* a, global float* b, global float* c) {
2 int d, e, f;
3 d = get_local_id(0);
4 for (int g = 0; g < 100000; g++)
5 barrier[1];
6 }

(b) Testbed 1+ hangs during optimization of kernels with large loop bounds. Testbeds 1± and 2± compile in under 1 second.

1 kernel void A(global int*: a) {
2 int b = get_global_id(0);
3 while (b < 512) { }
4 }

(c) Testbeds 4±, 5±, 6±, 7+ hang during optimization of kernels with non-
terminating loops.

1 kernel void A(global unsigned char*: a, global unsigned b) {
2 for (int i = 0; i < 42; i++)
3 barrier[1];
4 }

(d) Testbeds 7± loops indefinitely, consuming 100% CPU usage.

Figure 5: Example kernels which hang compilers.

DeepSmith the first random program generator to induce a defect through exploiting compiler-specific functionality.

Parser Failures. Parser development is a mature and well under-
stood practice. We uncovered parser errors in several compilers. Each of the code samples in Figure 4 induce crash errors during parsing of compound statements in both Testbeds 5± and 7±. For space, we have hand-reduced the listings to minimal code samples, which we have reported to Intel. Each reduction took around 6 edit-compile steps, taking less than 10 minutes. In total, we have generated 100 distinct programs which crash compilers during parsing.

Compiler Hangs. As expected, some compile-time defects are optimization sensitive. Testbed 1+ hangs on large loop bounds, shown in Figure 5b. All commercial Intel compilers we tested hang during optimization of non-terminating loops (Figure 5c).

Testbeds 7± loop indefinitely during compilation of the simple kernel in Figure 5d.

Other errors. Some compilers are more permissive than others. Testbeds 4±, 6±, 9± reject out-of-range literal values e.g. int i = 0xfffffffffffffffffffffff, whilst Testbeds 3±, 5±, 7±, 8±, and 10± interpret the literal as an unsigned long long and implicitly cast to an integer value of -1. Testbeds 1±, 2± emit no warning.

Testbeds 1±, 2±, 3± rejected address space qualifiers on automatic variables, where all other testbeds successfully compiled and executed.

4.2 Runtime Defects
Prior work on compiler test case generation has focused on extensive stress-testing of compiler middle-ends to uncover uncompilations [6]. Csmith, and by extension, CLSmith, specifically targets this class of bugs. Grammar based enumeration is highly effective at this task, yet is bounded by the expressiveness of the grammar. Here we provide examples of bugs which cannot currently be discovered by CLSmith.

Thread-dependent Flow Control. A common pattern in OpenCL is to obtain the thread identity, often as an int, and to compare this against some fixed value to determine whether or not to complete a unit of work (46% of OpenCL kernels on GitHub use this (tid → int,
if (tid < ...) {...}) (pattern). DeepSmith, having modeled the frequency with which this pattern occurs in real handwritten code, generates many permutations of this pattern. And in doing so, exposed a bug in the optimizer of Testbeds 4+ and 6+ which causes the if branch in Figure 6a to be erroneously executed when the kernel is compiled with optimizations enabled. We have reported this issue to Intel. CLSmith does not permit the thread identity to modify control flow, rendering such productions impossible.

Figure 6b shows a simple program in which thread identity determines the program output. We found that this test case would sporadically crash Testbeds 10±, an OpenCL device simulator and debugger. Upon reporting to the developers, the underlying cause was quickly diagnosed as a race condition in switch statement evaluation, and fixed within a week.

Kernel Inputs. CLSmith kernels accept a single buffer parameter into which each thread computes its result. This fixed prototype limits the ability to detect bugs which depend on input arguments. Figure 6c exposes a bug of this type. Testbeds 3± will silently miscompile ternary operators when the ternary operands consist of values stored in multiple different global buffers. CLSmith, with its fixed single input prototype, is unable to discover this bug.

Latent Compile-time Defects. Sometimes, invalid compiler inputs may go undetected, leading to runtime defects only upon program execution. Since CLSmith enumerates only well-formed programs, this class of bugs cannot be discovered.

Figure 6d exposes a bug in which a kernel containing an undefined symbol will successfully compile without warning on Testbeds 8±, then crash the program when attempting to run the kernel. This issue has been reported to the developers and fixed.

4.3 Comparison to State-of-the-art

In this section, we provide a quantitative comparison of the bug-finding capabilities of DeepSmith and CLSmith.

Results Overview. Table 2 shows the results of 48 hours of consecutive testing for all Testbeds. An average of 15k CLSmith and 91k DeepSmith test cases were evaluated on each Testbed, taking 12.1s and 1.90s per test case respectively. There are three significant factors providing the sixfold increase in testing throughput achieved by DeepSmith over CLSmith: test cases are faster to generate, test cases are less likely to timeout (execute for 60 seconds without termination), and the test cases which do not timeout execute faster.

Figure 7a shows the generation and execution times of DeepSmith and CLSmith test cases, excluding timeouts. DeepSmith generation time grows linearly with program length, and is on average 2.45× faster than CLSmith. Test case execution is on average 4.46× faster than CLSmith.

The optimization level generally does not affect testing throughput significantly, with the exception of Testbed 7+. Optimization of large structs is expensive on Testbed 7+, and CLSmith test cases use global structs extensively. This is a known issue — in [20] the authors omit large-scale testing on this device for this reason.

If timeouts are included then the performance improvement of DeepSmith is 6.5x with the execution times being 11x faster. However, this number grows as we change the arbitrary timeout threshold, so for fairness we have chosen to exclude it.

Latent Compile-time Defects. Sometimes, invalid compiler inputs may go undetected, leading to runtime defects only upon program execution. Since CLSmith enumerates only well-formed programs, this class of bugs cannot be discovered.

Comparison of Test Cases. The average CLSmith program is 1189 lines long (excluding headers). CLSmith test cases require reduction in order to expose the underlying bug. An automated approach to OpenCL test case reduction is presented in [25], though it requires on average 100 minutes for each test case using a parallelized implementation (and over 6 hours if this parallelization is not available); the authors also suggest a final manual pass after automated reduction. In contrast, DeepSmith learned to program from humans, and humans do not typically write such large kernel functions. The average DeepSmith kernel is 20 lines long, which is interpretable without reduction, either manual or automatic.

Comparison of Results. Both testing systems found anomalous results of all types. In 48 hours of testing, CLSmith discovered compile-time crashes (be) in 8 of the 20 testbeds, DeepSmith crashed all of them. DeepSmith triggered 31 distinct compiler assertions, CLSmith 2. Both of the assertions triggered by CLSmith were also triggered by DeepSmith. DeepSmith also triggered 3 distinct unreachable! compile-time crashes, CLSmith triggered 0. The ratio of build failures is higher in the token-level generation of DeepSmith (51%) than the grammar-based generation of CLSmith (26%).

The Intel CPU Testbeds (4±, 5±, 6±, and 7±) would occasionally emit a stack trace upon crashing, identifying the failure point in a specific compiler pass. CLSmith triggered such crashes in 4 distinct passes. DeepSmith triggered crashes in 10 distinct passes, including 3 of the 4 in which CLSmith did. Figure 8 provides examples. Many of these crashes are optimization sensitive, and are more likely to occur when optimizations are enabled. CLSmith was able to induce a crash in only one of the Intel testbeds with optimizations disabled. DeepSmith crashed all of the compilers with both optimizations enabled and disabled.
Table 2: Results from 48 hours of testing using CLSmith and DeepSmith. System #. as per Table 1.

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<th>±</th>
<th>bc</th>
<th>bto</th>
<th>abf</th>
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<th>awo</th>
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CLSmith produced many bto results across 13 Testbeds. Given the large kernel size, it is unclear how many of those are infinite loops or simply a result of slow compilation of large kernels. The average size of CLSmith bto kernels is 1558 lines. Automated test case reduction — in which thousands of permutations of a program are executed — may be prohibitively expensive for test cases with very long runtimes. DeepSmith produced bto results across 11 Testbeds and with an average kernel size of 9 lines, allowing for rapid identification of the underlying problem.

The integrated GPU Testbeds (3±) frequently failed to compile CLSmith kernels, resulting in over 10k bc and bto results. Of the build crashes, 68% failed silently, and the remainder were caused by the same two compiler assertions for which DeepSmith generated 4 line test cases, shown in Figure 9. DeepSmith also triggered silent build crashes in Testbeds 3±, and a further 8 distinct compiler assertions.

The 4719 abf results for CLSmith on Testbeds 4± and 6± are all a result of compilers rejecting empty declarations, (e.g. int;) which CLSmith occasionally emits. DeepSmith also generated these statements, but with a much lower probability, given that it is an unusual construct (0.6% of test cases, versus 7.0% of CLSmith test cases).

ComputeAorta (Testbeds 9±) defers kernel compilation so that it can perform optimizations dependent on runtime parameters. This may contribute to the relatively large number of arc results and few bc results of Testbeds 9±. Only DeepSmith was able to expose compile-time defects in this compiler.

Over the course of testing, a combined 3.4 × 10^8 lines of CLSmith code was evaluated, compared to 3.8 × 10^8 lines of DeepSmith code. This provides CLSmith a greater potential to trigger mis-compilations. CLSmith generated 33 programs with anomalous wrong-outputs. DeepSmith generated 30.

4.4 Compiler Stability Over Time

The Clang front-end to LLVM supports OpenCL, and is commonly used in OpenCL drivers. This in turn causes Clang-related defects to potentially affect multiple compilers, for example the one in Figure 3e. To evaluate the impact of Clang, we used debug+assert builds of every LLVM release in the past 24 months and processed 75,000 DeepSmith kernels through the Clang front-end (this includes the lexer, parser, and type checker, but not code generation).

Figure 10 shows that the crash rate of the Clang front-end is, for the most part, steadily decreasing over time. The number of failing compiler crashes decreased tenfold between 3.6.2 and 5.0.0. Table 3 shows the 7 distinct assertions triggered during this experiment. Assertion 1 (Uncorrected type!) is raised on all compiler versions — see Figure 3a for an example. The overall rate at which the assertion is triggered has decreased markedly, although there are slight increases between some releases. Notably, the current development trunk has the second lowest crash rate, but is joint first in terms of the number of unique assertions. Assertions 3 (Addr == 0 || hasTargetSpecificAddressSpace()) and 4 (isScalarType()) were triggered by some kernels in the development trunk but not under any prior release. We have submitted bug reports for each of the three assertions triggered in the development trunk, as well as for two distinct unreachablees.

The results emphasize that compiler validation is a moving target. Every change and feature addition has the potential to introduce regressions or new failure cases. Since LLVM will not release unless their compiler passes their own extensive test suites, this also reinforces the case for compiler fuzzing. We believe our approach provides an effective means for the generation of such fuzzers, at a fraction of the cost of existing techniques.
Compiler Fuzzing through Deep Learning

Figure 8: Example kernels which crash Intel compiler passes.

```c
1 kernel void A() {
2   while (!true) {
3     barrier();
4   }
}
```

(a) Post-Dominance Frontier Construction pass.

```c
1 kernel void A(global float a, global float b, int c) {
2   for (int d = 0; d < c; d++)
3     a[d] = b / b * get_global_id(0);
4   barrier();
5     a[get_global_id(0)] = a[get_global_id(0)] > b;
6 }
```

(b) Simplify the CFG pass.

```c
1 kernel void A(int a) {
2   int b = get_global_id(0);
3   while (b < a)
4     if (a[i] < 0)
5       a[i] = a[i] + b / b + get_local_id(0);
6 }
```

(c) Predicator pass.

```c
1 kernel void A(int a, global int b) {
2   global float c, const int d (;
3     for (unsigned int e = get_global_id(0);
4       e < d; e += get_global_size(0))
5       for (unsigned f = 0; f < d; f++)
6         e += a[f];
7 }
```

(d) Combine redundant instructions pass.

```c
1 kernel void A(int a, global int b) {
2   if (a[i] < 0)
3     a[i] = 0;
4 }
```

(e) PrepareKernelArgs pass.

```c
1 kernel void A() {
2   local float a; A(a);
3 }
```

(f) Add SPIR related module scope metadata pass.

```c
1 kernel void A() {
2   local int a[10];
3   local int b[16][16];
4   a[1024 + (2 * get_local_id(1) +
5     get_local_id(0)) + get_local_id(0)] = 6;
6   barrier(b);)
7 }
```

(g) Intel OpenCL RemoveDuplicationBarrier pass.

```c
1 kernel void A(global half* a) {
2   int b = get_global_id(0);
3   a[b] = b * b;
4 }
```

(h) X86 DAG→DAG Instruction Selection pass.

Figure 9: Example kernels which trigger compiler assertions which both CLSmith and DeepSmith exposed.

```c
1 kernel void A(global int* a, global int* b, global int* c) {
2   a[get_global_id(0)] = a[get_global_id(0)] > b;
3 }
```

(a) Assertion storing/loading pointers only support private array.

```c
1 kernel void A() {
2   global int* a = ((void*)0);
3   b[0] = a;
4 }
```

(b) Assertion iter != pointerOrigMap.end()
with existing code, it inserts or deletes statements that will not be
triggered. Generally expressed as a probability of correctness.

Program Generation. In the foundational work on differential
testing for compilers, McKeeman et al. present generators capable of
equating programs of a range of qualities, from random ASCII
sequences to C model conforming programs [22]. Subsequent works
have presented increasingly complex generators which improve in
some metric of interest, generally expressiveness or probability of
the chances of discovering vulnerabilities in scripting language engines.

Program Mutation. Equivalence Modulo Inputs (EMI) testing [19,
29] follows a different approach to test case generation. Starting
with existing code, it inserts or deletes statements that will not be
executed, so functionality should remain the same. If it is affected,
it is due to a compiler bug. While a powerful technique able to
find hard to detect bugs, it relies on having a very large number
of programs to mutate. As such, it still requires an external code
generator. Similarly to CSMith, EMI favors very long test programs.
LangFuzz [13] also uses mutation but does this by inserting code
segments which have previously exposed bugs. This increases the
chances of discovering vulnerabilities in scripting language engines.

5 RELATED WORK

The random generation of test cases is a well established approach
to the compiler validation problem. Prior approaches are surveyed in
[3, 16] and empirically contrasted in [6]. The main question of
interest is in how to efficiently generate codes which trigger bugs.

There are two main approaches: program generation, where inputs
are synthesized from scratch; and program mutation, where existing
codes are modified so as to identify anomalous behavior.

Program Generation. In the foundational work on differential
testing for compilers, McKeeman et al. present generators capable of
equating programs of a range of qualities, from random ASCII
sequences to C model conforming programs [22]. Subsequent works
have presented increasingly complex generators which improve in
some metric of interest, generally expressiveness or probability of
correctness. CSMith [32] is a widely known and effective generator
which enumerates programs by pairing infrequently combined
language features. In doing so, it produces correct programs with
clearly defined behavior but very unlikely functionality, increasing
the chances of triggering a bug. Achieving this required extensive
eering work, most of it not portable across languages, and
ignoring some language features. Subsequent generators influenced
by CSMith, like Orange3 [23], focus on features and bug types
beyond the scope of CSMith, arithmetic bugs in the case of Orange3.

Program Mutation. Equivalence Modulo Inputs (EMI) testing [19,
29] follows a different approach to test case generation. Starting
with existing code, it inserts or deletes statements that will not be
executed, so functionality should remain the same. If it is affected,
ACKNOWLEDGEMENTS

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REFERENCES